

CS 5363, Fall 2013

Attribute Grammars

1 Reading / Source Material

- Cooper: Chapter 4
- Scott: Chapter 4

2 Objectives: Be able to ...

1. Given a simple attributed grammar and a sentence in the language described by the grammar, calculate the attributes.
2. Given an attributed grammar, identify attributes of non-terminals for which the attribute is always inherited
3. Given an attributed grammar, identify attributes of non-terminals for which the attribute is always synthesized
4. Given an expression grammar in which the productions for left-associative operations have undergone recursion elimination to produce an LL(1) grammar, be able to write an L-attribute grammar that evaluates the expression applying left-associativity.
5. write an L-attributed grammar that generates an AST for a typical procedural programming language described by a BNF and informal semantics
6. given an L-attributed grammar, write a recursive descent parser with ad-hoc rules implementing the attributed grammar
7. write a tree grammar to type-check the AST of programs written in a typical procedural programming language described by a BNF, informal semantics, and informal type rules

3 Outline

1. Semantic Elaboration ("Context-Sensitive" Analysis)
 - (a) beyond syntax, beyond context-free grammars
 - (b) examples:
 - i. declaration before use
 - ii. number of parameters in procedural call
 - iii. types of variables / expressions
 - iv. return statements
 - (c) purpose:
 - i. discover information needed to generate efficient code

- ii. detect errors
- 2. Attribute Grammar
 - (a) context free grammar + attributes + semantic functions
 - (b) expression example
 - (c) synthesized attributes
 - i. rules defining the attribute, define it for the non-terminal on the left-hand side (lhs) of the production in which the rule occurs
 - ii. that is, the attribute of a non-terminal is defined in terms of the non-terminal's own attributes and those of its children
 - (d) inherited attributes
 - i. rules defining the attribute, define it for a non-terminal on the right-hand side (rhs) of the production in which the rule occurs
 - ii. that is, the attribute of a non-terminal is defined in terms of the non-terminal's own attributes, those of its parents, and those of its siblings
 - (e) S-attributed grammars
 - i. an attribute grammar where all attributes are synthesized
 - (f) L-attributed grammars
 - i. an attribute grammar where all attributes are inherited or synthesized and satisfy the following additional constraints:
 - A. the synthesized attributes of a lhs non-terminal depend only on inherited attributes of that lhs non-terminal and on (synthesized or inherited) attributes of the rhs non-terminals.
 - B. the inherited attributes of a rhs non-terminal depend only on inherited attributes of the lhs non-terminal and on (synthesized or inherited) attributes of rhs non-terminals that occur to the left of the non-terminal for which the attribute is being defined.
 - ii. relation to recursive descent parser
 - A. inherited attributes can be passed as parameters
 - B. synthesized attributes can be returned as values
- 3. tree grammars

4 Vocabulary

Attribute Grammar, attribute, synthesized attribute, inherited attribute, L-attributed grammar, S-attributed grammar

5 Questions

1. Consider the following grammar describe the binary numbers that are multiples of four.

```

<Number> -> 0 | <HighestBit> <List> <DoubleZero>
<HighestBit> -> 1
<DoubleZ> -> 00
<List> -> 1 <List> | 0 <List> |

```

For example, 1000 is a binary number satisfying the above grammar. Now, suppose we change it into an attribute grammar with the following attributes, such that 1000's value is 8 and its HighestBit's position is 3, that is, the position of the rightmost 0 is 0.

| Symbol | Attribute |
|--------------|-----------------|
| <Number> | value |
| <List> | value, position |
| <HighestBit> | position |

- (a) Fill in the Attribution Rules for each production.

| Production | Attribution Rules |
|---|---|
| 1 <Number> -> 0 | Number.value := |
| 2 <Number> -> <HighestBit> <List> <DoubleZ> | HighestBit.position := Number.value := |
| 3 <List0> -> 1 <List1> | List0.position := List0.value := |
| 4 <List0> -> 0 <List1> | List0.position := List0.value := |
| 5 <List> -> | List.value := List.position := |

- (b) Using the above attribute grammar, build the syntax tree for the binary number 101100, annotating all attributes with the corresponding value.
- (c) Indicate which attributes are inherited and which are synthesized.
- (d) Is your attributed grammar L-attributed?

2. Given the attribute grammar:

```

<A> ::= <B>
^ A.x = B.x
^ B.y = 0

<B1> ::= c <B2>

```

```

^ B1.x = B2.x + 1
^ B2.y = B1.y + 1

```

```

<B> ::=
^ B.x = 0

```

- (a) Which attributes are synthesized and which are inherited?
- (b) Draw an attribute-annotated parse tree for the sentence: ccc
Add arrows showing how attribute information flows between pairs of attributes.

3. Consider the attribute grammar:

```

<A> ::= x <B>
^ A.result = B.result
<B1> ::= y <B2>
^ B1.result = B2.result + 2
<B1> ::= z <B2>
^ B1.result = B2.result + 1
<B> ::=

```

- (a) Which attributes are synthesized and which are inherited?
- (b) Draw an attribute-annotated parse tree for the sentence: xzyy
Add arrows showing how attribute information flows between pairs of attributes.

4. Consider the following attributed tree grammar for evaluating an expression with an integer value of X whose value is 5 and Y whose value is 3.

```

start: <start> ::= <expr>
^ <expr>.env = <'X', 5, <'Y', 3, nil>>

id: <expr> ::=
^ <expr>.val = lookup(id.$literalText$)

int_const: [expr] ::=
^ <expr>.val = text2Integer(int_const.$literalText$)

'+': <expr1> ::= <expr2> <expr3>
^ <expr2>.symtab = <expr1>.symtab
^ <expr3>.symtab = <expr1>.symtab
^ <expr1>.val = <expr2>.val + <expr3>.val

'*': <expr1> ::= <expr2> <expr3>
^ <expr2>.symtab = <expr1>.symtab
^ <expr3>.symtab = <expr1>.symtab
^ <expr1>.val = <expr2>.val * <expr3>.val

```

Assume that int_const and id nodes have a pre-defined attribute \$literalText\$ that corresponds to the text of the id or integer constant, respectively.

Also assume an auxiliary function lookup, defined as follows:

```

lookup(X, env) = if env = nil
                  then error
                  else let <Y, n env2> = env in
                       if Y = X then n else lookup(X, env2)
end

```