

CS450 – Introduction to Networking

Lecture 14 – TCP

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Feb 13, 2015

Next lecture (Feb 16)

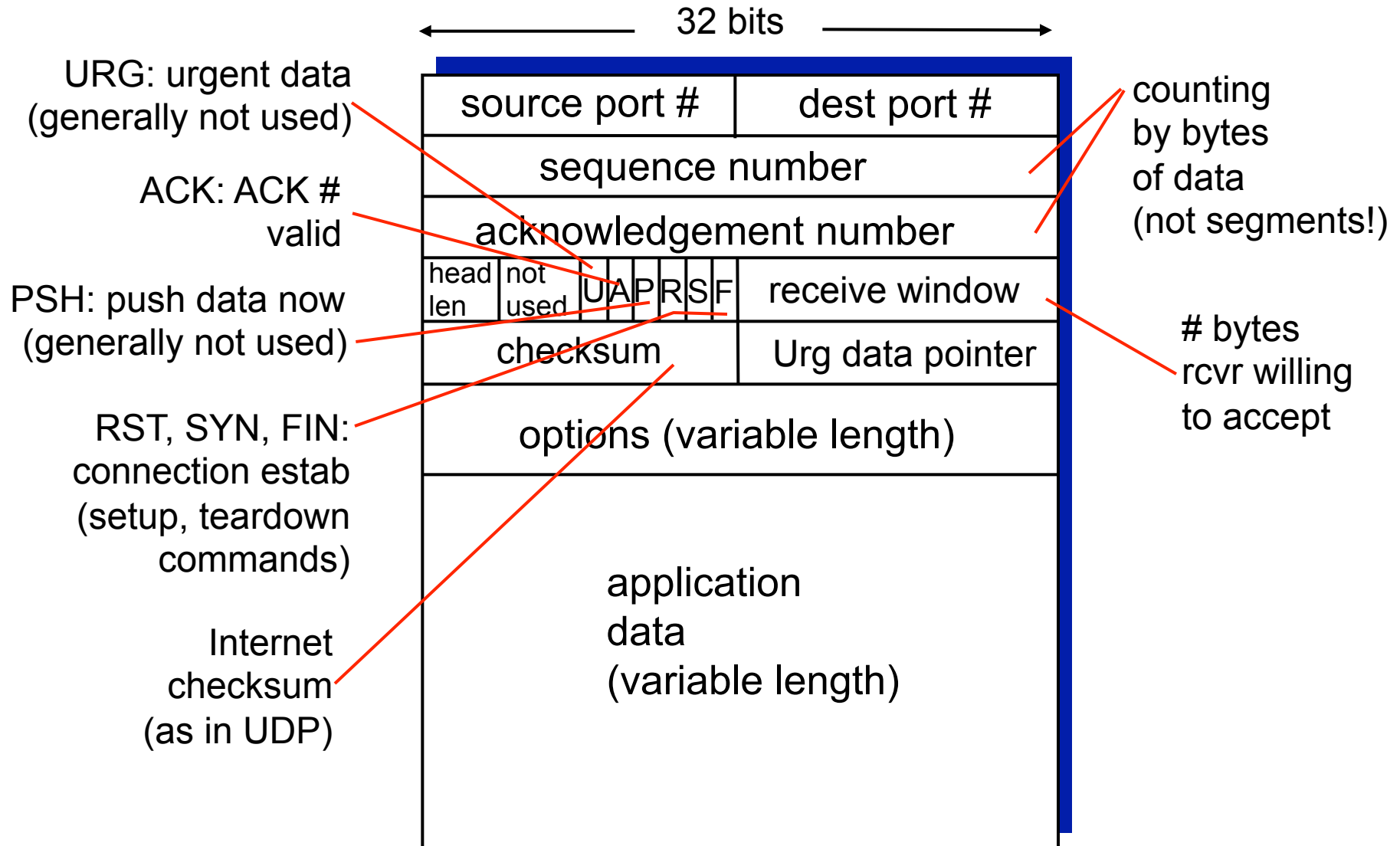
- Assignment 3 (No iclicker questions)
 - Wireshark links
- Guest lecture on Monday Feb 23rd
 - DNS Security
- Midterm exam in class
 - 1 PM Friday, March 6th

TCP: Overview

RFCs: 793, 1122, 1323, 2018, 2581

- **point-to-point:**
 - one sender, one receiver
- **reliable, in-order *byte stream*:**
 - no “message boundaries”
- **pipelined:**
 - TCP congestion and flow control set window size
- **full duplex data:**
 - bi-directional data flow in same connection
 - MSS: maximum segment size
- **connection-oriented:**
 - handshaking (exchange of control msgs) initiates sender, receiver state before data exchange
- **flow controlled:**
 - sender will not overwhelm receiver

TCP segment structure



TCP seq. numbers, ACKs

sequence numbers:

- byte stream “number” of first byte in segment’s data

acknowledgements:

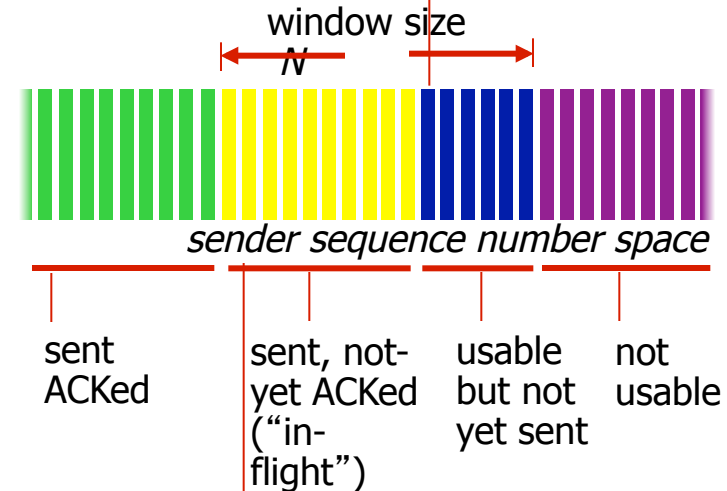
- seq # of next byte expected from other side
- cumulative ACK

Q: how receiver handles out-of-order segments

- A: TCP spec doesn’t say, - up to implementor

outgoing segment from sender

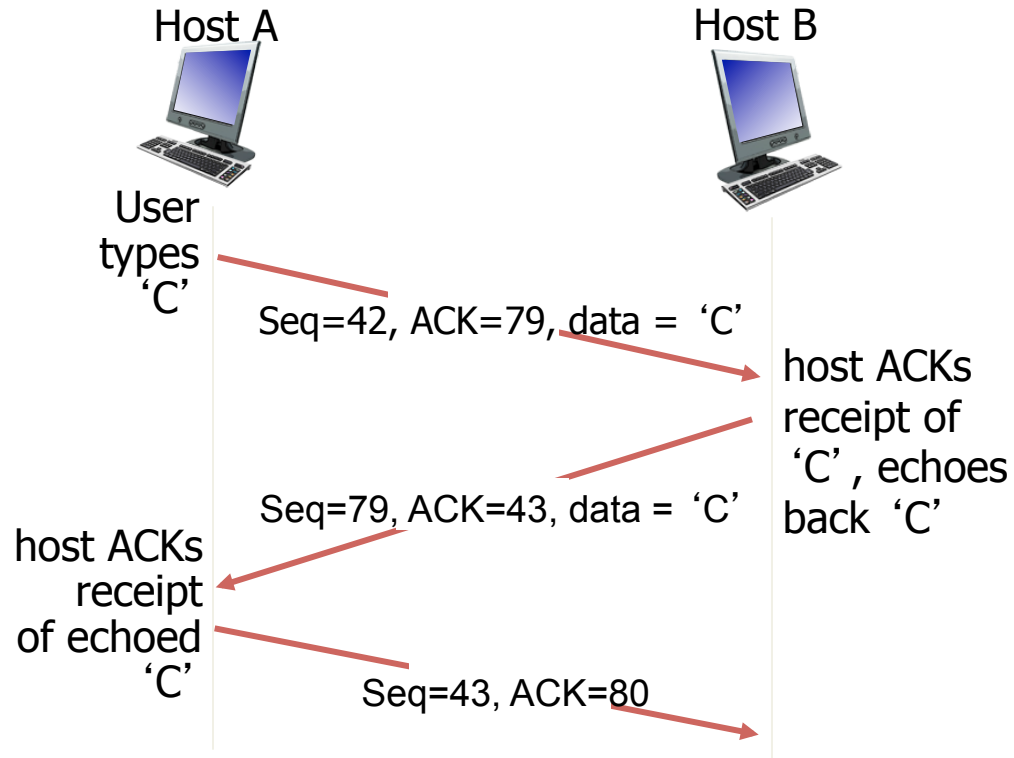
source port #	dest port #
sequence number	
acknowledgement number	
	rwnd
checksum	urg pointer



incoming segment to sender

source port #	dest port #
sequence number	
acknowledgement number	
	A
checksum	urg pointer

TCP seq. numbers, ACKs



simple telnet scenario

Sender sends seq#=101, with data length = 20 bytes, and receives ack = 91. What is true?

- A. The receiver has successfully received up to seq# 90
- B. The packet with seq#=101 was lost
- C. The packet with seq#=91 was lost
- D. A and B
- E. A and C

TCP round trip time, timeout

Q: how to set TCP timeout value?

- longer than RTT
 - but RTT varies
- *too short*: premature timeout, unnecessary retransmissions
- *too long*: slow reaction to segment loss

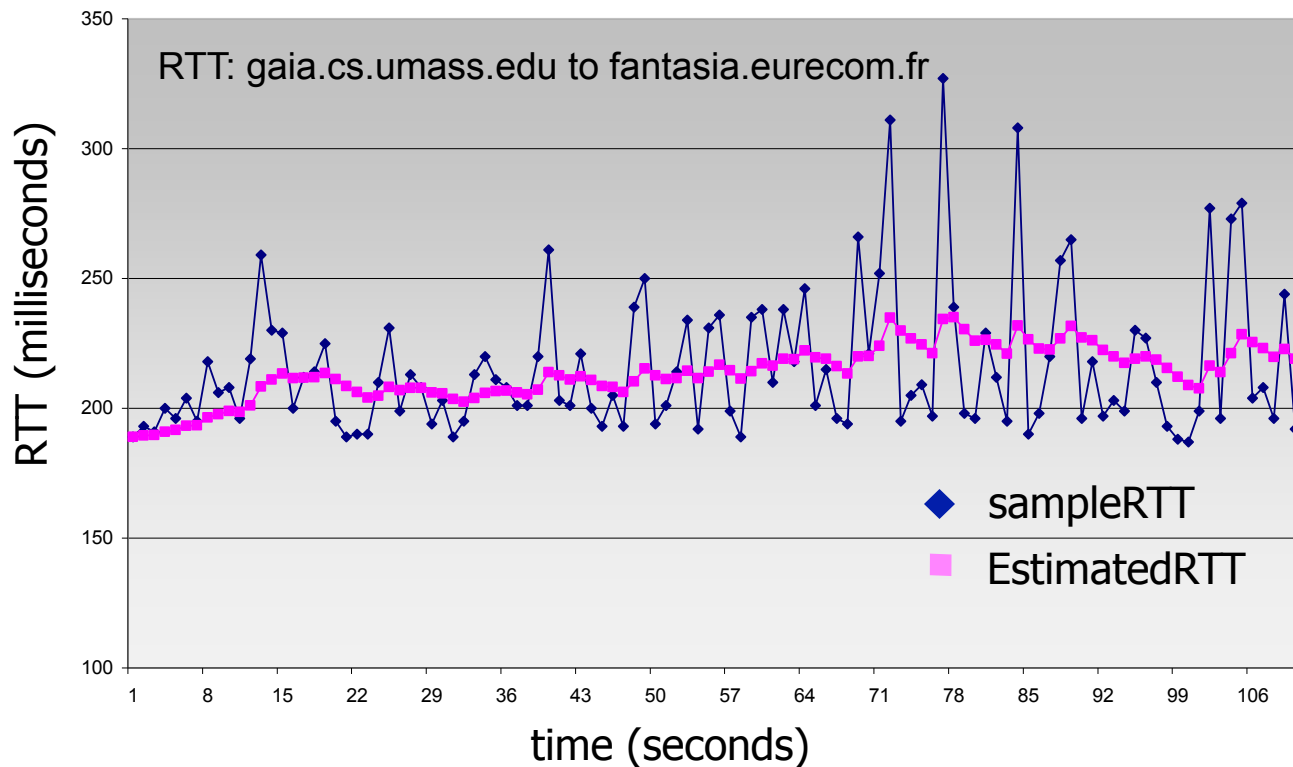
Q: how to estimate RTT?

- **SampleRTT**: measured time from segment transmission until ACK receipt
 - ignore retransmissions
- **SampleRTT** will vary, want estimated RTT “smoother”
 - average several *recent* measurements, not just current **SampleRTT**

TCP round trip time, timeout

$$\text{EstimatedRTT} = (1 - \alpha) * \text{EstimatedRTT} + \alpha * \text{SampleRTT}$$

- ❖ exponential weighted moving average
- ❖ influence of past sample decreases exponentially fast
- ❖ typical value: $\alpha = 0.125$




TCP round trip time, timeout

- **timeout interval:** `EstimatedRTT` plus “safety margin”
 - large variation in `EstimatedRTT` → larger safety margin
- estimate `SampleRTT` deviation from `EstimatedRTT`:

$$\text{DevRTT} = (1-\beta) * \text{DevRTT} + \beta * |\text{SampleRTT} - \text{EstimatedRTT}|$$

(typically, $\beta = 0.25$)

$$\text{TimeoutInterval} = \text{EstimatedRTT} + 4 * \text{DevRTT}$$

 ↑
estimated RTT ↑
“safety margin”

What is DevRTT?

- A. Relative difference between SampleRTT and EstimatedRTT
- B. Is used as a parameter for safe margin of timeout interval
- C. Is used to re-calculate SampleRTT
- D. A and B
- E. A and C

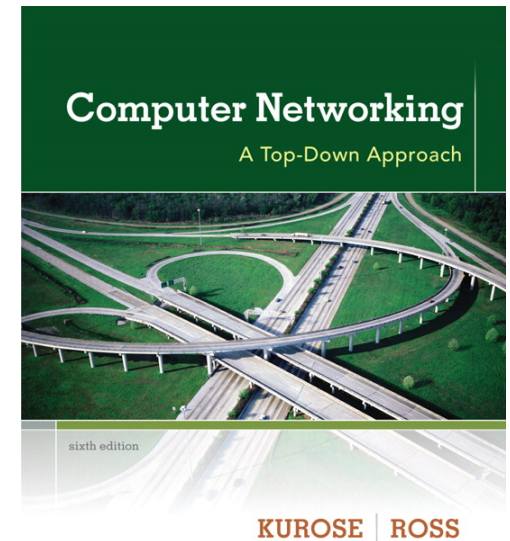
Next lecture

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 - Readings 3.5
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Computer Networking: A Top Down Approach

6th edition

Jim Kurose, Keith Ross

Addison-Wesley

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