### Project I

#### Overview, Guidelines, and JavaCC Tutorial

#### Project I

- **Due**: Monday Feb 18 (2 weeks+1 weekend)
- **Groups**: 2-3 People
- Languages: Java, JavaCC (like yacc)
- Goal: Familiarize yourself with manipulating RA
- Project I update posted on of Jan 31
  - Revisions to group submission instructions.

### Project I: Grading

- Test cases are provided
  - •edu.buffalo.cse.sql.test.\*
- Tests for both parts are provided
  - Part I:60 pts (55 from provided tests)
  - Part 2:40 pts (35 from provided tests)

### Part I

A Relational Algebra Interpreter

- **Given**: Data File(s), Relational Algebra Tree
- **Compute**: Evaluate the RA expression
- How? Use Naive In-Memory Algorithms
  - Selection/Projection/Union: Basic Pipelined Impl.
  - File Scan: Java IO (or NIO for better perf.)
  - Join: Nested-Loop (or others for better perf.)
  - Aggregates: Hash-table grouping



#### Part I

#### Keep in mind that in Project 3, **performance will be important** So starting now can't hurt.

#### Part I

- I. Implement relational operators.
  - The iterator interface is the simplest way.
  - init(), readNext(), hasMore(), close()
- 2. Implement a translator to assemble your relational operators from the RA plan.
- 3. Implement an evaluator for arithmetic expressions.

### Projection Example

- ProjectionOperator(ColNames, Schema, Source)
  this.OutCols = ColNames.map(Schema.getIndex);
  this.Source = Source;
- hasNext()

return this.Source.hasNext();

```
• readNext()
```

Datum[] next = this.Source.readNext();

return this.OutCols.map(next.get);

 $\frac{\text{Bindings}}{A \rightarrow 5}$  $B \rightarrow 2$  $C \rightarrow \text{`Foo'}$ 



eval(node, bindings) switch(node.type) / \* / • case return eval(node.l) \* eval(node.r) **'+':** ... case case const: return node.value case var: return bindings.get( node.var) 9

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### Part 2: Parsing SQL

- Design a **Grammar** for your language.
- Use a compiler compiler (e.g., Yacc, Bison, Javacc)
  - Takes a grammar as input
  - Generates "trees" for expressions in your language
- You specify how to interpret nodes in this tree.

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How do you specify a language?

$$A := BA | B | C$$

An **A** is a **B** followed by an **A**, or just a **B**, or just a **C** 

12

Friday, February 1, 13

B, C, BC, and BBBC are valid As

BCC is not valid, as although we can get 'C' as an A, no rule exists to let us prepend another 'C' before it.

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Α

B, C, BC, and BBBC are valid As

BCC is not valid, as although we can get 'C' as an A, no rule exists to let us prepend another 'C' before it.

CB is not valid, since we can only prepend 'B's to an 'A'

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How do you specify a language?

An **A** is a **B** followed by an **A**, or just a **B**, or just a **C** 

A := BA | B |

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Friday, February 1, 13

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BCC is not valid, as although we can get 'C' as an A, no rule exists to let us prepend another 'C' before it.

How do you specify a language?

A := BA | B | CB := 'B' C := 'C'

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B, C, BC, and BBBC are valid As

BCC is not valid, as although we can get 'C' as an A, no rule exists to let us prepend another 'C' before it.

How do you specify a language?

A := 
$$BA | B | C$$
  
B := 'B'  
C := 'C'

#### Which of the following are valid As?

'B' 'C' 'BC' 'BCC' 'BBBC' 'CB'

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B, C, BC, and BBBC are valid As

BCC is not valid, as although we can get 'C' as an A, no rule exists to let us prepend another 'C' before it.

#### Group Exercise

Define a grammar for basic arithmetic using addition, multiplication, and negation?

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**Hint:** This problem is underspecified. What additional information do you need?

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Use integers, and assume <INT> is defined

#### Tokens

- Strings are complex!
- Tokenize strings first
  - Simple, regular expression matching on strings to make 'tokens'.
    - e.g., <INT> := /-?[0-9]+/
    - An optional '-', followed by  $\geq 1$  digits
- Build parse trees with tokens as leaves.

JavaCC

- Java Compiler Compiler (Part of Sun JDK)
  - Reads in a grammar file.
  - Generates a Java class that parses your language.
  - The Java class can be compiled as normal.
- JavaCC files have a .jj suffix.

### A JavaCC .jj File

#### <u>Header</u>

Define code to explicitly copy into the java file

#### <u>Grammar Rules</u>

#### Declare rules each (roughly) of the form: A := B | C | D | ...

#### Tokenizer Rules

Declare all token types and how to identify them

### A JavaCC Header





```
A JavaCC Grammar Rule
                Rule Name
  Program Program()
    { Program p = new Program();
      Statement s;}
    {
       s = Statement() <EOS>
            { p.addStatement(s); }
      )+ <EOF>
        { return p; }
    }
```

```
Program Program() :
    { Program p = new Program();
    Statement s;}
    {
        ( s = Statement() <EOS>
            { p.addStatement(s); }
        Rule )+ <EOF>
Body { return p; }
    }
}
```

```
A JavaCC Grammar Rule
       Rule Type
  Program Program() :
    { Program p = new Program();
      Statement s;}
    {
       s = Statement() <EOS>
            { p.addStatement(s); }
      )+ <EOF>
        { return p; }
    }
```

```
Program Program() :
    { Program p = new Program();
Rule Statement s;}
Header {
    ( s = Statement() <EOS>
        { p.addStatement(s); }
    )+ <EOF>
        { return p; }
    }
}
```

```
Program Program(+) :--- Rule Arguments
{ Program p = new Program();
   Statement s;}
{
   ( s = Statement() <EOS>
        { p.addStatement(s); }
   )+ <EOF>
        { return p; }
}
```

#### Each JavaCC rule is a function definition The Rule Body is a **Pattern**

A JavaCC Grammar Rule Rule Type Rule Name Program Program (+) ---- Rule Arguments Program p = new Program(); Rule Statement s;} Header { ( s = Statement() <EOS> { p.addStatement(s); } Rule )+ <EOF> Body { return p; } Each JavaCC rule is a function definition The Rule Body is a **Pattern** 

#### Pattern :=

[var =] <TOKEN>

#### Meaning:

Note:

### Match a token of the indicated type optionally assign it to variable var

var must be of the Java class Token
var.image is the string matched by the token's regex

#### EXAMPLE

- IntBase :=
  - $i = \langle INT \rangle$

Meaning:

match an <INT> token assign it to the variable i

Note:

You'll soon see how to use i to assign a 'meaning' to this token (i.e, an integer)

#### Pattern :=

[var =] Rule(args)

# Invoke the indicated rule (with arguments)Meaning:The pattern matches if the rule finds a match<br/>Optionally assign the matched value to varNote:var must be of the type 'returned' by Rule()

Pattern :=

Pattern1 Pattern2 Pattern3 ...

Meaning:

Match an input consisting of something matching pattern 1, followed by a match for pattern 2, followed by pattern 3, ...

#### EXAMPLE

Program :=

s = Statement() <EOS>

Meaning:

Find a match for the 'Statement' rule followed by an EOS token. Assign the matched Statement to s

#### Pattern :=

Pattern { java code }

#### Meaning:

If the Pattern is matched Then Execute the Java Code

#### EXAMPLE

IntBase :=

i = <INT> {return parseInt(i.image);}

Meaning:

match an <INT> token and assign it to the variable i Return the parsed integer as the value for this rule.

**Note:** Remember, i.image is the string matched by the token

Pattern :=

Pattern[op] where [op] is ?, + or \*

Optional/Repeating pattern. Match... ? : at most once Heaning: + : at least once (greedy) \* : any number of times (greedy) Code literals and variable assignments in the pattern

are executed once for each match

#### EXAMPLE

Program :=

(Statement() <EOS>)+ <EOF>

Match one or more Statements as long as eachMeaning:statement is followed by an EOS TokenAfter all statements, there should be an EOF

**Note:** Parenthesis may be used to group patterns together

#### Pattern :=

Pattern1 | Pattern2

# Meaning:Match either of the two patternsNote:The generated parser looks only at the next N tokens<br/>before committing to this decision (default N=1)<br/>N is known as the lookahead parameter



Friday, February 1, 13







### A JavaCC Tokenizer Rule

SKIP: { regex | regex | ... }

Don't generate a token for expressions that match a regex

Define the regex to generate each token type

#### A JavaCC Tokenizer Rule

Ignore whitespace (unless it's the 2nd+ character in a token)



### A JavaCC Tokenizer Rule





### Invoking JavaCC Parsers

Instantiate your parser with a java.io.InputStream

```
public static List<List<Datum[]>> execFile(File fname)
    throws SqlException, IOException, ParseException
{
    Sql parser = new Sql(new FileInputStream(fname));
    Program p = parser.Program();
    return Runtime.execFile(p);
}
Each grammar rule defines a class method
```

Call the class method for the 'root' of your grammar.

### Testing

- SQL examples in sql/test
- Use edu.buffalo.cse.sql.test.\* to test both RA and SQL execution
  - Successful tests = guaranteed  $\geq$  90% grade
- Make sure to create your own test cases!

