

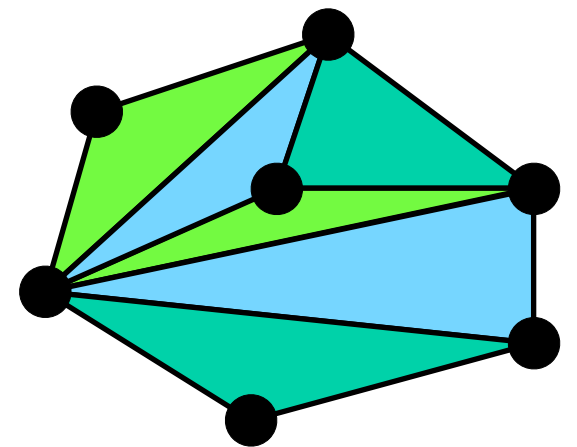
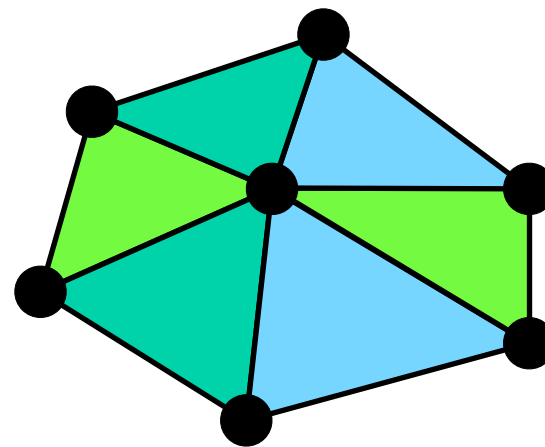
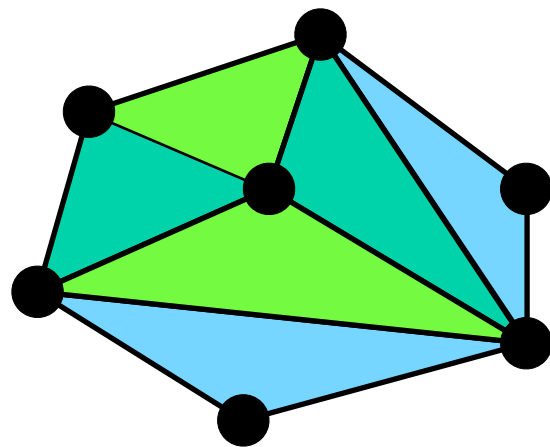


Triangulation

# Triangulations

A *triangulation* of set of points in the plane is a *partition* of the convex hull to triangles whose vertices are the points, and do not contain other points.

There are an exponential number of triangulations of a point set.

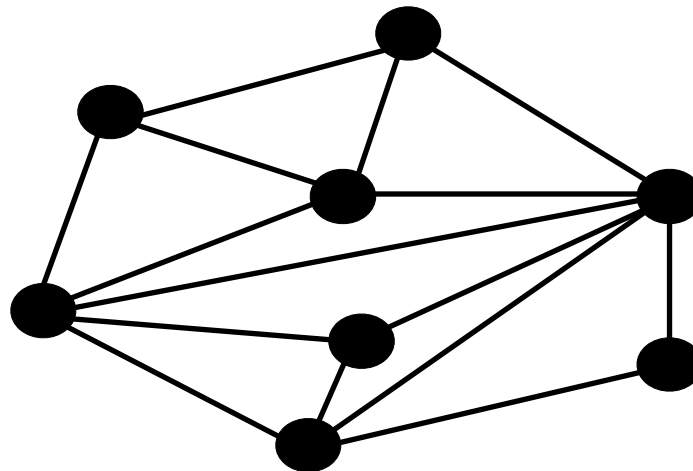


# An $O(n^3)$ Triangulation Algorithm

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Repeat until impossible:

- Select two sites.
- If the edge connecting them does not intersect previous edges, keep it.



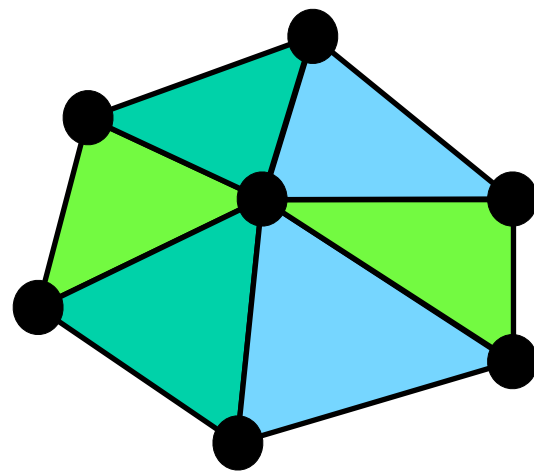
*Do you get consistent results?*

*Which ones are better?*

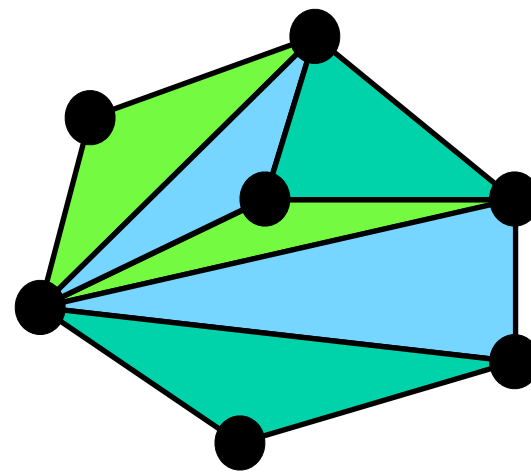
# “Quality” Triangulations

Let  $\alpha(T_i) = (\alpha_{i1}, \alpha_{i2}, \dots, \alpha_{i3})$  be the vector of angles in the triangulation  $T$  in increasing order:

- A triangulation  $T_1$  is “better” than  $T_2$  if the smallest angle of  $T_1$  is larger than the smallest angle of  $T_2$
- Delaunay triangulation is the “best” (maximizes the smallest angles)



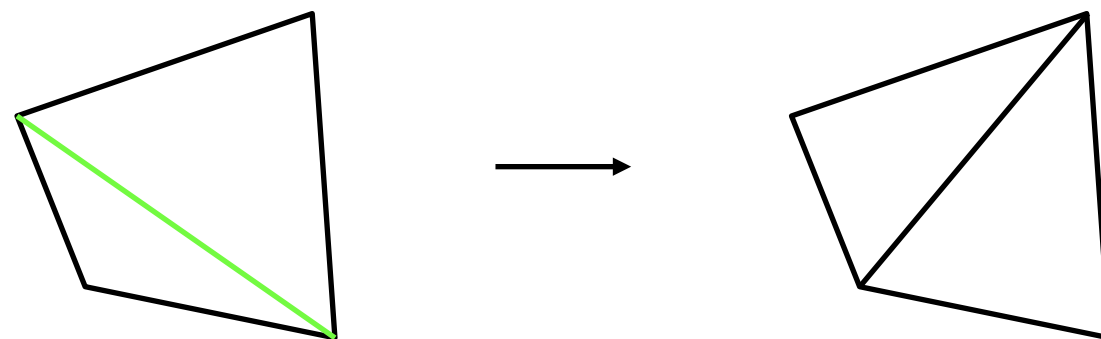
good



bad

# Improving a Triangulation

In any convex quadrangle, an *edge flip* is possible. If this flip *improves* the triangulation locally, it also improves the global triangulation.

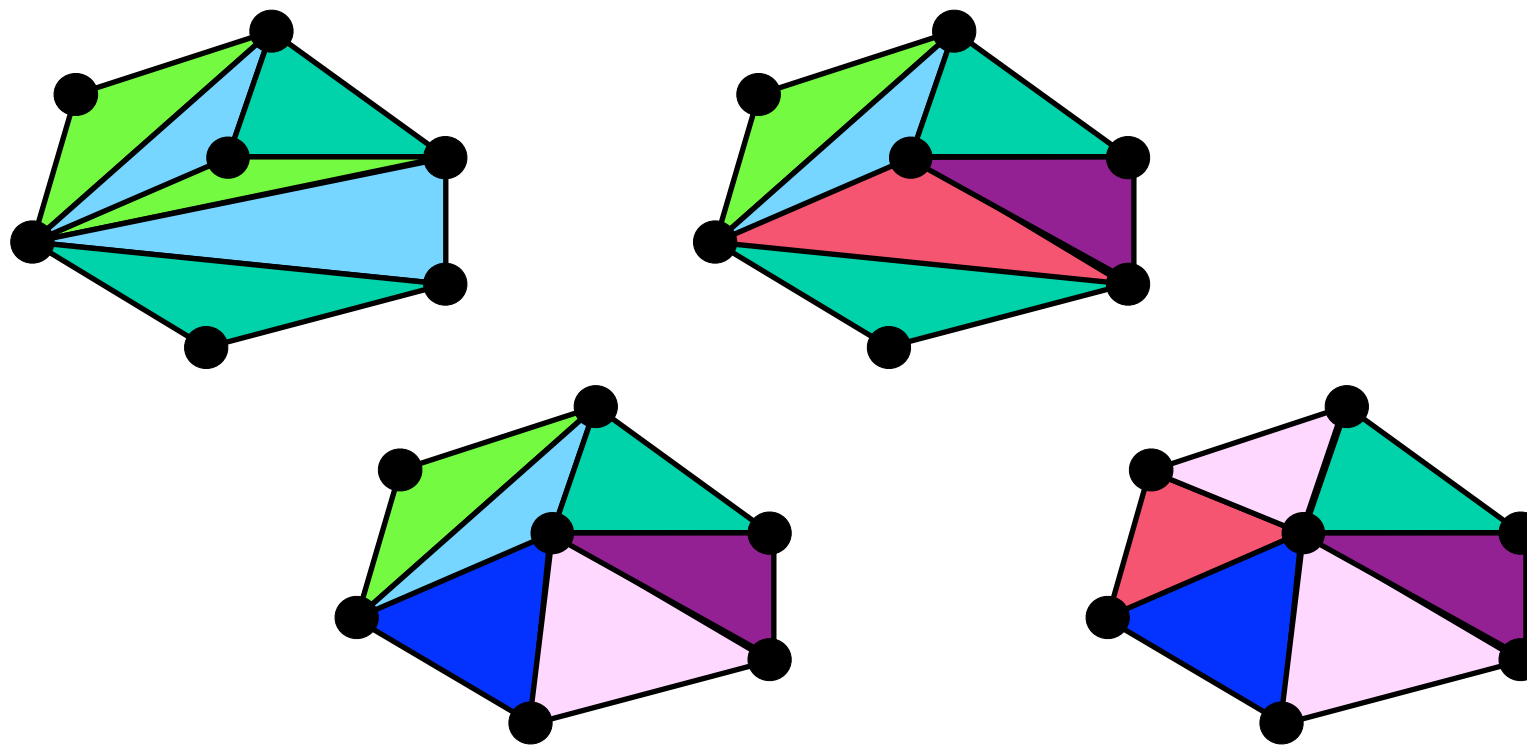


If an edge flip improves the triangulation, the first edge is called “*illegal*”.

# Naïve Delaunay Algorithm

Start with an arbitrary triangulation. Flip any illegal edge until no more exist.

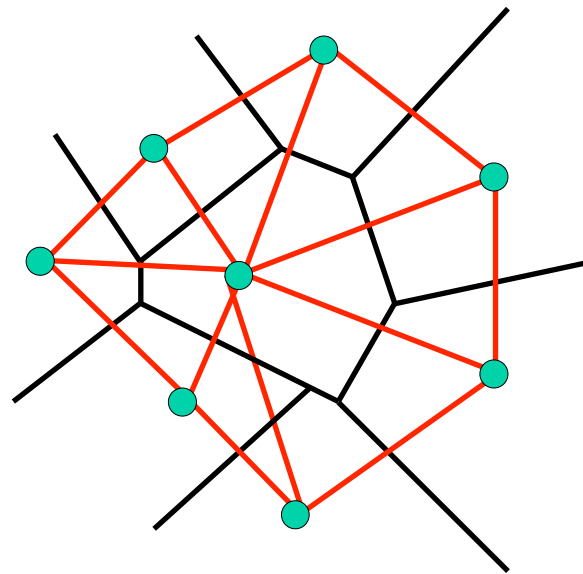
Could take a long time to terminate.



# Delaunay Triangulation by Duality

Draw the dual to the Voronoi diagram by connecting each two neighboring sites in the Voronoi diagram.

- The DT may be constructed in  $O(n \log n)$  time
- This is what Matlab's `delaunay` function uses



Demo: <http://www.cs.cornell.edu/home/chew/Delaunay.html>