

Analytics and Visualization of Big Data

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Lecture 10: Similarity of Sets (LSH)



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Spring 13

Applications of Near-Neighbor Search

- **Jaccard Similarity of Sets**
- **Similarity of Documents**
- **Collaborative Filtering**

Locality-Sensitive Hashing

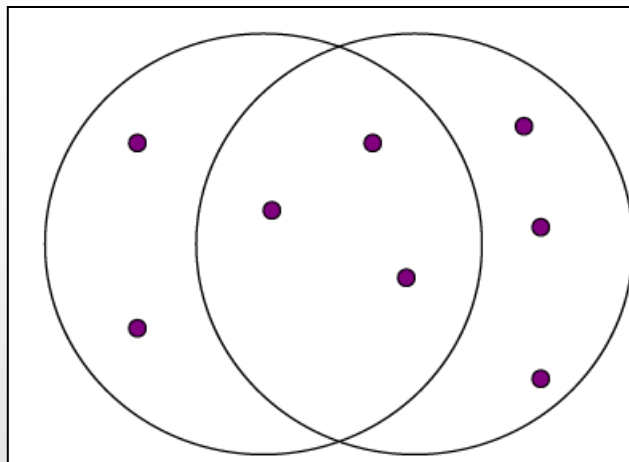
- **Shingling**
- **Minhashing**
- **LSH for Minhash Signatures**
- **Combining the Techniques**



How do we Define Similarity?

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- Typically, we want to have items that have common features → we use this to say there are similar ☺
- The **Jaccard Similarity** of two sets is:
 - $Sim(C_1, C_2) = |C_1 \cap C_2| / |C_1 \cup C_2|$
- The **Jaccard Distance** between sets is 1 minus their Jaccard similarity: $d(C_1, C_2) = 1 - |C_1 \cap C_2| / |C_1 \cup C_2|$



3 in intersection

8 in union

Jaccard similarity = $3/8$

Jaccard distance = $5/8$

How does this Concept Relate to “Big Data” Analytics?

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- **Goal:** Finding textually similar documents in a collection of news articles or web pages
 - Character-level similarity vs. similar meaning?
- Two levels of **similarity**:
 - **Exactness:** Easy, character-by-character comparison
 - **Near duplicates:** More involved; topic of today's class
- Typical applications in **Big Data Analytics**:
 - Plagiarism detection
 - Articles from the same source
 - Collaborative filtering



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Locality-Sensitive Hashing

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- Combining the Techniques



Problem Statement:

- Given a large number (N in the millions or billions) of text documents, find pairs that are “near duplicates”

Issues:

- Many small pieces of one doc can appear out of order in another
- Too many docs to compare all pairs
- Docs are so large or so many that they cannot fit in main memory

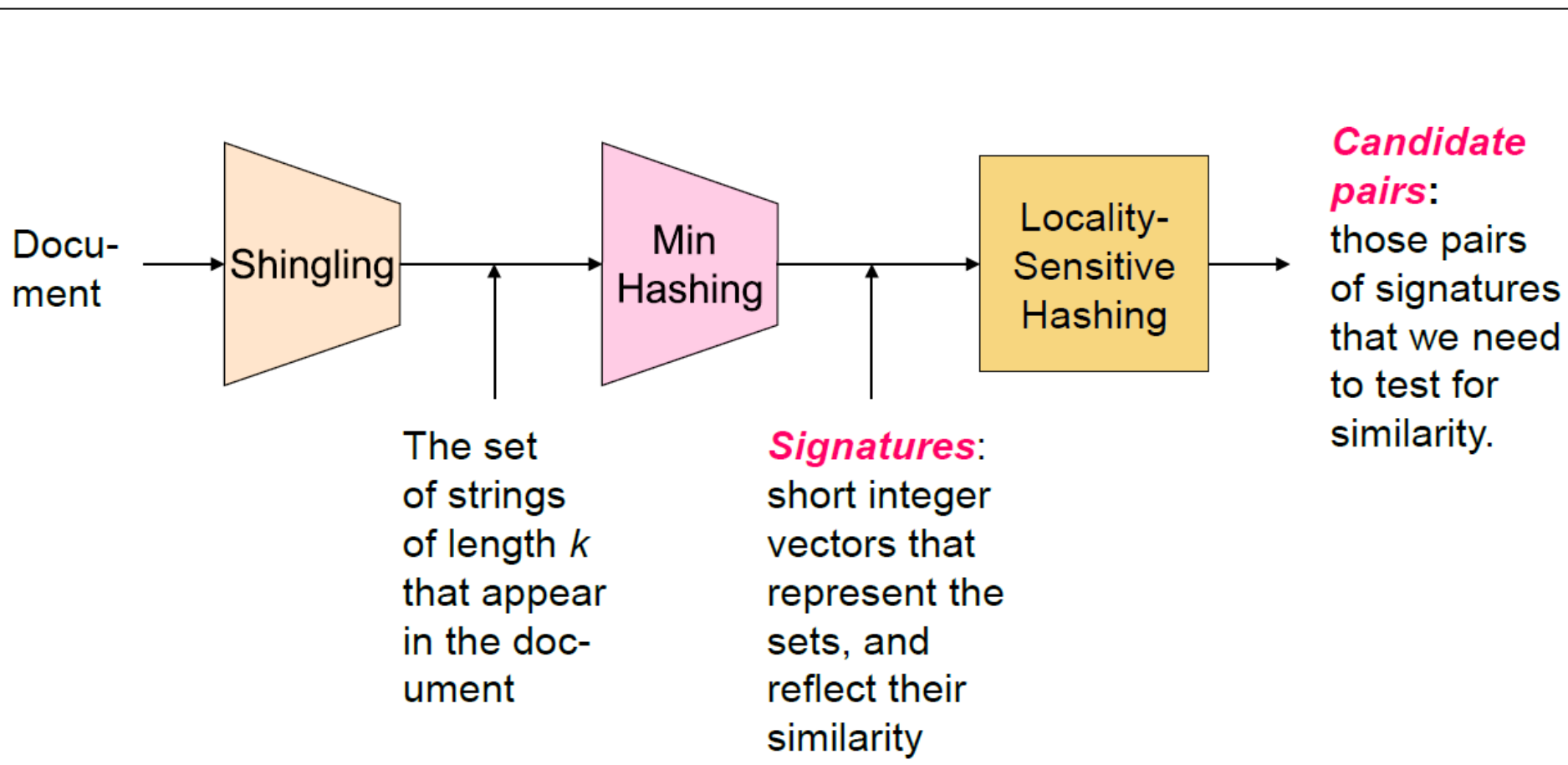


The Step-by-Step Guideline for Finding Similar Documents ⁷

- **Shingling**: Convert documents, emails, etc., to sets of short strings that appear within it
- **Minhashing**: Convert large sets to short signatures, while preserving similarity
- **Locality-sensitive hashing**: Focus on pairs of signatures likely to be from similar documents



The Step-by-Step Guideline for Finding Similar Documents 8



Source: Slide Adapted Jure Leskovic, Stanford CS246, Lecture Notes, see <http://cs246.stanford.edu>

Shingling – What is Shingling?

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- A **k -shingle** (or **k -gram**) for a document is a sequence of k tokens that appears within the document
 - Tokens can be characters, words or something else, depending on application
 - Assume tokens = characters for reading the book examples
- **Example:** $k=2$; $D_1 = \text{abcbab}$
 - Set of 2-shingles: $S(D_1) = \{\text{ab}, \text{bc}, \text{ca}\}$
 - **Option:** Shingles as a bag
- **Represent a doc by the set of hash values of its k -shingles**

How do we pick k ?



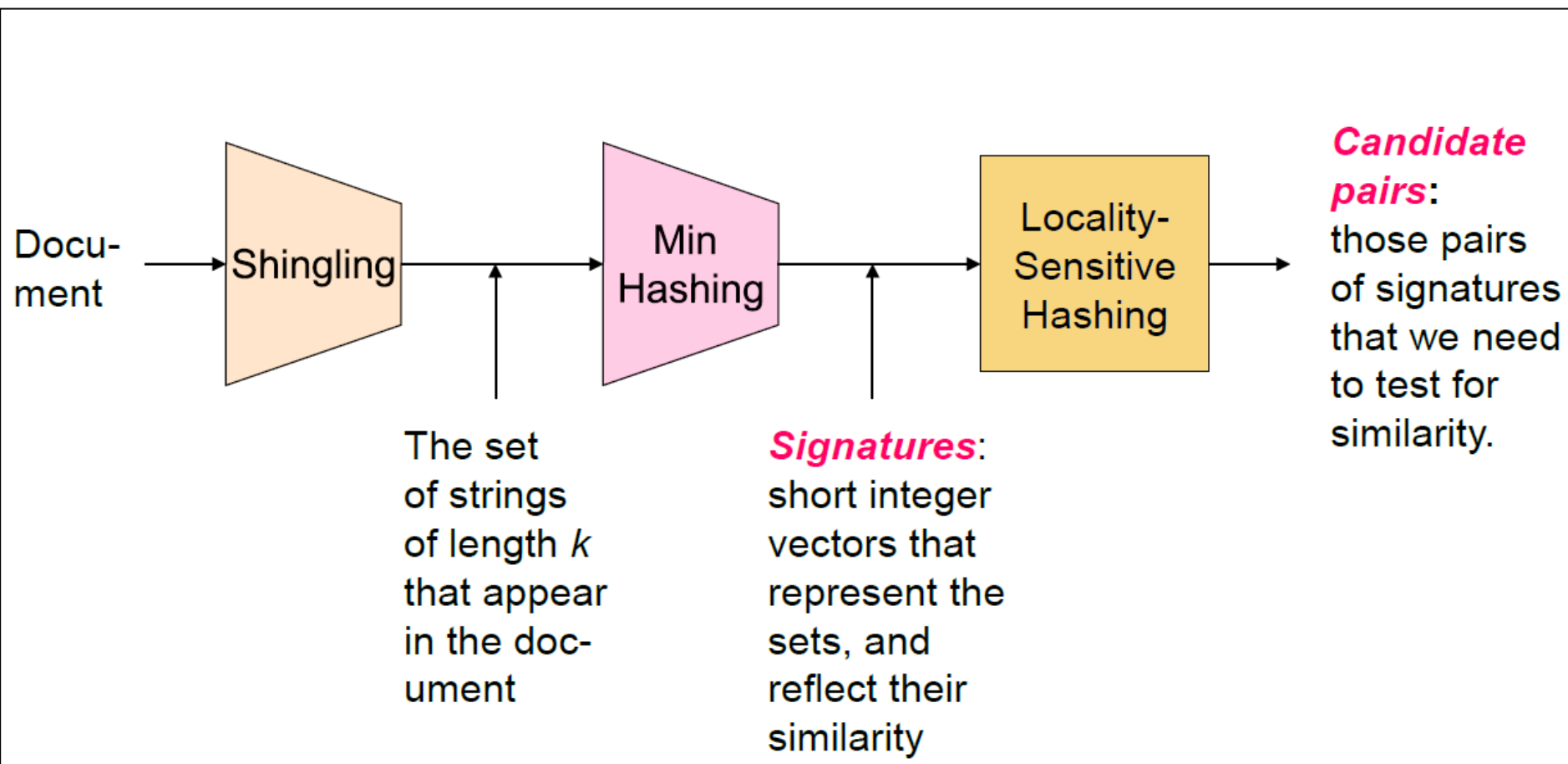
- Document D_1 = set of k -shingles $C_1 = S(D_1)$
- Equivalently, each document is a 0/1 vector in the space of k -shingles
 - Each unique shingle is a dimension
 - Vectors are very sparse
- A natural similarity measure is the **Jaccard similarity**:
 - $Sim(D_1, D_2) = |C_1 \cap C_2| / |C_1 \cup C_2|$
- **Assumption**: Documents that have lots of shingles in common have similar text, even if the text appears in different order



- Suppose we need to find near-duplicate documents among $N=1$ million documents
- Naïvely, we'd have to compute pairwise Jaccard similarities for every pair of docs



The Step-by-Step Guideline for Finding Similar Documents 12



Source: Slide Adapted Jure Leskovic, Stanford CS246, Lecture Notes, see <http://cs246.stanford.edu>

- Properties of the Matrix:
 - Rows = elements of the universal set
 - Columns = sets
- 1 if and only if the token is a member of the set
- Column similarity is the Jaccard similarity of the sets of their rows with 1
- Typical matrix is sparse



- Suppose that we have the following sets:
 - Universal set $\{a, b, c, d, e\}$.
 - $S1 = \{a, d\}$, $S2 = \{c\}$, $S3 = \{b, d, e\}$, and $S4 = \{a, c, d\}$.
- What is the *characteristic matrix* for this problem? 😊



- **So far:**

- Documents → Sets of shingles
- Represent sets as boolean vectors in a matrix

- **Next Goal:** Find similar columns

- **Approach:**

1. **Signatures of columns:** small summaries of columns
2. **Examine pairs of signatures** to find similar columns –**Essential property:** Similarities of signatures & columns are related
3. **Optional:** check that columns with similar sigs. are really similar

- **Warnings:**

- Comparing all pairs may take too long: **job for LSH/Minhash**



- Imagine the rows of the boolean matrix permuted under **random permutation** π
- Define a “hash” function $h_{\pi}(C)$ = the number of the first (in the permuted order π) row in which column C has value 1:

$$h_{\pi}(C) = \min \pi(C)$$

- Use several (e.g., 100) independent hash functions to create a signature of a column



Minhashing – An Example

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Suppose we pick the order of rows *beadc*. This permutation defines a minhash fn h that maps sets to rows. Compute the minhash value for all S according to h (i.e. for S_1 , S_2 , S_3 , and S_4).

<i>Element</i>	S_1	S_2	S_3	S_4
<i>a</i>	1	0	0	1
<i>b</i>	0	0	1	0
<i>c</i>	0	1	0	1
<i>d</i>	1	0	1	1
<i>e</i>	0	0	1	0

Note that: It is typical to replace the letters naming the rows by integers 0, 1, 2, etc.



Minhashing and Jaccard Similarity – A Surprising Property 18

- There is a remarkable connection between minhashing and Jaccard similarity of the sets that are minhashed.
 - The probability that the minhash function for a random permutation of rows produces the same value for two sets equals the Jaccard similarity of those sets.
- To see why, check p. 80 in the book ☺

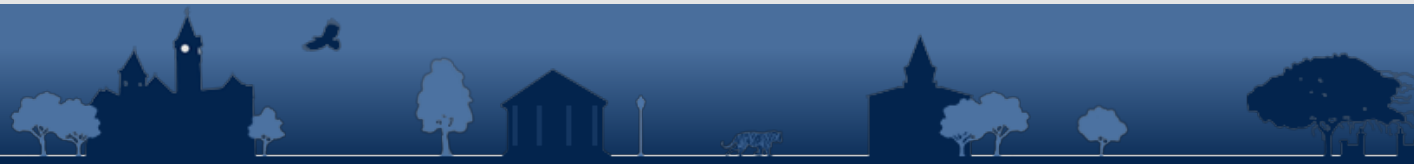


Minhash Signatures – An Example

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Using the new indices, let us return the signature matrix using these two hash functions ☺

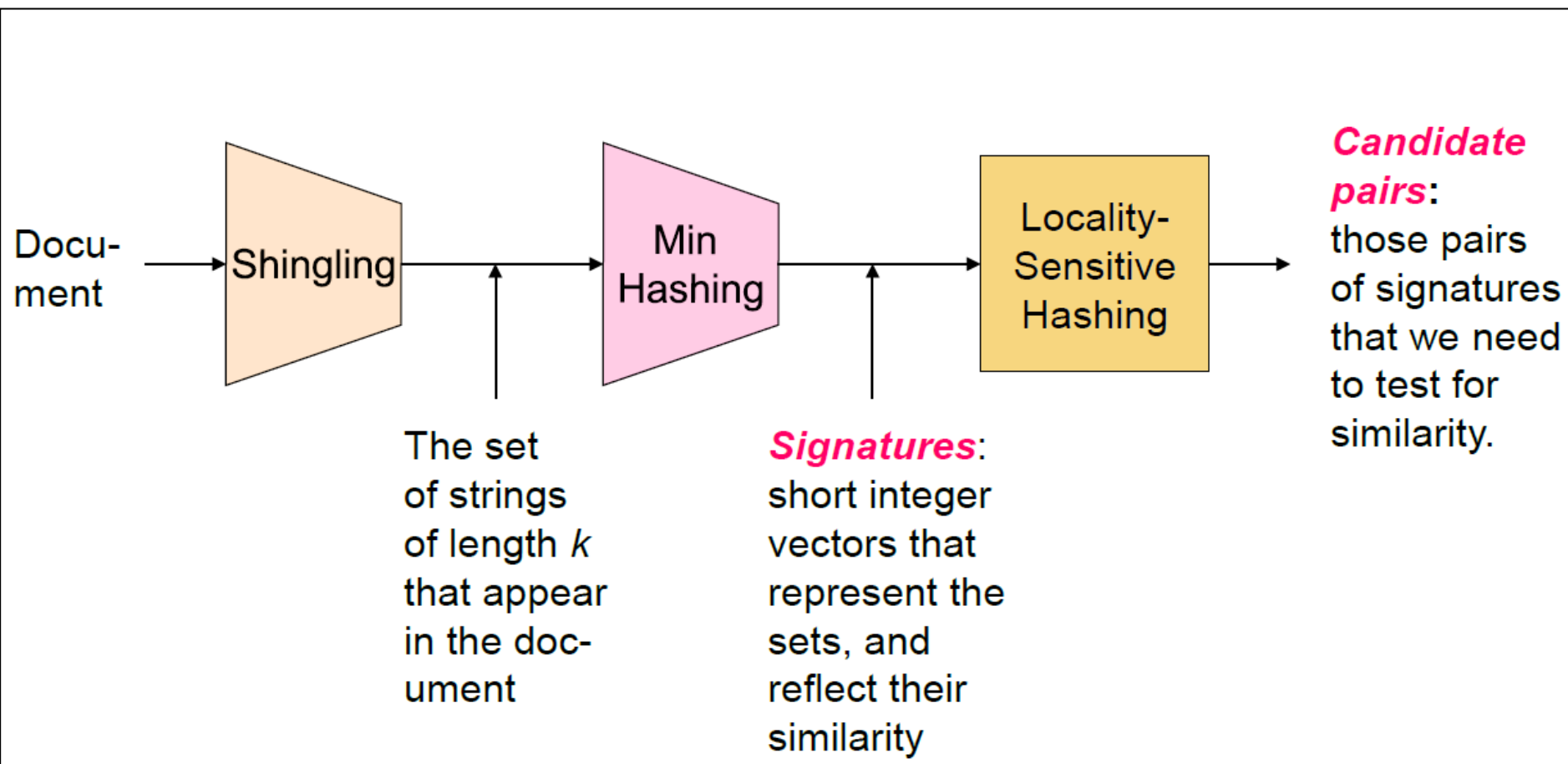
<i>Row</i>	S_1	S_2	S_3	S_4	$x + 1 \bmod 5$	$3x + 1 \bmod 5$
0	1	0	0	1	1	1
1	0	0	1	0	2	4
2	0	1	0	1	3	2
3	1	0	1	1	4	0
4	0	0	1	0	0	3



- We know: $\Pr[h_{\pi}(C1) = h_{\pi}(C2)] = \text{sim}(C1, C2)$
- Now generalize to multiple hash functions
- The similarity of two signatures is the fraction of the hash functions in which they agree
- Note: Because of the minhash property, the similarity of columns is the same as the expected similarity of their signatures



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- **Goal:** Find documents with Jaccard similarity at least s (for some similarity threshold, e.g., $s=0.8$)
- **LSH – General idea:** Use a function $f(x,y)$ that tells whether x and y is a **candidate pair**: a pair of elements whose similarity must be evaluated
- **For minhash matrices:**
 - Hash columns of signature matrix M to many buckets
 - Each pair of documents that hashes into the same bucket is a candidate pair



- Pick a similarity threshold s , a fraction < 1
- Columns x and y of M are a **candidate pair** if their signatures agree on at least fraction s of their rows:
 $M(i, x) = M(i, y)$ for at least frac. s values of i
- Note: We expect documents x and y to have the same similarity as their signatures (see previous slides)



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