

CS450 – Introduction to Networking

Lecture 5 – HTTP

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Live demo with HTTP (client side) using telnet

1. Telnet to your favorite Web server:

```
telnet www.cs.uic.edu 80
```

opens TCP connection to port 80
(default HTTP server port) at www.cs.uic.edu.
anything typed in sent
to port 80 at www.cs.uic.edu

2. type in a GET HTTP request:

```
GET /~phu/cs450.html HTTP/1.0  
Host: www.cs.uic.edu
```

by typing this in (hit carriage
return twice), you send
this minimal (but complete)
GET request to HTTP server

3. look at response message sent by HTTP server!

Q: To save data from
a HTTP response, you
should save from:

```
PhuMAC:~ phu$ telnet www.cs.uic.edu 80
Trying 131.193.32.29...
Connected to www.cs.uic.edu.
Escape character is '^J'.
```

A → GET /~phu/cs450.html HTTP/1.0
Host: www.cs.uic.edu

B → HTTP/1.1 200 OK

C → Date: Fri, 23 Jan 2015 05:31:17 GMT
Server: Apache/2.2.3 (Red Hat)
Last-Modified: Tue, 20 Jan 2015 17:05:17 GMT
ETag: "6f60b27-121-716e8d40"
Accept-Ranges: bytes
Content-Length: 289
Content-Type: text/html; charset=UTF-8
Connection: close

D →

E → <HTML>
<BODY>
 <h1>Welcome! </h1>
 This is just a test page for HTTP protocol.
</BODY>
</HTML>
Connection closed by foreign host.

Live demo with wireshark

possible structure of applications:

1. Install wireshark: `$sudo apt-get install`
(wireshark link on the course's homepage)
2. Run: `$sudo wireshark`
 1. Select the connected network interface (In VirtualBox, configure Bridge Adaptor)
3. Open a web browser and browse a URL
4. Examine the captured HTTP request/
response in wireshark

HTTP overview

uses TCP:

- client initiates TCP connection (creates socket) to server, port 80
- server accepts TCP connection from client
- HTTP messages (application-layer protocol messages) exchanged between browser (HTTP client) and Web server (HTTP server)
- TCP connection closed

HTTP is “stateless”

- server maintains no information about past client requests

aside
protocols that maintain
“state” are complex!

- ❖ past history (state) must be maintained
- ❖ if server/client crashes, their views of “state” may be inconsistent, must be reconciled

HTTP connections

non-persistent HTTP

- at most one object sent over TCP connection
 - connection then closed
- downloading multiple objects required multiple connections

persistent HTTP

- multiple objects can be sent over single TCP connection between client, server

Non-persistent HTTP

suppose user enters URL:

`www.someSchool.edu/someDepartment/home.index`

(contains text,
references to 10
jpeg images)

1a. HTTP client initiates TCP

connection to HTTP server
(process) at
`www.someSchool.edu` on port
80

1b. HTTP server at host

`www.someSchool.edu` waiting for
TCP connection at port 80.
“accepts” connection, notifying
client

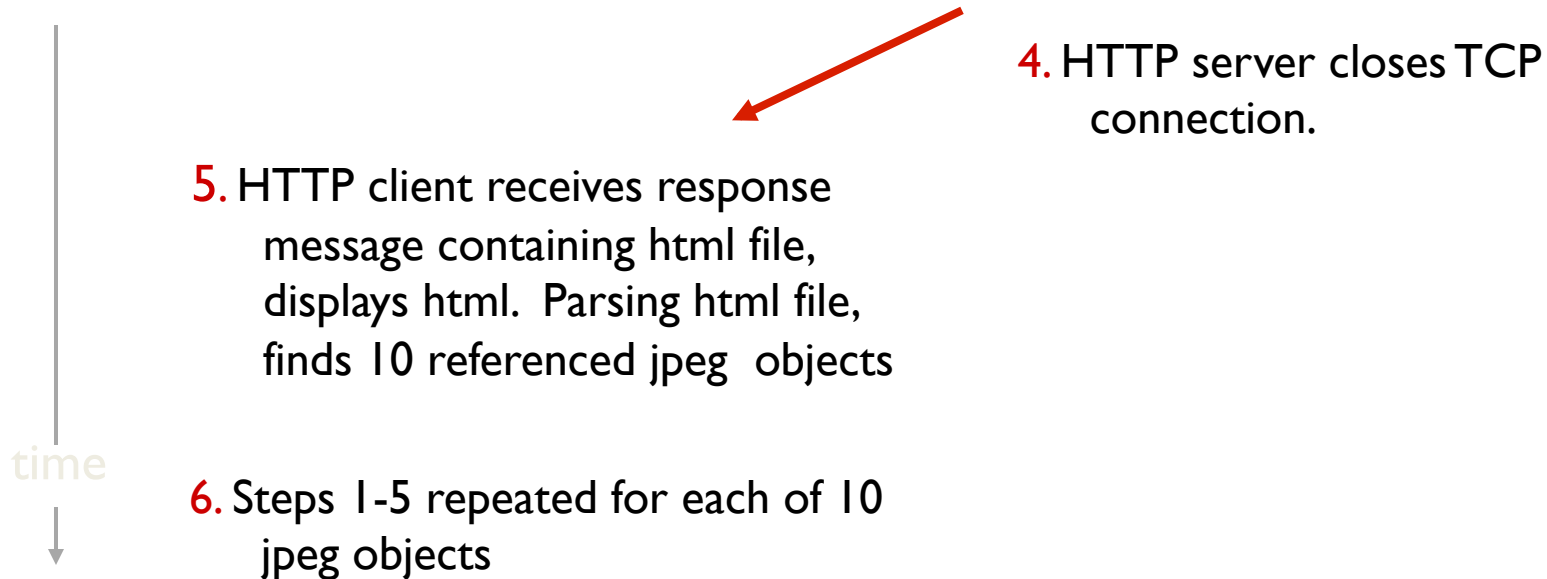
2. HTTP client sends HTTP *request
message* (containing URL) into TCP
connection socket. Message
indicates that client wants object
`someDepartment/home.index`

3. HTTP server receives request
message, forms *response message*
containing requested object, and
sends message into its socket

time



Non-persistent HTTP (cont.)



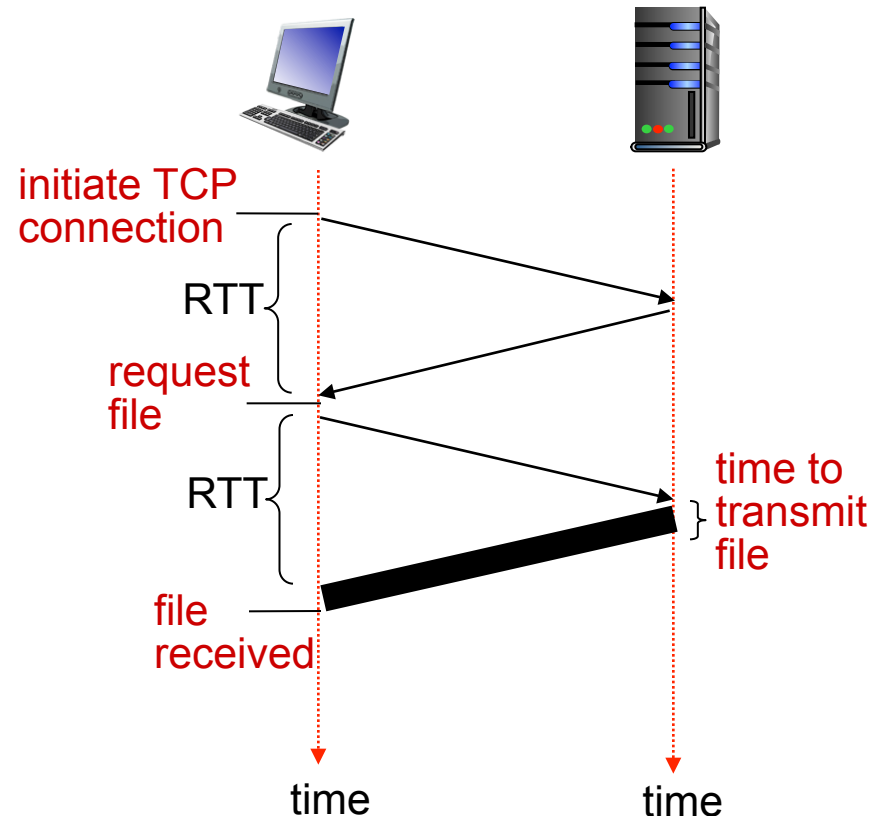
Non-persistent HTTP: response time

RTT (definition): time for a small packet to travel from client to server and back

HTTP response time:

- one RTT to initiate TCP connection
- one RTT for HTTP request and first few bytes of HTTP response to return
- file transmission time
- non-persistent HTTP response time =

$2RTT + \text{file transmission time}$



Persistent HTTP

non-persistent HTTP issues:

- requires 2 RTTs per object
- OS overhead for *each* TCP connection
- browsers often open parallel TCP connections to fetch referenced objects

persistent HTTP:

- server leaves connection open after sending response
- subsequent HTTP messages between same client/server sent over open connection
- client sends requests as soon as it encounters a referenced object
- as little as one RTT for all the referenced objects

Select a correct statement

- A. Persistent HTTP uses more connections.
- B. Persistent HTTP sends fewer HTTP requests
- C. Persistent HTTP sends more HTTP requests
- D. Persistent HTTP uses less connections to improve download time
- E. Both B and D are correct

HTTP request message

- two types of HTTP messages: *request, response*
- **HTTP request message:**
 - ASCII (human-readable format)

request line
(GET, POST,
HEAD commands)

header
lines

carriage return,
line feed at start
of line indicates
end of header lines

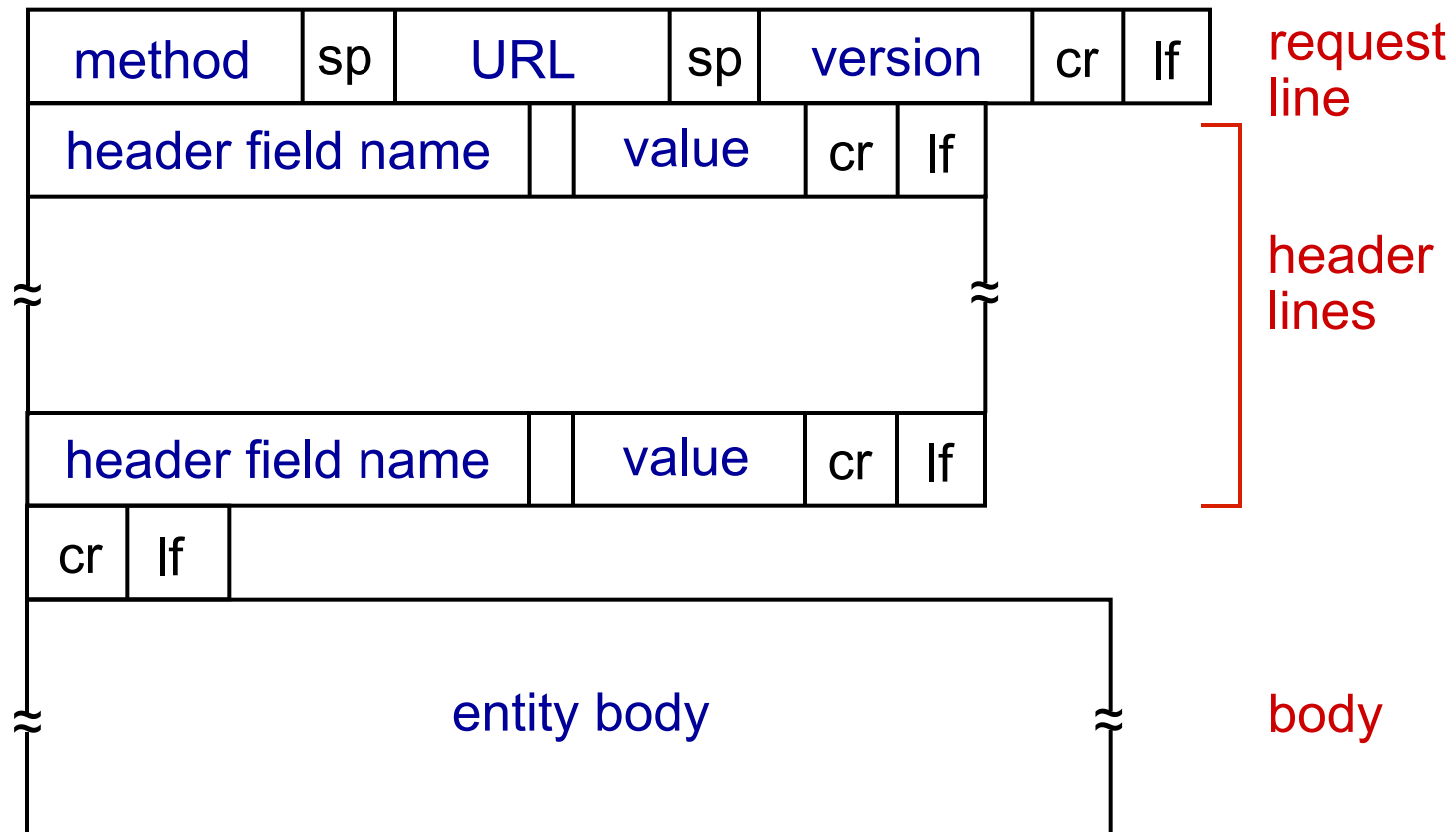
```
GET /index.html HTTP/1.1\r\n
Host: www-net.cs.umass.edu\r\n
User-Agent: Firefox/3.6.10\r\n
Accept: text/html,application/xhtml+xml\r\n
Accept-Language: en-us,en;q=0.5\r\n
Accept-Encoding: gzip,deflate\r\n
Accept-Charset: ISO-8859-1,utf-8;q=0.7\r\n
Keep-Alive: 115\r\n
Connection: keep-alive\r\n
\r\n
```

carriage return character
line-feed character

Why \r\n (CRLF) is used in HTTP request?

- A. Just unnecessary protocol format
- B. They are automatically generated when user hits ENTER
- C. They are used to indicate the end of a header field or section
- D. None of above is correct

HTTP request message: general format



Uploading form input

POST method:

- web page often includes form input
- input is uploaded to server in entity body

URL method:

- uses GET method
- input is uploaded in URL field of request line:

`www.somesite.com/animalsearch?monkeys&banana`

Method types

HTTP/1.0:

- GET
- POST
- HEAD
 - asks server to leave requested object out of response

HTTP/1.1:

- GET, POST, HEAD
- PUT
 - uploads file in entity body to path specified in URL field
- DELETE
 - deletes file specified in the URL field

HTTP response message

status line
(protocol
status code
status phrase)

header
lines

data, e.g.,
requested
HTML file

```
HTTP/1.1 200 OK\r\n
Date: Sun, 26 Sep 2010 20:09:20 GMT\r\n
Server: Apache/2.0.52 (CentOS)\r\n
Last-Modified: Tue, 30 Oct 2007 17:00:02 GMT\r\n
ETag: "17dc6-a5c-bf716880"\r\n
Accept-Ranges: bytes\r\n
Content-Length: 2652\r\n
Keep-Alive: timeout=10, max=100\r\n
Connection: Keep-Alive\r\n
Content-Type: text/html; charset=ISO-8859-1\r\n
\r\n
data data data data data ...
```

HTTP response status codes

- ❖ status code appears in 1st line in server-to-client response message.
- ❖ some sample codes:

200 OK

- request succeeded, requested object later in this msg

301 Moved Permanently

- requested object moved, new location specified later in this msg (Location:)

400 Bad Request

- request msg not understood by server

404 Not Found

- requested document not found on this server

505 HTTP Version Not Supported

Different features between HTTP 1.0 vs 1.1

- Extensibility
- Caching
- Bandwidth optimization
- Network connection management
- Message transmission
- Internet address conservation
- Error notification
- Security, integrity, and authentication
- Content negotiation

Reference: <http://www8.org/w8-papers/5c-protocols/key/key.html>

Why cookies are used in HTTP

- A. It is required by law
- B. To improve non-persistent HTTP
- C. Because HTTP is stateless
- D. All above are correct

User-server state: cookies

many Web sites use cookies

four components:

- 1) cookie header line of HTTP *response* message
- 2) cookie header line in next HTTP *request* message
- 3) cookie file kept on user's host, managed by user's browser
- 4) back-end database at Web site

example:

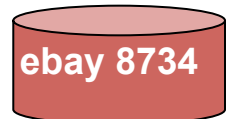
- Susan always access Internet from PC
- visits specific e-commerce site for first time
- when initial HTTP requests arrives at site, site creates:
 - unique ID
 - entry in backend database for ID

Cookies: keeping “state” (cont.)

client



server



cookie file



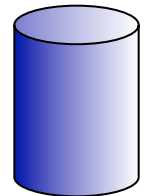
usual http request msg

Amazon server
creates ID
1678 for user

usual http response
set-cookie: 1678

create
entry

backend
database



usual http request msg
cookie: 1678

cookie-
specific
action

access

usual http response msg

access

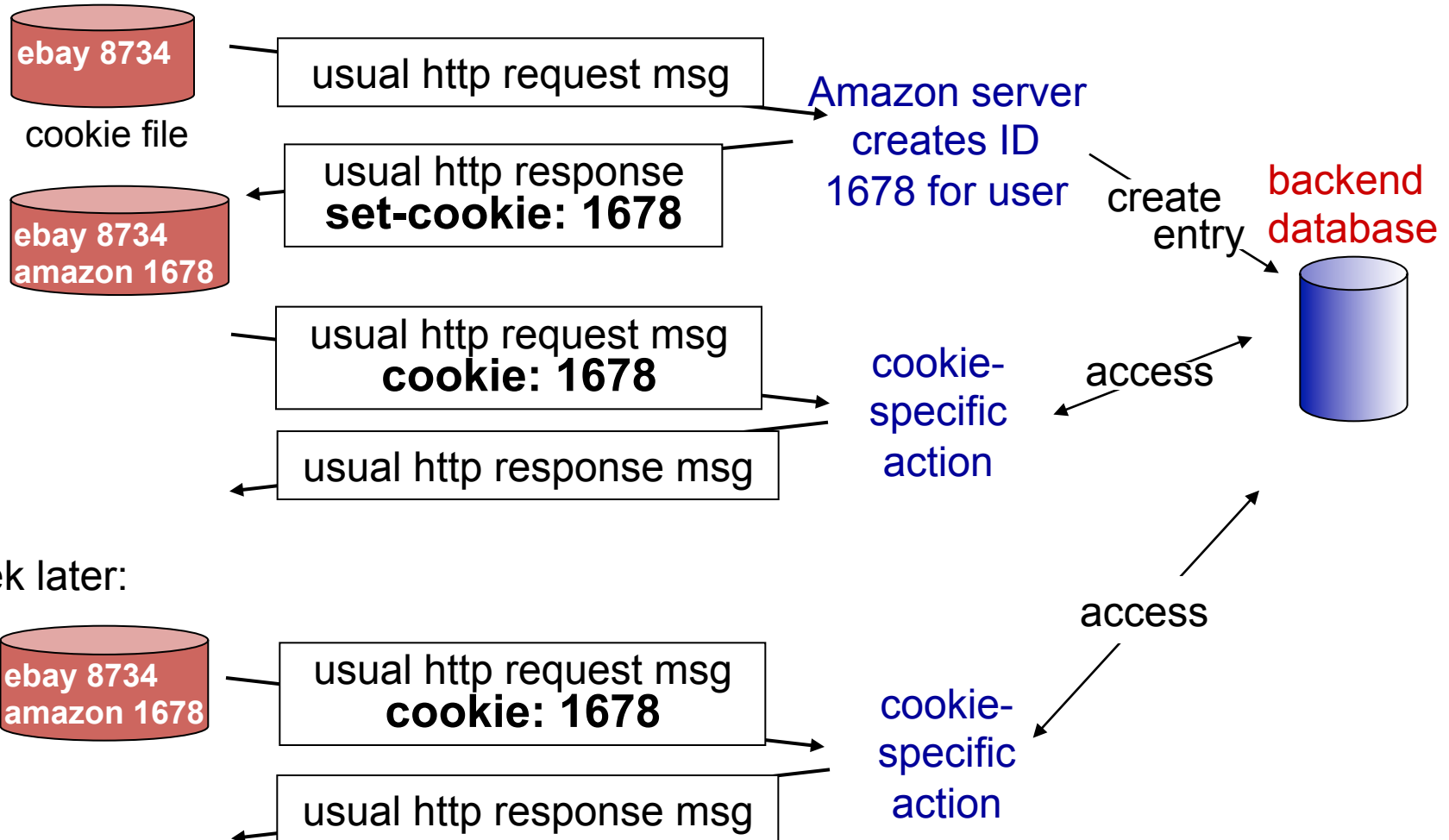
cookie-
specific
action

one week later:



usual http request msg
cookie: 1678

usual http response msg



Cookies (continued)

what cookies can be used for:

- authorization
- shopping carts
- recommendations
- user session state (Web e-mail)

cookies and privacy: aside

- ❖ cookies permit sites to learn a lot about you
- ❖ you may supply name and e-mail to sites

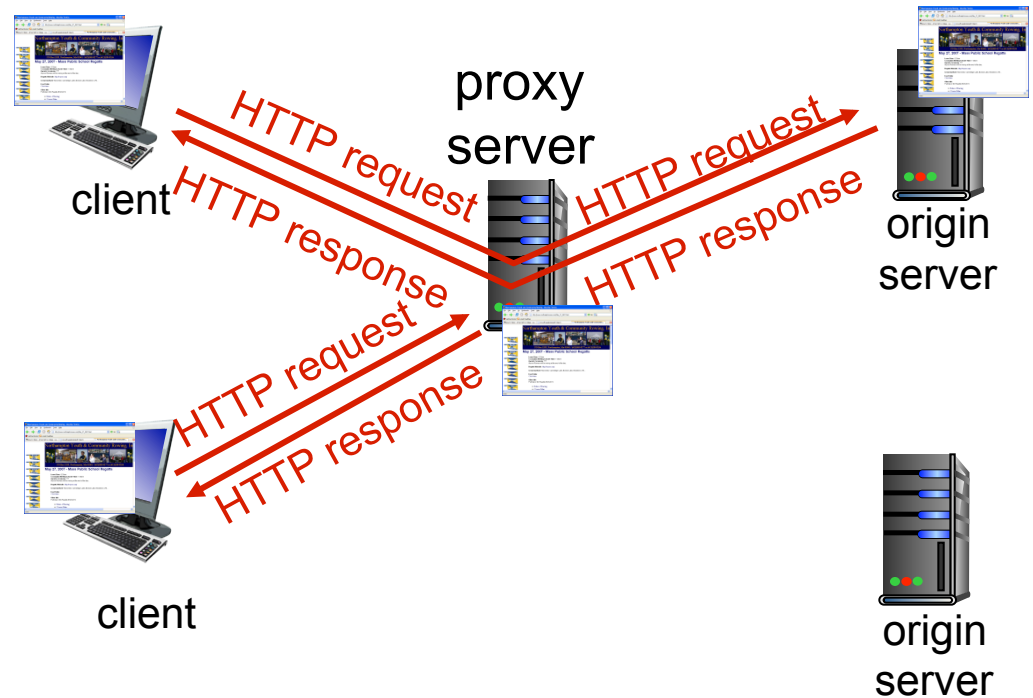
how to keep “state”:

- ❖ protocol endpoints: maintain state at sender/receiver over multiple transactions
- ❖ cookies: http messages carry state

Web caches (proxy server)

goal: satisfy client request without involving origin server

- user sets browser: Web accesses via cache
- browser sends all HTTP requests to cache
 - object in cache: cache returns object
 - else cache requests object from origin server, then returns object to client



More about Web caching

- cache acts as both client and server
 - server for original requesting client
 - client to origin server
- typically cache is installed by ISP (university, company, residential ISP)

why Web caching?

- reduce response time for client request
- reduce traffic on an institution's access link
- Internet dense with caches: enables “poor” content providers to effectively deliver content (so too does P2P file sharing)

Conditional GET

Let's see in
wireshark

client



server



- **Goal:** don't send object if cache has up-to-date cached version
 - no object transmission delay
 - lower link utilization
- **cache:** specify date of cached copy in HTTP request

If-modified-since:
<date>

- **server:** response contains no object if cached copy is up-to-date:

**HTTP/1.0 304 Not
Modified**

HTTP request msg
If-modified-since: <date>

object
not
modified
before
<date>

HTTP response
**HTTP/1.0
304 Not Modified**

HTTP request msg
If-modified-since: <date>

object
modified
after
<date>

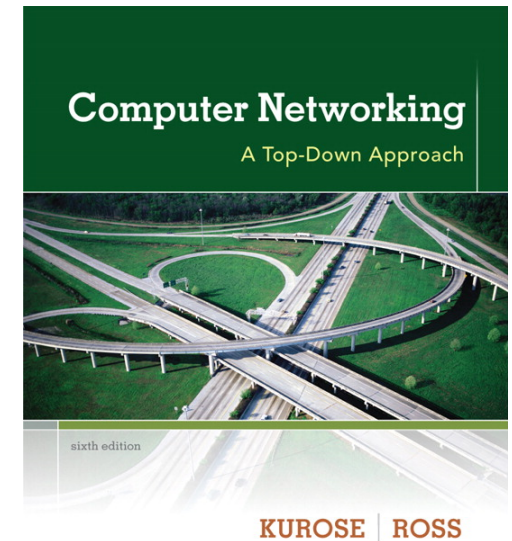
HTTP response
**HTTP/1.0 200 OK
<data>**

A quick check on assignment 1 progress

- A. I have completed/nearly completed the project
- B. I have done about $\geq 75\%$
- C. I have done about $\geq 50\%$
- D. I have done about $\geq 25\%$
- E. I want the deadline extended

Next lecture

- Client-Server Model and Assignment 2
- SMTP
 - Readings 2.4
- Questions regarding Assignment 1?
- Assignment 2 will be posted by next lecture



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6th edition

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