

# CS635: Tools and Environments for Optimization. Lecture 10

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## New Problem!

- Given network  $G = (N, A)$
- Each node  $i \in N$  has a “supply”  $b_i$
- $b_i < 0 \Rightarrow$  node  $i$  “demands” an amount  $b_i$ .
- For feasibility, we must have that  $\sum_{i \in N} b_i = 0$ . (Supply = Demand)
- Arcs  $a \in A$  may have costs  $c_a$  or capacities  $u_a$

### Min-Cost Network Flow

Find a minimum cost flow of the commodity from supply nodes to demand nodes without exceeded the arc capacity



## MCNF: Mathematical Formulation

$$\min \sum_{(i,j) \in A} c_{ij} x_{ij}$$

$$\sum_{j \mid (i,j) \in A} x_{ij} - \sum_{j \mid (i,j) \in A} x_{ji} = b_i \quad \forall i \in N$$

$$0 \leq x_{ij} \leq u_{ij} \quad \forall (i,j) \in A$$

GAMS time: `mincost.gms`, `mincost2.gms`

- `abort`
- **Important:** Note use of dynamic set `arc`



## mincost2.gms

- You can loop over sets in GAMS. `loop`
- **Important GAMS attributes:** `modelstat`, `solvestat`



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## Attributes Controlled by the User

- `iterlim` iteration limit
- `limcol` number of columns displayed for each block of variables
- `limrow` number of rows displayed for each block of equations
- `optfile` option file usage
- `reslim` time limit for solver. Usually in CPU seconds
- `solprint` solution print option



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## Attributes Controlled by the Solver

- `iterusd`: number of iterations used
- `modelstat`: model status (cf. Section 10.5.4)
- `numequ`: number of single equations generated
- `numvar`: number of single variables generated
- `resusd`: resource units (in CPU seconds) used to solve model
- `solvestat`: solver status (cf. Section 10.5.4)



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## Model Status (10.5.4)

- 1 OPTIMAL
- 2 LOCALLY OPTIMAL
- 3 UNBOUNDED
- 4 INFEASIBLE
- 5 LOCALLY INFEASIBLE
- 6 INTERMEDIATE INFEASIBLE
- 7 INTERMEDIATE NONOPTIMAL
- 8 INTEGER SOLUTION
- 9 INTERMEDIATE NON-INTEGER
- 10 INTEGER INFEASIBLE



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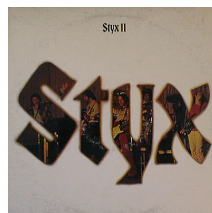
## Solver Status (10.5.5)

- 1 Normal Completion
- 2 Iteration Interrupt (`option iterlim`)
- 3 Resource Interrupt (`option redlim`)
- 4 Terminated By Solver
- 5 Evaluation Error Limit
- 6 Unknown



## Sailco, Cont.

## Come Sail Away, Come Sail Away



- Sailco manufactures sailboats. During the next 4 months the company must meet the following demands for their sailboats:
  - 40, 60, 75, 25
- At the beginning of Month 1, Sailco has 10 boats in inventory.
- Each month it must determine how many boats to produce.

- During any month, Sailco can produce up to 40 boats with regular labor and an unlimited number of boats with overtime labor.
  - Boats produced with regular labor cost \$400 each to produce,
  - while boats produced with overtime labor cost \$450 each.
- It costs \$20 to hold a boat in inventory from one month to the next. Find the production and inventory schedule that minimizes the cost of meeting the next 4 months' demands on time.



## Sailco

- Applications that are not obviously network problems can be modeled as such.

## Model as a network

- One node per month – given demand at each month.
- Inflow to each month node from a "regular labor" node, with capacity 40 and cost 400
- Inflow to each month node from a "overtime labor" node, with unlimited capacity and cost 450
- Month-to-month arcs representing inventory, (cost 20)
- Balance constraints applied only at month nodes.

- GAMS Time: [sailco2.gms](#)



## Max Flow

## Maximum Flow Problem

Given a network capacitated  $G = (N, A)$ , with capacities  $u \in \mathbb{R}_+^{|A|}$ , a **source node**  $s \in N$ , and a **sink node**  $t \in V$ , what is the **maximum flow** that can be sent from  $s$  to  $t$ .

- Model "what goes in = what come out".

$$\sum_{j:(i,j) \in A} x_{ij} - \sum_{j:(j,i) \in A} x_{ji} = 0$$

- This constraint holds for all nodes except the source and sink:  $(\forall i \in N \setminus \{s\} \setminus \{t\})$

## Max Flow Problem

$$\begin{aligned} \max \quad & \sum_{j:(s,j) \in A} x_{sj} \\ \text{s.t.} \quad & \sum_{j:(i,j) \in A} x_{ij} - \sum_{j:(j,i) \in A} x_{ji} = 0 \quad \forall i \in N \setminus \{s\} \setminus \{t\} \\ & x_{ij} \leq u_{ij} \quad \forall (i,j) \in A \\ & x_{ij} \geq 0 \quad \forall (i,j) \in A \end{aligned}$$

- Note that due to flow balance, objective is the same as

$$\sum_{j:(j,t) \in A} x_{jt}$$



## Let's Have a Picnic!

- The Hatfields, Montagues, McCoys and Capulets are going on their annual family picnic.
- Four cars are available to transport the families to the picnic.
- The cars can carry the following numbers of people: car 1, 4; car 2, 3; car 3, 3; car 4, 4.
- There are four people in each family, and no car can carry more than two people from any one family.
- Determine the maximum number of people that can be transported to the picnic.

## Max Flow Problem

- Let's try and model this is a max flow problem

### The \$0.0001 Question

- Who Can Make This A Max Flow Problem?

### GAMS Break



## SPP

### Shortest Path Problem

Find the shortest path through a network from a specified origin to a specified destination.

- This is a special case of min cost flow, where costs are distances, there is a single supply source with supply +1 and a single demand (sink)
- There are specialized algorithms for shortest path that do not exploit the relationship with min-cost flow and in fact are more efficient from the complexity viewpoint:  $O(|A| + |N| \log |N|)$



## Examples

- `short1.gms`
  - Note: `alias` statement
  - Note: `sqr` and `sqrt`
- `short2.gms`
  - Note: `option seed=`
  - Note: `loop` (This is `cool`)
  - Note: difference between assignments to `arcs(i,j)` and `arcs(i,i)`. Use of the `alias` is essential!
  - Note: `modelstat`

### Count the CPU Time:

- `totalSolveTime = totalSolveTime + short.resusd;`



## More Goodness

### Homework #4

- Due Friday @11AM.
- Three problems.

