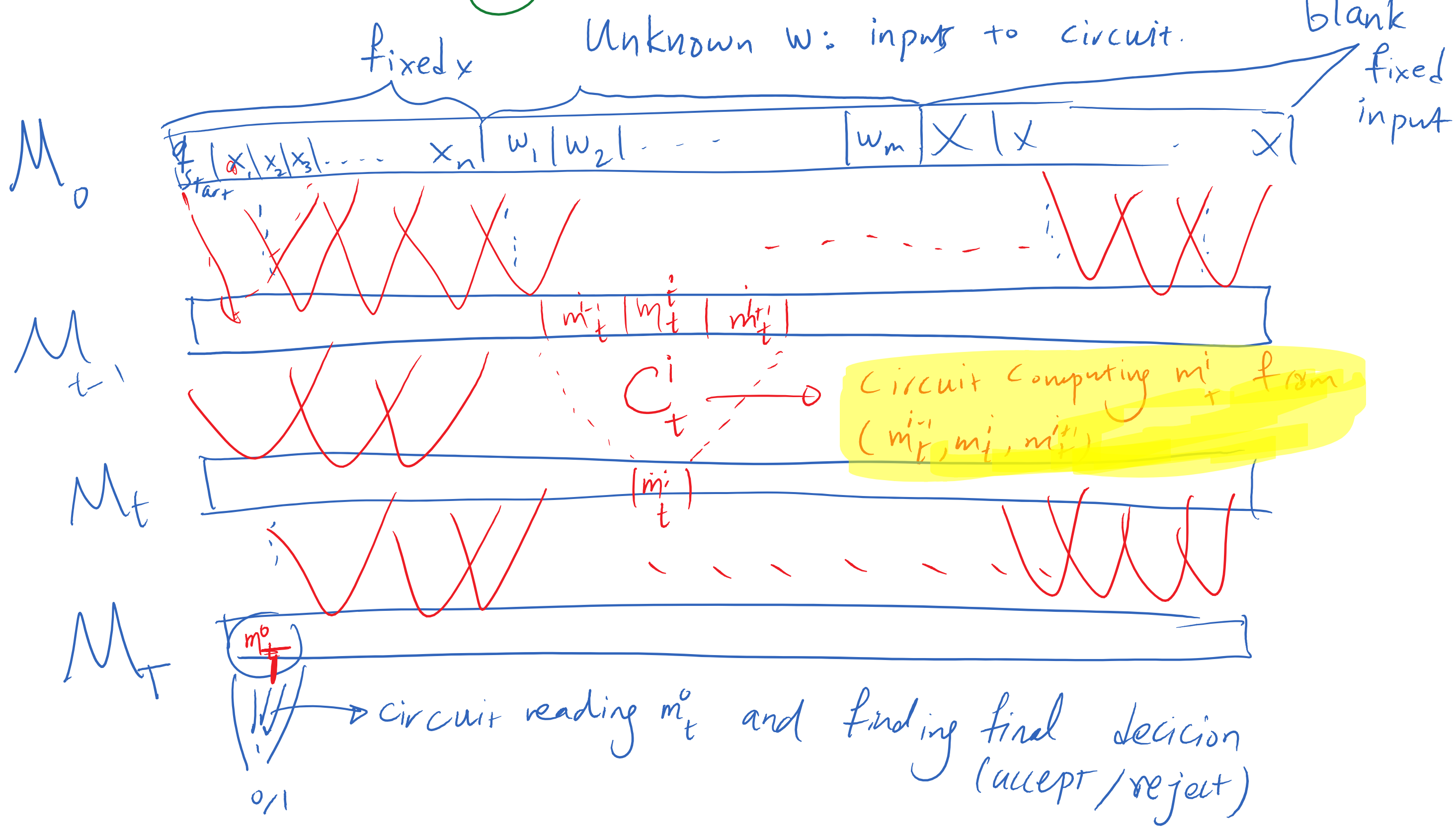
Cook-Levin Theorem in Picture:

Let M be Turing machine checking witnesses for $L \in NP$.

The reduction from L to CircSAT will take x and maps it (efficiently) to a circuit C such that $x \in L$ iff C is satisfiable (i.e $C(\alpha)=1$ for some α).

The circuit C is explained in the note on Piazza has the following form:



Note that lots of Computation done by C is not needed: for example the first layer can simply ignore what is after "w" because they are all blank. Also in the last layer we ~~only~~ care about $\underline{m'}_T$ not the rest. However, what is important for us is that C has size $T^2 \leq poly(|x|)$ and can be computed in $poly(|x|)$ time.