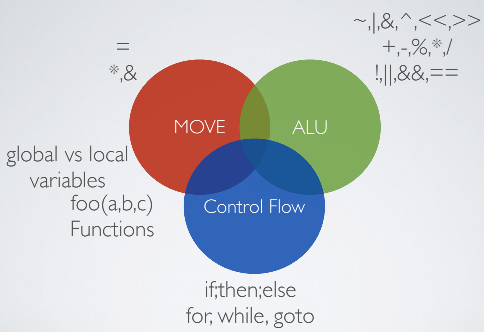
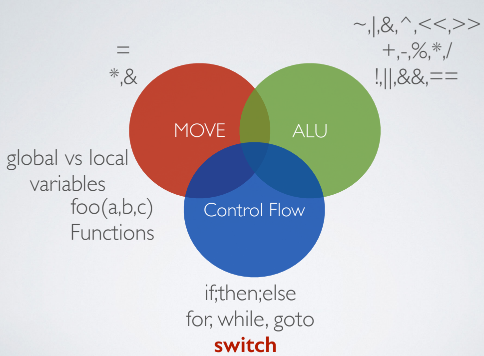
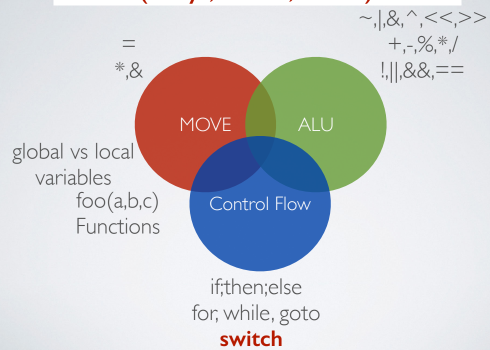
MACHINE LEVEL REPRESENTATION 3

CAS CS 210 Computer Systems
Based on CMU Slides
CS:APP2e





Programmer Defined Types
(arrays, structs, unions)



Basic Data Types

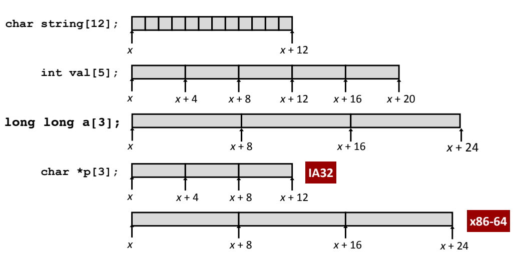
- **Integral**
  - Stored & operated on in general (integer) registers
  - Signed vs. unsigned depends on instructions used

Intel	GAS	Bytes	C
byte	**b**	1	**[unsigned] char**
word	**w**	2	**[unsigned] short**
double word	**l**	4	**[unsigned] int**
quad word	**q**	8	**[unsigned] long int** (x86-64)

Array Allocation

- **Basic Principle**

  T **A[***L***];**
  - Array of data type *T* and length *L*
  - Contiguously allocated region of *L* * **sizeof**(*T*) bytes



Array Access

- **Basic Principle**

  T **A**[*L*];
  - Array of data type *T* and length *L*
  - Identifier **A** can be used as a pointer to array element 0: Type *T**

`int val[5];`

	1	5	2	1	3
x	*x* + 4	*x* + 8	*x* + 12	*x* + 16	*x* + 20

- **Reference** **Type** **Value**

```
val[4]
val
val+1
&val[2]
val[5]
*(val+1)
val + i
```

Array Access

- **Basic Principle**

  T **A**[*L*];
  - Array of data type *T* and length *L*
  - Identifier **A** can be used as a pointer to array element 0: Type *T**

`int val[5];`

	1	5	2	1	3
x	*x* + 4	*x* + 8	*x* + 12	*x* + 16	*x* + 20

Reference	Type	Value
`val[4]`	`int`	3
`val`	`int *`	*x*
`val+1`	`int *`	*x* + 4
`&val[2]`	`int *`	*x* + 8
`val[5]`	`int`	??
`*(val+1)`	`int`	5
`val + i`	`int *`	*x* + 4 *i*

Array Example

```
typedef int zip_dig[5];

zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

`zip_dig cmu;`

	1	5	2	1	3
16	20	24	28	32	36

`zip_dig mit;`

	0	2	1	3	9
36	40	44	48	52	56

`zip_dig ucb;`

	9	4	7	2	0
56	60	64	68	72	76

- **Declaration "`zip_dig cmu`" equivalent to "`int cmu[5]`"**
- **Example arrays were allocated in successive 20 byte blocks**
  - Not guaranteed to happen in general

Array Accessing Example

`zip_dig cmu;`

1	5	2	1	3

16 20 24 28 32 36

```
int get_digit
  (zip_dig z, int dig)
{
  return z[dig];
}
```

IA32

```
  # %edx = z
  # %eax = dig
  movl (%edx,%eax,4),%eax  # z[dig]
```

- Register `%edx` contains starting address of array
- Register `%eax` contains array index
- Desired digit at `4*%eax + %edx`
- Use memory reference `(%edx,%eax,4)`

Referencing Examples

`zip_dig cmu;`

1	5	2	1	3

16 20 24 28 32 36

`zip_dig mit;`

0	2	1	3	9

36 40 44 48 52 56

`zip_dig ucb;`

9	4	7	2	0

56 60 64 68 72 76

- **Reference** **Address** **Value** **Guaranteed?**

`mit[3]`
`mit[5]`
`mit[-1]`
`cmu[15]`

Referencing Examples

`zip_dig cmu;`

1	5	2	1	3

16 20 24 28 32 36

`zip_dig mit;`

0	2	1	3	9

36 40 44 48 52 56

`zip_dig ucb;`

9	4	7	2	0

56 60 64 68 72 76

Reference	Address	Value	Guaranteed?
`mit[3]`	`36 + 4* 3 = 48`	3	Yes
`mit[5]`	`36 + 4* 5 = 56`	9	No
`mit[-1]`	`36 + 4*-1 = 32`	3	No
`cmu[15]`	`16 + 4*15 = 76`	??	No

- No bound checking
- Out of range behavior implementation-dependent
- No guaranteed relative allocation of different arrays

Array Loop Example

- **Original**

```
int zd2int(zip_dig z)
{
  int i;
  int zi = 0;
  for (i = 0; i < 5; i++) {
    zi = 10 * zi + z[i];
  }
  return zi;
}
```

- **Transformed**
  - As generated by GCC
  - Eliminate loop variable `i`
  - Convert array code to pointer code
  - Express in do-while form (no test at entrance)

```
int zd2int(zip_dig z)
{
  int zi = 0;
  int *zend = z + 4;
  do {
    zi = 10 * zi + *z;
    z++;
  } while (z <= zend);
  return zi;
}
```

Array Loop Implementation (IA32)

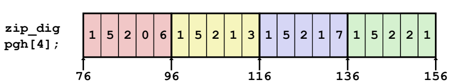
- **Registers**
  - `%ecx z`
  - `%eax zi`
  - `%ebx zend`
- **Computations**
  - `10*zi + *z` implemented as `*z + 2*(zi+4*zi)`
  - `z++` increments by 4

```
int zd2int(zip_dig z)
{
  int zi = 0;
  int *zend = z + 4;
  do {
    zi = 10 * zi + *z;
    z++;
  } while(z <= zend);
  return zi;
}
```

```
  # %ecx = z
  xorl %eax,%eax           # zi = 0
  leal 16(%ecx),%ebx       # zend  = z+4
.L59:
  leal (%eax,%eax,4),%edx  # 5*zi
  movl (%ecx),%eax         # *z
  addl $4,%ecx             # z++
  leal (%eax,%edx,2),%eax  # zi = *z + 2*(5*zi)
  cmpl %ebx,%ecx           # z : zend
  jle .L59                 # if <= goto loop
```

Nested Array Example

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
  {{1, 5, 2, 0, 6},
   {1, 5, 2, 1, 3 },
   {1, 5, 2, 1, 7 },
   {1, 5, 2, 2, 1 }};
```



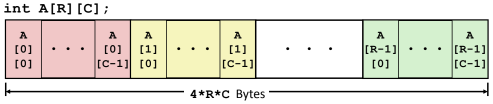
- "`zip_dig pgh[4]`" **equivalent to** "`int pgh[4][5]`"
  - Variable `pgh`: array of 4 elements, allocated contiguously
  - Each element is an array of 5 `int`'s, allocated contiguously
- **"Row-Major" ordering of all elements guaranteed**

Multidimensional (Nested) Arrays

- **Declaration**
  - *T* **A**[*R*][*C*];
  - 2D array of data type *T*
  - *R* rows, *C* columns
  - Type *T* element requires *K* bytes
- **Array Size**
  - *R* * *C* * *K* bytes
- **Arrangement**
  - Row-Major Ordering

$$\begin{bmatrix} \texttt{A[0][0]} & \cdots & \texttt{A[0][C-1]} \\ \vdots & & \vdots \\ \texttt{A[R-1][0]} & \cdots & \texttt{A[R-1][C-1]} \end{bmatrix}$$

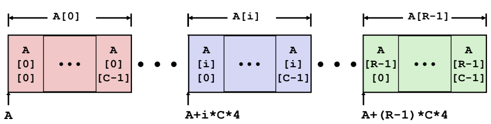
```
int A[R][C];
```



Nested Array Row Access

- **Row Vectors**
  - `A[i]` is array of *C* elements
  - Each element of type *T* requires *K* bytes
  - Starting address `A +` *i* * (*C* * *K*)

```
int A[R][C];
```



Nested Array Row Access Code

```
int *get_pgh_zip(int index)
{
  return pgh[index];
}
```

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
  {{1, 5, 2, 0, 6},
   {1, 5, 2, 1, 3 },
   {1, 5, 2, 1, 7 },
   {1, 5, 2, 2, 1 }};
```

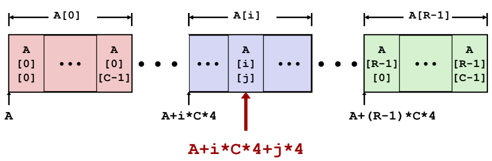
```
  # %eax = index
  leal (%eax,%eax,4),%eax  # 5 * index
  leal pgh(,%eax,4),%eax   # pgh + (20 * index)
```

- **Row Vector**
  - `pgh[index]` is array of 5 `int`'s
  - Starting address `pgh+20*index`
- **IA32 Code**
  - Computes and returns address
  - Compute as `pgh + 4*(index+4*index)`

Nested Array Row Access

- **Array Elements**
  - `A[i][j]` is element of type *T*, which requires *K* bytes
  - Address `A + ` *i* * (*C* * *K*) + *j* * *K* = *A* + (*i* * *C* + *j*)* *K*

```
int A[R][C];
```



Nested Array Element Access Code

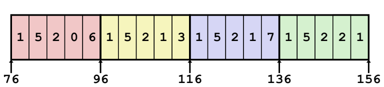
```
int get_pgh_digit
  (int index, int dig)
{
  return pgh[index][dig];
}
```

```
# %ecx = dig
# %eax = index
leal 0(,%ecx,4),%edx          # 4*dig
leal (%eax,%eax,4),%eax       # 5*index
movl pgh(%edx,%eax,4),%eax    # *(pgh + 4*dig + 20*index)
```

- **Array Elements**
  - `pgh[index][dig]` is `int`
  - Address: `pgh + 20*index + 4*dig`
- **IA32 Code**
  - Computes address `pgh + 4*dig + 4*(index+4*index)`
  - `movl` performs memory reference

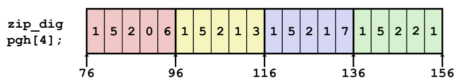
Strange Referencing Examples

```
zip_dig
pgh[4];
```



Reference	Address	Value	Guaranteed?
`pgh[3][3]`			
`pgh[2][5]`			
`pgh[2][-1]`			
`pgh[4][-1]`			
`pgh[0][19]`			
`pgh[0][-1]`			

Strange Referencing Examples



Reference	Address	Value	Guaranteed?
`pgh[3][3]`	`76+20*3+4*3 = 148`	2	Yes
`pgh[2][5]`	`76+20*2+4*5 = 136`	1	Yes
`pgh[2][-1]`	`76+20*2+4*-1 = 112`	3	Yes
`pgh[4][-1]`	`76+20*4+4*-1 = 152`	1	Yes
`pgh[0][19]`	`76+20*0+4*19 = 152`	1	Yes
`pgh[0][-1]`	`76+20*0+4*-1 = 72`	??	No

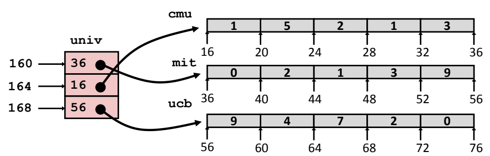
- Code does not do any bounds checking
- Ordering of elements within array guaranteed

Multi-Level Array Example

```
zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

```
#define UCOUNT 3
int *univ[UCOUNT] = {mit, cmu, ucb};
```

- **Variable `univ` denotes array of 3 elements**
- **Each element is a pointer**
  - 4 bytes
- **Each pointer points to array of `int`'s**



Element Access in Multi-Level Array

```
int get_univ_digit
  (int index, int dig)
{
  return univ[index][dig];
}
```

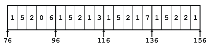
```
  # %ecx = index
  # %eax = dig
  leal 0(,%ecx,4),%edx      # 4*index
  movl univ(%edx),%edx      # Mem[univ+4*index]
  movl (%edx,%eax,4),%eax  # Mem[...+4*dig]
```

- **Computation (IA32)**
  - Element access `Mem[Mem[univ+4*index]+4*dig]`
  - Must do two memory reads
    - First get pointer to row array
    - Then access element within array

Array Element Accesses

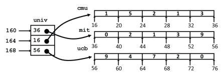
Nested array

```
int get_pgh_digit
  (int index, int dig)
{
  return pgh[index][dig];
}
```



Multi-level array

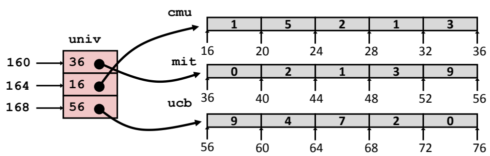
```
int get_univ_digit
  (int index, int dig)
{
  return univ[index][dig];
}
```



Access looks similar, but element:

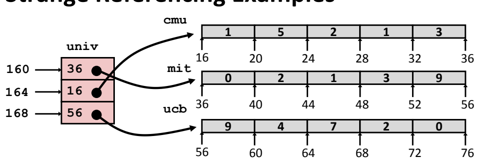
`Mem[pgh+20*index+4*dig]`   `Mem[Mem[univ+4*index]+4*dig]`

Strange Referencing Examples



- **Reference** **Address** **Value** **Guaranteed?**
  - `univ[2][3]`
  - `univ[1][5]`
  - `univ[2][-1]`
  - `univ[3][-1]`
  - `univ[1][12]`

Strange Referencing Examples



Reference	Address	Value	Guaranteed?
`univ[2][3]`	`56+4*3 = 68`	2	Yes
`univ[1][5]`	`16+4*5 = 36`	0	No
`univ[2][-1]`	`56+4*-1 = 52`	9	No
`univ[3][-1]`	`??`	??	No
`univ[1][12]`	`16+4*12 = 64`	7	No

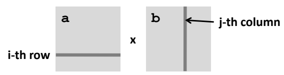
- Code does not do any bounds checking
- Ordering of elements in different arrays not guaranteed

Using Nested Arrays

- **Strengths**
  - C compiler handles doubly subscripted arrays
  - Generates very efficient code
  - Avoids multiply in index computation
- **Limitation**
  - Only works for fixed array size

```
#define N 16
typedef int fix_matrix[N][N];
```

```
/* Compute element i,k of
   fixed matrix product */
int fix_prod_ele
(fix_matrix a, fix_matrix b,
 int i, int k)
{
  int j;
  int result = 0;
  for (j = 0; j < N; j++)
    result += a[i][j]*b[j][k];
  return result;
}
```



Dynamic Nested Arrays

- **Strength**
  - Can create matrix of any size
- **Programming**
  - Must do index computation explicitly
- **Performance**
  - Accessing single element costly
  - Must do multiplication

```
int * new_var_matrix(int n)
{
  return (int *)
    calloc(sizeof(int), n*n);
}
```

```
int var_ele
  (int *a, int i, int j, int n)
{
  return a[i*n+j];
}
```

```
movl 12(%ebp),%eax        # i
movl 8(%ebp),%edx         # a
imull 20(%ebp),%eax       # n*i
addl 16(%ebp),%eax        # n*i+j
movl (%edx,%eax,4),%eax   # Mem[a+4*(i*n+j)]
```

Dynamic Array Multiplication

- **Without Optimizations**
  - Multiplies: 3
    - 2 for subscripts
    - 1 for data
  - Adds: 4
    - 2 for array indexing
    - 1 for loop index
    - 1 for data

```
/* Compute element i,k of
   variable matrix product */
int var_prod_ele
  (int *a, int *b,
   int i, int k, int n)
{
  int j;
  int result = 0;
  for (j = 0; j < n; j++)
    result +=
      a[i*n+j] * b[j*n+k];
  return result;
}
```

Optimizing Dynamic Array Multiplication

- **Optimizations**
  - Performed when set optimization level to `-O2`
- **Code Motion**
  - Expression `i*n` can be computed outside loop
- **Strength Reduction**
  - Incrementing `j` has effect of incrementing `j*n+k` by `n`
- **Operations count**
  - 4 adds, 1 mult
- **Compiler can optimize regular access patterns**

```
{
  int j;
  int result = 0;
  for (j = 0; j < n; j++)
    result +=
      a[i*n+j] * b[j*n+k];
  return result;
}
```

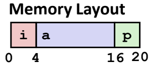
```
{
  int j;
  int result = 0;
  int iTn = i*n;
  int jTnPk = k;
  for (j = 0; j < n; j++) {
    result +=
      a[iTn+j] * b[jTnPk];
    jTnPk += n;
  }
  return result;
}
```

STRUCTURES

Structures

```
struct rec {
  int i;
  int a[3];
  int *p;
};
```

Memory Layout



- **Concept**
  - Contiguously-allocated region of memory
  - Refer to members within structure by names
  - Members may be of different types
- **Accessing Structure Member**

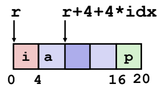
```
void
set_i(struct rec *r,
      int val)
{
  r->i = val;
}
```

IA32 Assembly

```
# %eax = val
# %edx = r
movl %eax,(%edx)    # Mem[r] = val
```

Generating Pointer to Structure Member

```
struct rec {
  int i;
  int a[3];
  int *p;
};
```



- **Generating Pointer to Array Element**
  - Offset of each structure member determined at compile time

```
int *find_a
 (struct rec *r, int idx)
{
  return &r->a[idx];
}
```

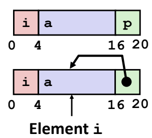
```
# %ecx = idx
# %edx = r
leal 0(,%ecx,4),%eax   # 4*idx
leal 4(%eax,%edx),%eax # r+4*idx+4
```

Structure Referencing (Cont.)

- **C Code**

```
struct rec {
  int i;
  int a[3];
  int *p;
};
```

```
void
set_p(struct rec *r)
{
  r->p =
   &r->a[r->i];
}
```




```
# %edx = r
movl (%edx),%ecx        # r->i
leal 0(,%ecx,4),%eax    # 4*(r->i)
leal 4(%edx,%eax),%eax  # r+4+4*(r->i)
movl %eax,16(%edx)      # Update r->p
```

ALIGNMENT

Alignment

- **Aligned Data**
  - Primitive data type requires K bytes
  - Address must be multiple of K
  - Required on some machines; advised on IA32
    - treated differently by IA32 Linux, x86-64 Linux, and Windows!
- **Motivation for Aligning Data**
  - Memory accessed by (aligned) chunks of 4 or 8 bytes (system dependent)
    - Inefficient to load or store datum that spans quad word boundaries
    - Virtual memory very tricky when datum spans 2 pages
- **Compiler**
  - Inserts gaps in structure to ensure correct alignment of fields

Specific Cases of Alignment (IA32)

- **1 byte: `char`, …**
  - no restrictions on address
- **2 bytes: `short`, …**
  - lowest 1 bit of address must be 0_2
- **4 bytes: `int`, `float`, `char *`, …**
  - lowest 2 bits of address must be 00_2
- **8 bytes: `double`, …**
  - Windows (and most other OS's & instruction sets):
    - lowest 3 bits of address must be 000_2
  - Linux:
    - lowest 2 bits of address must be 00_2
    - i.e., treated the same as a 4-byte primitive data type
- **12 bytes: `long double`**
  - Windows, Linux:
    - lowest 2 bits of address must be 00_2
    - i.e., treated the same as a 4-byte primitive data type

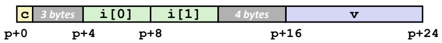
Specific Cases of Alignment (x86-64)

- **1 byte: `char`, …**
  - no restrictions on address
- **2 bytes: `short`, …**
  - lowest 1 bit of address must be 0_2
- **4 bytes: `int`, `float`, …**
  - lowest 2 bits of address must be 00_2
- **8 bytes: `double`, `char *`, …**
  - Windows & Linux:
    - lowest 3 bits of address must be 000_2
- **16 bytes: `long double`**
  - Linux:
    - lowest 3 bits of address must be 000_2
    - i.e., treated the same as a 8-byte primitive data type

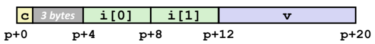
Different Alignment Conventions

```
struct S1 {
  char c;
  int i[2];
  double v;
} *p;
```

- **x86-64 or IA32 Windows:**
  - K = 8, due to `double` element



- **IA32 Linux**
  - K = 4; `double` treated like a 4-byte data type



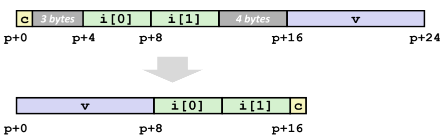
Saving Space

- **Put large data types first**

```
struct S1 {
  char c;
  int i[2];
  double v;
} *p;
```

```
struct S2 {
  double v;
  int i[2];
  char c;
} *p;
```

- **Effect (example x86-64, both have *K=8*)**



Arrays of Structures

- **Satisfy alignment requirement for every element**

```
struct S2 {
  double v;
  int i[2];
  char c;
} a[10];
```