

# CS450 – Introduction to Networking Lecture 4 – Network Applications & HTTP

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# News: Discussion board

• Piazza

https://piazza.com/uic/spring2015/cs450/home

- There will be bonus points up to 5%
  - for Piazza forums
    - contribution of helpful code to the common good of the class (e.g. test cases and/or testing scripts)
  - Thoughtful discussions during lecture.

# A quick check on assignment 1 progress

- A. I have completed/nearly completed the project
- B. I have done about >= 75%
- C. I have done about >= 50%
- D. I have done about >= 25%
- E. I want the deadline extended

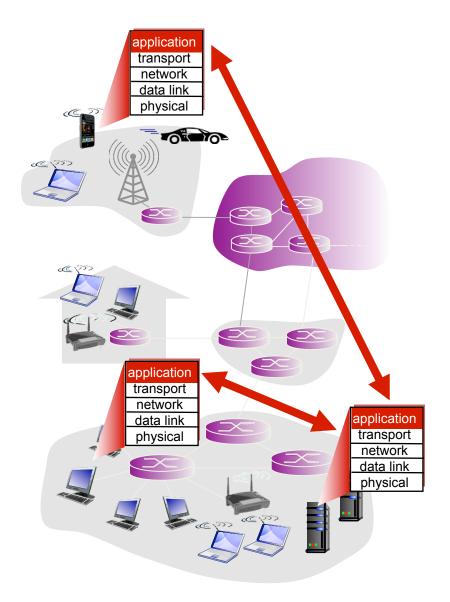
# Creating a network app

#### write programs that:

- run on (different) end systems
- communicate over network
- e.g., web server software communicates with browser software

#### no need to write software for network-core devices

- network-core devices do not run user applications
- applications on end systems allows for rapid app development, propagation

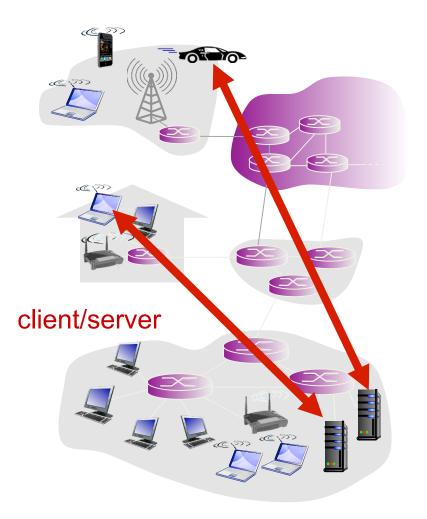


# Application architectures

possible structure of applications:

- client-server
- peer-to-peer (P2P)

# Client-server architecture



#### server:

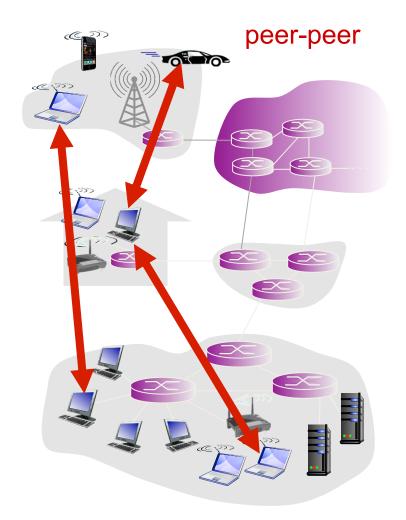
- always-on host
- permanent IP address
- data centers for scaling

#### clients:

- communicate with server
- may be intermittently connected
- may have dynamic IP addresses
- do not communicate directly with each other

# P2P architecture

- no always-on server
- arbitrary end systems directly communicate
- peers request service from other peers, provide service in return to other peers
  - self scalability new peers bring new service capacity, as well as new service demands
- peers are intermittently connected and change IP addresses
  - complex management



## Processes communicating

#### process: program running within a host

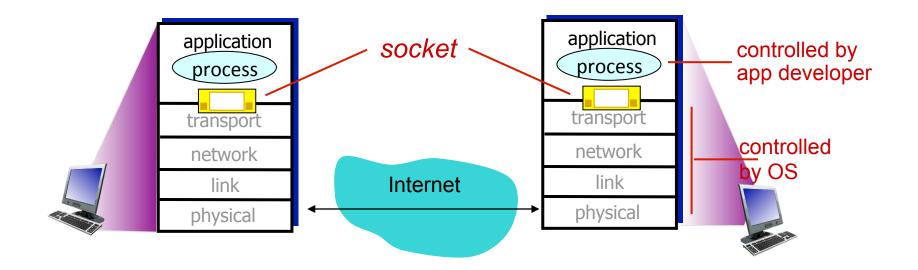
- within same host, two processes communicate using inter-process communication (defined by OS)
- processes in different hosts communicate by exchanging messages

 clients, servers
 client process: process that initiates communication
 server process: process that waits to be contacted

 aside: applications with P2P architectures have client processes & server processes

# Sockets

- process sends/receives messages to/from its socket
- socket analogous to door
  - sending process shoves message out door
  - sending process relies on transport infrastructure on other side of door to deliver message to socket at receiving process



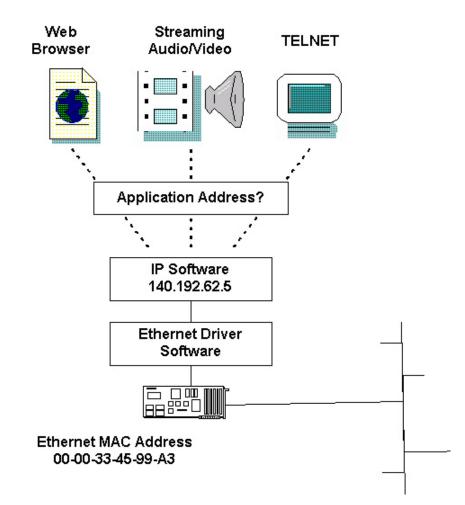
### Addressing processes

- to receive messages, process must have *identifier*
- host device has unique 32-bit IP address
- <u>Q</u>: does IP address of host on which process runs suffice for identifying the process?

A: no, many processes can be running on same host

## Addressing processes

- identifier includes both IP address and port numbers associated with process on host.
- example port numbers:
  - HTTP server: 80
  - mail server: 25
- to send HTTP message to www.cs.uic.edu web server:
  - IP address: 131.193.32.29
  - port number: 80
- more shortly...



### Process ports

- 16 bit integer
  - Port numbers: 0..65535
- 3 ranges
  - 0..1023: for common, well-known services
  - 1024..49151 Registered port: vendors use for applications
  - >49151 dynamic / private ports

# App-layer protocol defines

- types of messages exchanged,
  - e.g., request, response
- message syntax:
  - what fields in messages & how fields are delineated
- message semantics
  - meaning of information in fields
- rules for when and how processes send & respond to messages

#### open protocols:

- defined in RFCs
- allows for interoperability
- e.g., HTTP, SMTP

proprietary protocols:

• e.g., Skype

### What transport service does an app need?

#### data integrity

- some apps (e.g., file transfer, web transactions) require
   100% reliable data transfer
- other apps (e.g., audio) can tolerate some loss

#### timing

 some apps (e.g., Internet telephony, interactive games) require low delay to be "effective"

#### throughput

- some apps (e.g., multimedia) require minimum amount of throughput to be "effective"
- other apps ("elastic apps")
  make use of whatever
  throughput they get

#### security

. . .

encryption, data integrity,

#### Transport service requirements: common apps

	application	data loss	throughput	time sensitive
	file transfer	no loss	elastic	no
	e-mail	no loss	elastic	no
V	Veb documents	no loss	elastic	no
real-ti	me audio/video	loss-tolerant	audio: 5kbps-1Mbps video:10kbps-5Mbps	
sto	red audio/video	loss-tolerant	same as above	
	eractive games	loss-tolerant	few kbps up	yes, few secs
	text messaging	no loss	elastic	yes, 100' s
				msec yes and no

### Internet transport protocols services

#### TCP service:

- reliable transport between sending and receiving process
- *flow control:* sender won't overwhelm receiver
- congestion control: throttle sender when network overloaded
- does not provide: timing, minimum throughput guarantee, security
- connection-oriented: setup required between client and server processes

#### UDP service:

- unreliable data transfer between sending and receiving process
- does not provide: reliability, flow control, congestion control, timing, throughput guarantee, security, or connection setup,

<u>Q</u>: why bother? Why is there a UDP?

#### Internet apps: application, transport protocols

application	application layer protocol	underlying transport protocol
e-mail	SMTP [RFC 2821]	TCP
remote terminal access	Telnet [RFC 854]	TCP
Web	HTTP [RFC 2616]	TCP
file transfer	FTP [RFC 959]	TCP
streaming multimedia	HTTP (e.g., YouTube), RTP [RFC 1889]	TCP or UDP
Internet telephony	SIP, RTP, proprietary (e.g., Skype)	TCP or UDP

# Securing TCP

#### TCP & UDP

- ✤ no encryption
- cleartext passwds sent into socket traverse Internet in cleartext

#### SSL

- provides encrypted TCP connection
- ✤ data integrity
- end-point authentication

#### SSL is at app layer

 Apps use SSL libraries, which "talk" to TCP

#### SSL socket API

- cleartext passwds sent
  into socket traverse
  Internet encrypted
- See Chapter 7

# Select a wrong statement

- A. TCP and UDP are in transport layer
- B. TCP provides reliable connections while UDP is unreliable connections
- C. HTTP uses TCP
- D. FTP uses UDP
- E. All network applications rely on either TCP or UDP

# A network application process is identified by

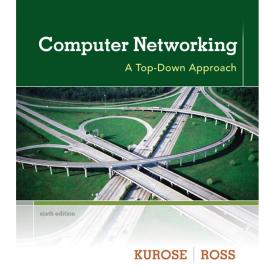
- A. MAC address
- B. IP address
- C. IP address and port number
- D. Domain name and IP address
- E. MAC address and port number

### UDP is unreliable, why we still need it?

- A. Because TCP is bad
- B. Because UDP is good
- C. Because in some cases, UDP is better than TCP
- D. Because UDP is speedy
- E. C and D are both correct

## Next lecture

- HTTP (cont')
  - Readings 2.2
- Questions regarding Assignment 1?



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Computer Networking: A Top Down Approach 6<sup>th</sup> edition Jim Kurose, Keith Ross Addison-Wesley March 2012