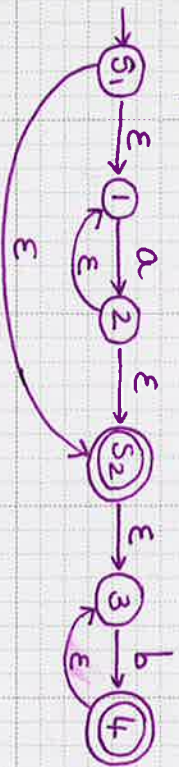


NFA to DFA algorithm: open on the transition table

example:



- about state of the DFA = ϵ -closure (start state of the NFA)

$$\epsilon\text{-closure}(S_1) = \{S_1, 1, S_2, 3\}$$

- we about the transition table with the new about state and we generate new states as we go (= the reachable states).

δ	a	b
$\{S_1, 1, S_2, 3\}$		

Now do we fill the transition table.

→ we do it in 2 steps for each cell.

Let's look at: $\delta(\{S_1, 1, S_2, 3\}, a)$ and $\delta(\{S_1, 1, S_2, 3\}, b)$.

- $\delta(\{S_1, 1, S_2, 3\}, a)$:

(1) for each state in $\{S_1, 1, S_2, 3\}$, we look at $\delta(\text{given state}, a)$:

$$\delta(S_1, a) = \emptyset \quad \text{* } S_1 \text{ only has } \epsilon \text{ transitions}$$

$$\delta(1, a) = 2$$

$$\delta(S_2, a) = \emptyset$$

$$\delta(3, a) = \emptyset$$

* S_2 only has ϵ transitions
* 3 only has b transition.

as a result, the true transitions in a from $\{S_1, 1, S_2, 3\}$ lead to $\{2\}$.

- (2) take the ϵ -closure of the set of states obtained at step (1):
 $\epsilon\text{-closure}(\{2\}) = \{2, 1, S_2, 3\}$.

we conclude that:

$$\delta(\{S_1, 1, S_2, 3\}, a) = \{2, 1, S_2, 3\}$$

- $\delta(\{S_1, 1, S_2, 3\}, b)$:

(1) for each state in $\{S_1, 1, S_2, 3\}$, we look at its transition through b :

$$\delta(S_1, b) = \emptyset$$

$$\delta(1, b) = \emptyset$$

$$\delta(S_2, b) = \emptyset$$

$$\delta(3, b) = 4$$

- (2) we take the ϵ -closure of $\{4\}$:

$$\epsilon\text{-closure}(\{4\}) = \{4, 3\}$$

we conclude that: $\delta(\{S_1, 1, S_2, 3\}, b) = \{4, 3\}$.

- Now, we have two more states, $\{2, 1, S_2, 3\}$ and $\{4, 3\}$, to add to the transition table, and we keep going by following the above 2 steps.