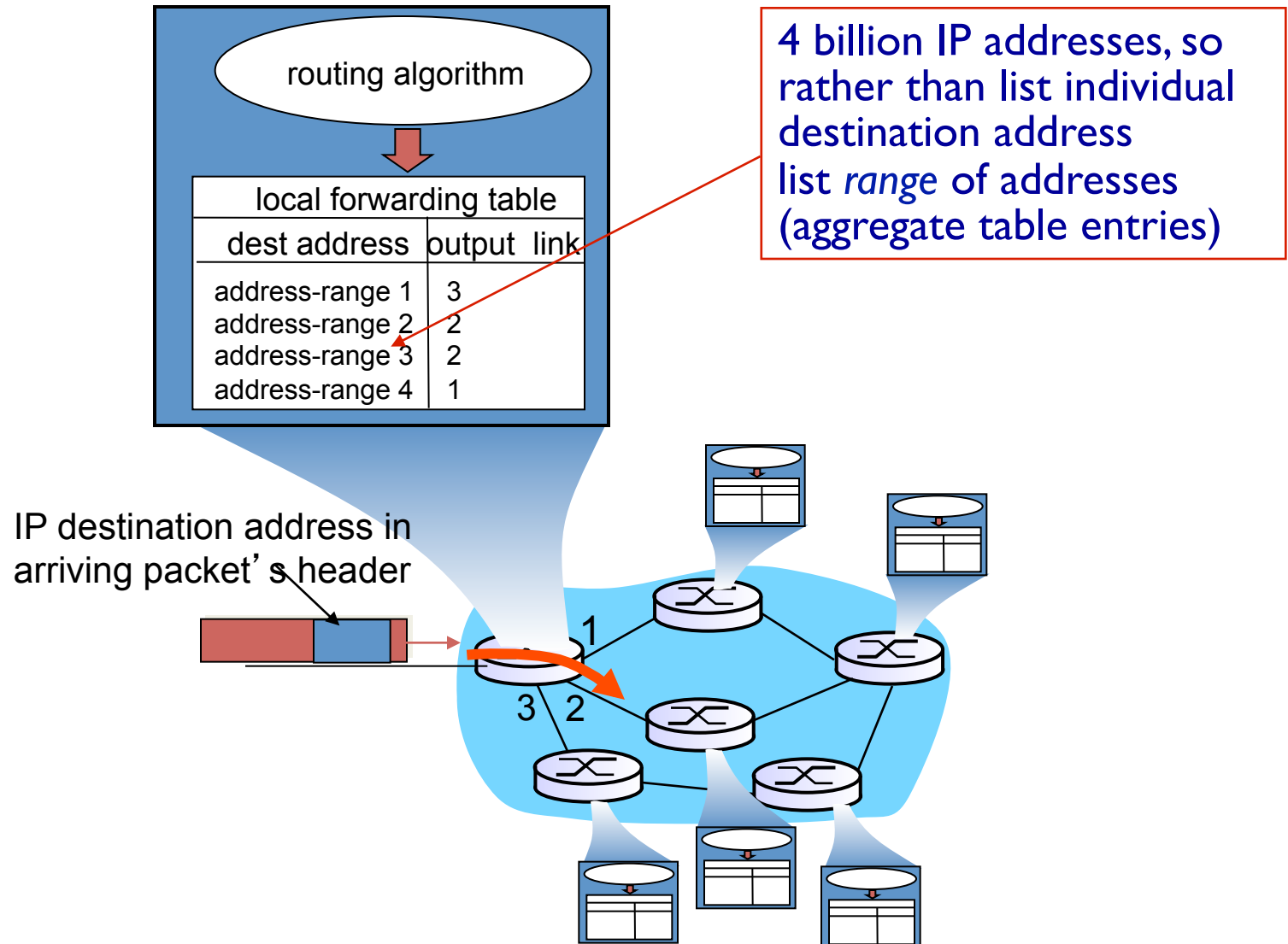


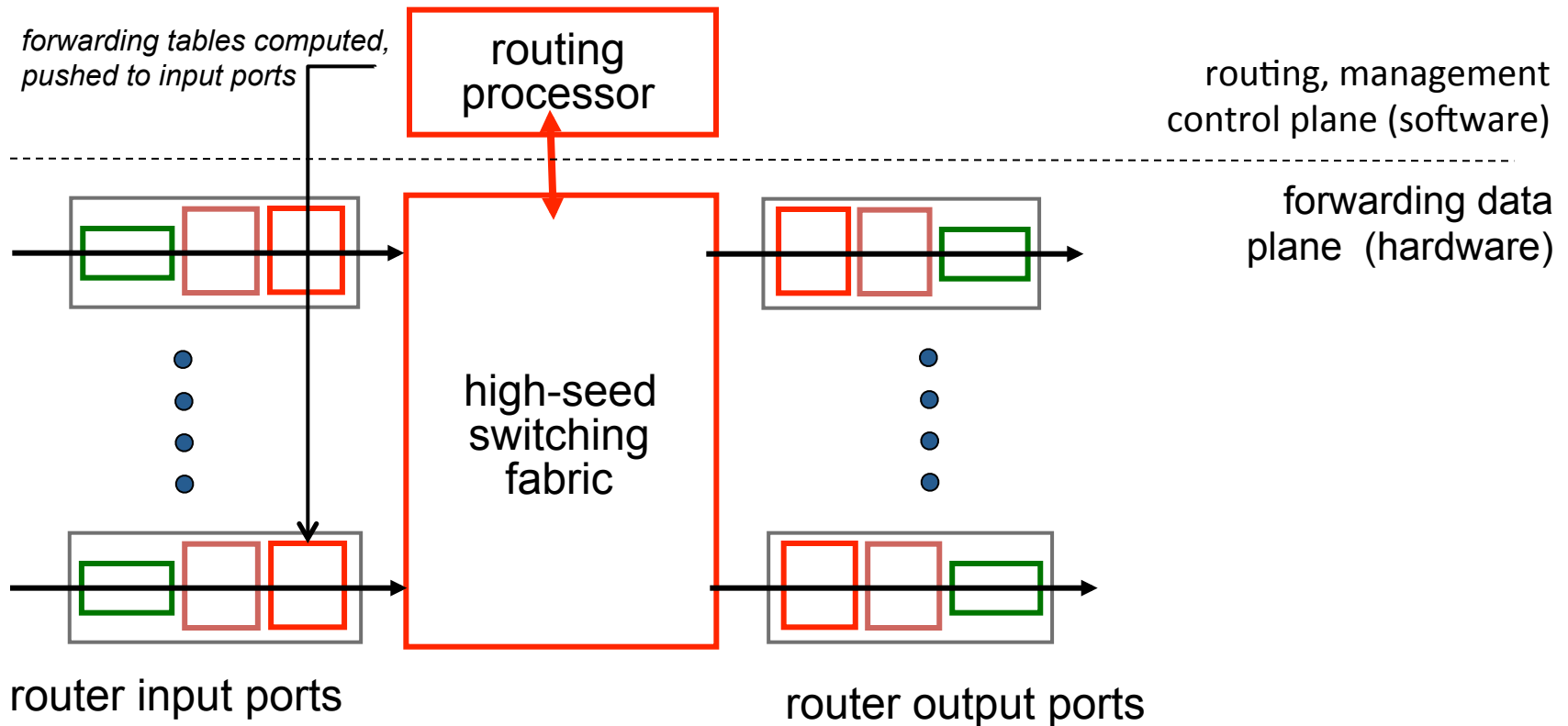
Datagram forwarding table review



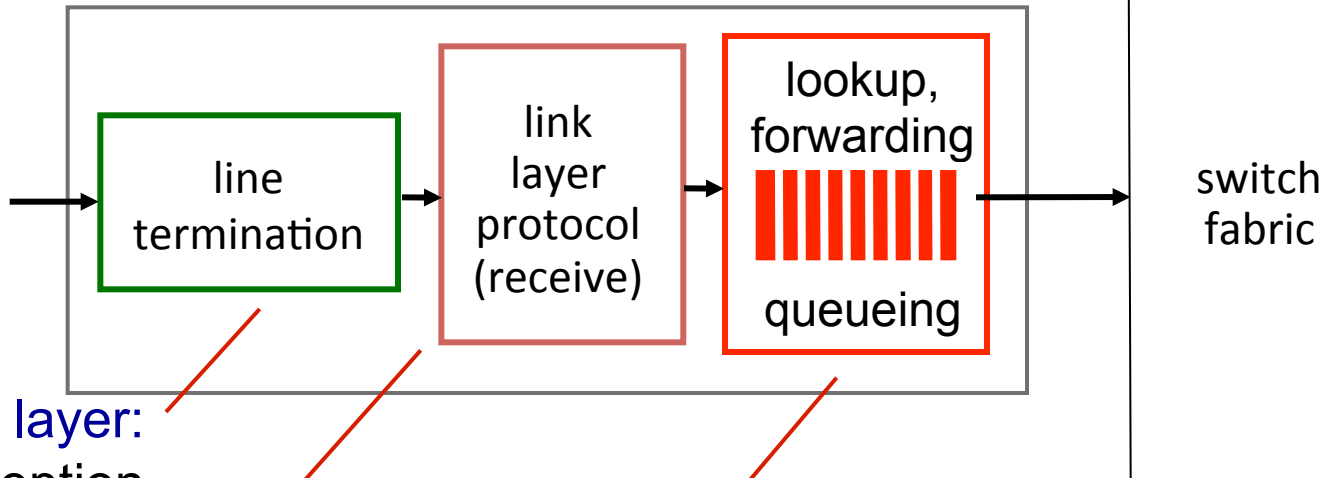
Router architecture overview

two key router functions:

- run routing algorithms/protocol (RIP, OSPF, BGP)
- *forwarding* datagrams from incoming to outgoing link



Input port functions



physical layer:
bit-level reception

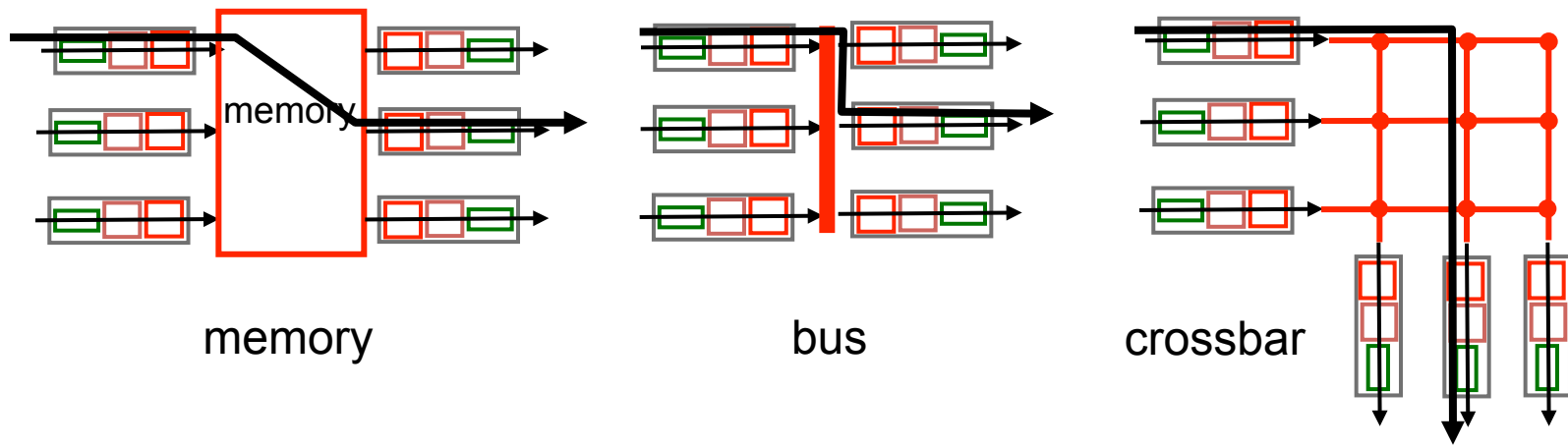
data link layer:
e.g., Ethernet
see chapter 5

decentralized switching:

- given datagram dest., lookup output port using forwarding table in input port memory (*“match plus action”*)
- goal: complete input port processing at ‘line speed’
- queuing: if datagrams arrive faster than forwarding rate into switch fabric

Switching fabrics

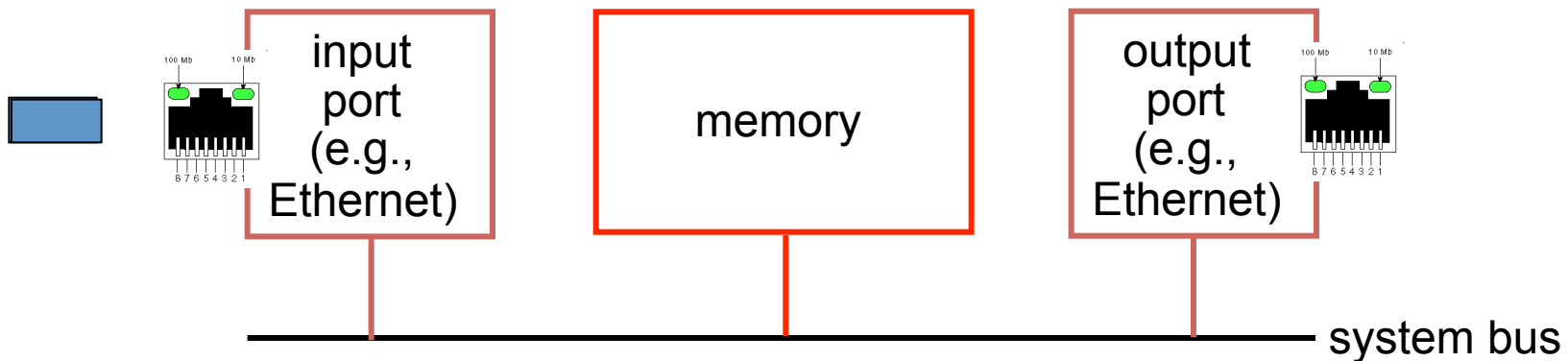
- transfer packet from input buffer to appropriate output buffer
- switching rate: rate at which packets can be transfer from inputs to outputs
 - often measured as multiple of input/output line rate
 - N inputs: switching rate N times line rate desirable
- three types of switching fabrics



Switching via memory

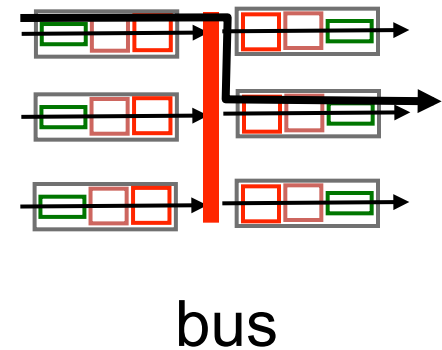
first generation routers:

- traditional computers with switching under direct control of CPU
- packet copied to system's memory
- speed limited by memory bandwidth (2 bus crossings per datagram)



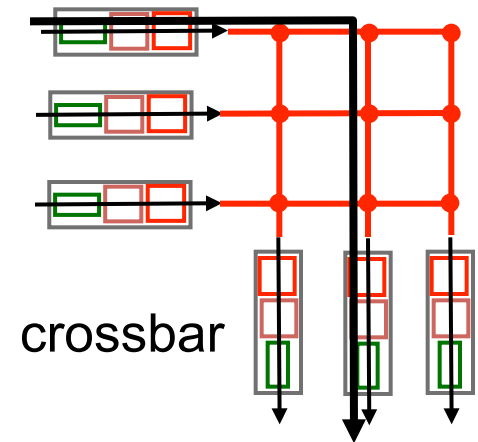
Switching via a bus

- datagram from input port memory to output port memory via a shared bus
- *bus contention*: switching speed limited by bus bandwidth
- 32 Gbps bus, Cisco 5600: sufficient speed for access and enterprise routers



Switching via interconnection network

- overcome bus bandwidth limitations
- banyan networks, crossbar, other interconnection nets initially developed to connect processors in multiprocessor
- advanced design: fragmenting datagram into fixed length cells, switch cells through the fabric.
- Cisco 12000: switches 60 Gbps through the interconnection network



Select a wrong statement

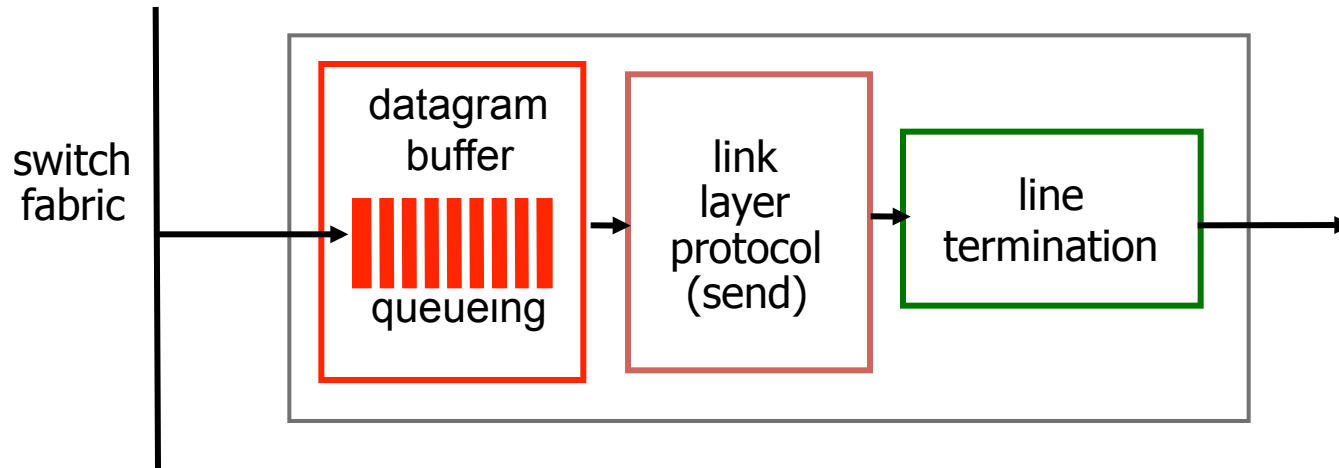
- A. Switching via memory handles one packet at the same time
- B. Switching via a bus is faster than switching via memory
- C. Switching via a bus can handle multiple packets at the same time
- D. Switching via interconnection network can handle multiple packets at the same time
- E. C and D

Switching via interconnection network

- A. Can handle multiple packets at the same time
- B. Faster than switching via memory or a bus
- C. Has no collision
- D. A and B
- E. A, B and C

This slide is HUGELY important!

Output ports

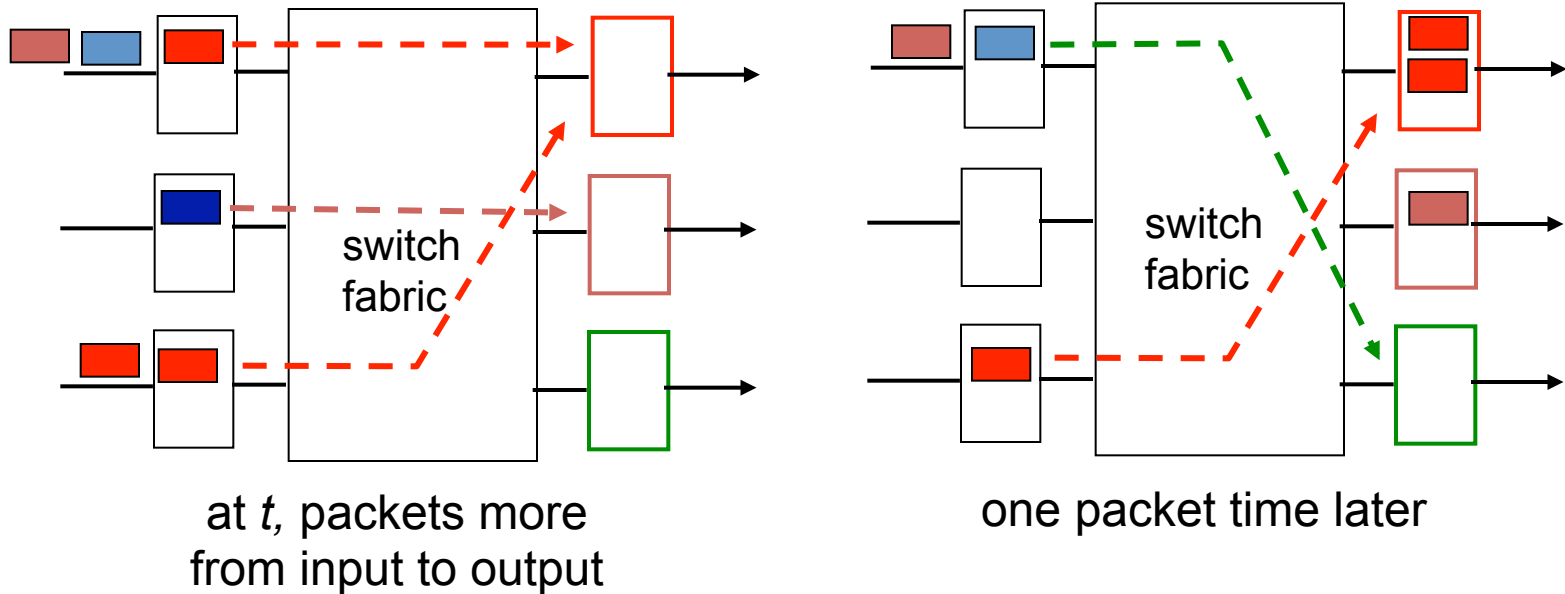


- *buffering* required when datagrams arrive from fabric faster than the transmission rate
- *scheduling discipline* chooses among queued datagrams for transmission

Datagram (packets) can be lost
due to congestion, lack of buffers

Priority scheduling – who gets best
performance, network neutrality

Output port queueing



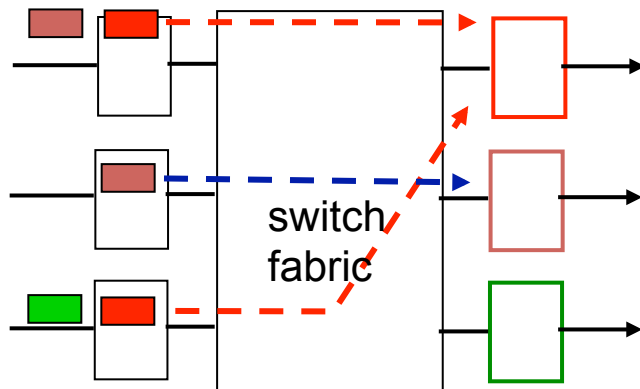
- buffering when arrival rate via switch exceeds output line speed
- *queueing (delay) and loss due to output port buffer overflow!*

How much buffering?

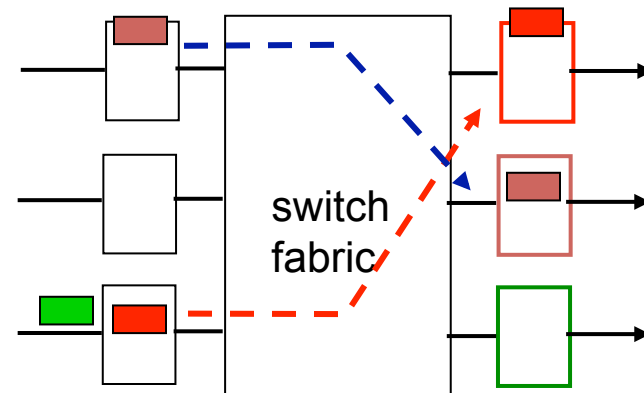
- RFC 3439 rule of thumb: average buffering equal to “typical” RTT (say 250 msec) times link capacity C
 - e.g., $C = 10$ Gpbs link: 2.5 Gbit buffer
- recent recommendation: with N flows, buffering equal to $\frac{RTT \cdot C}{\sqrt{N}}$

Input port queuing

- fabric slower than input ports combined -> queueing may occur at input queues
 - *queueing delay and loss due to input buffer overflow!*
- **Head-of-the-Line (HOL) blocking:** queued datagram at front of queue prevents others in queue from moving forward



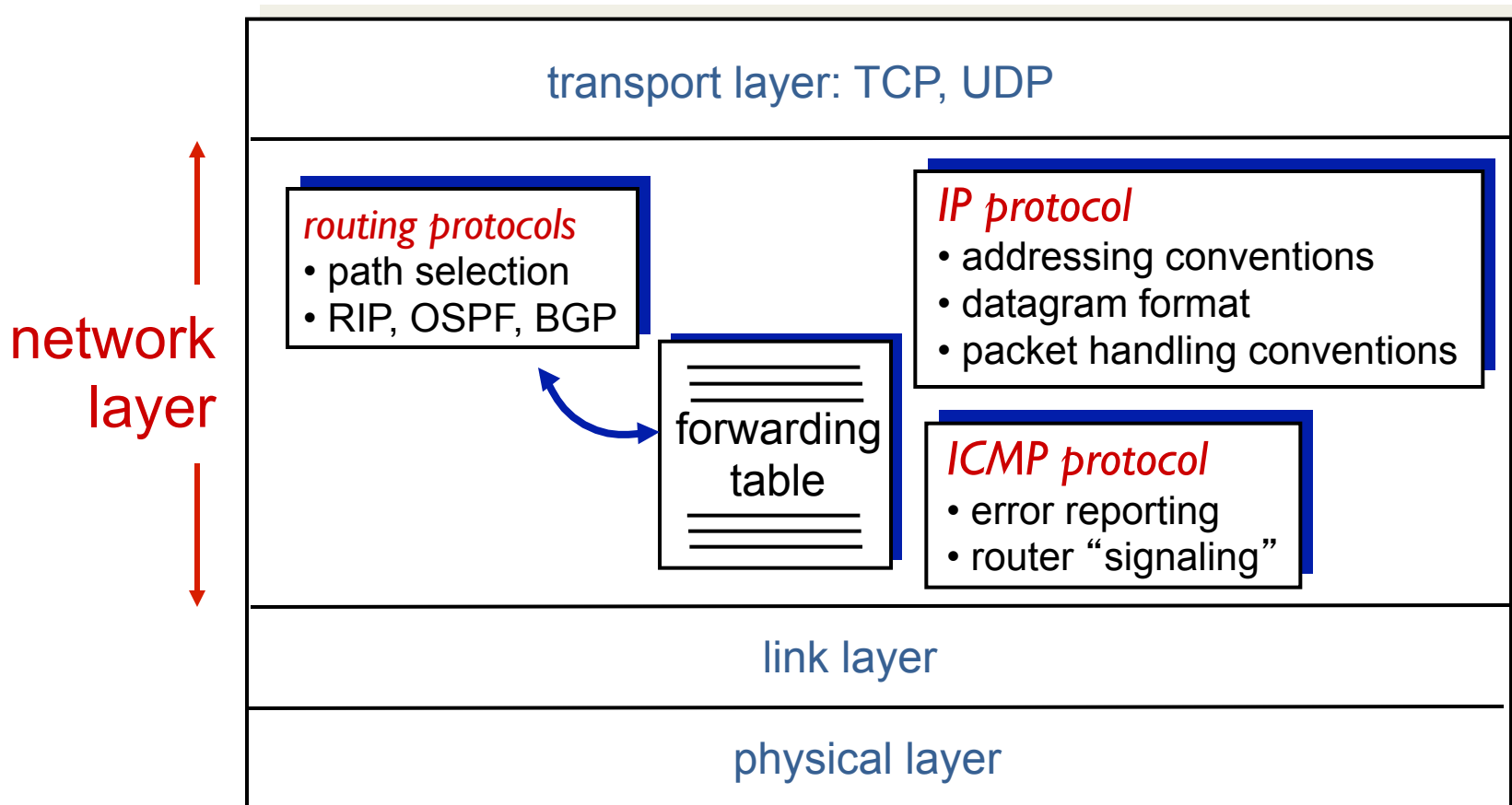
output port contention:
only one red datagram can be
transferred.
lower red packet is blocked



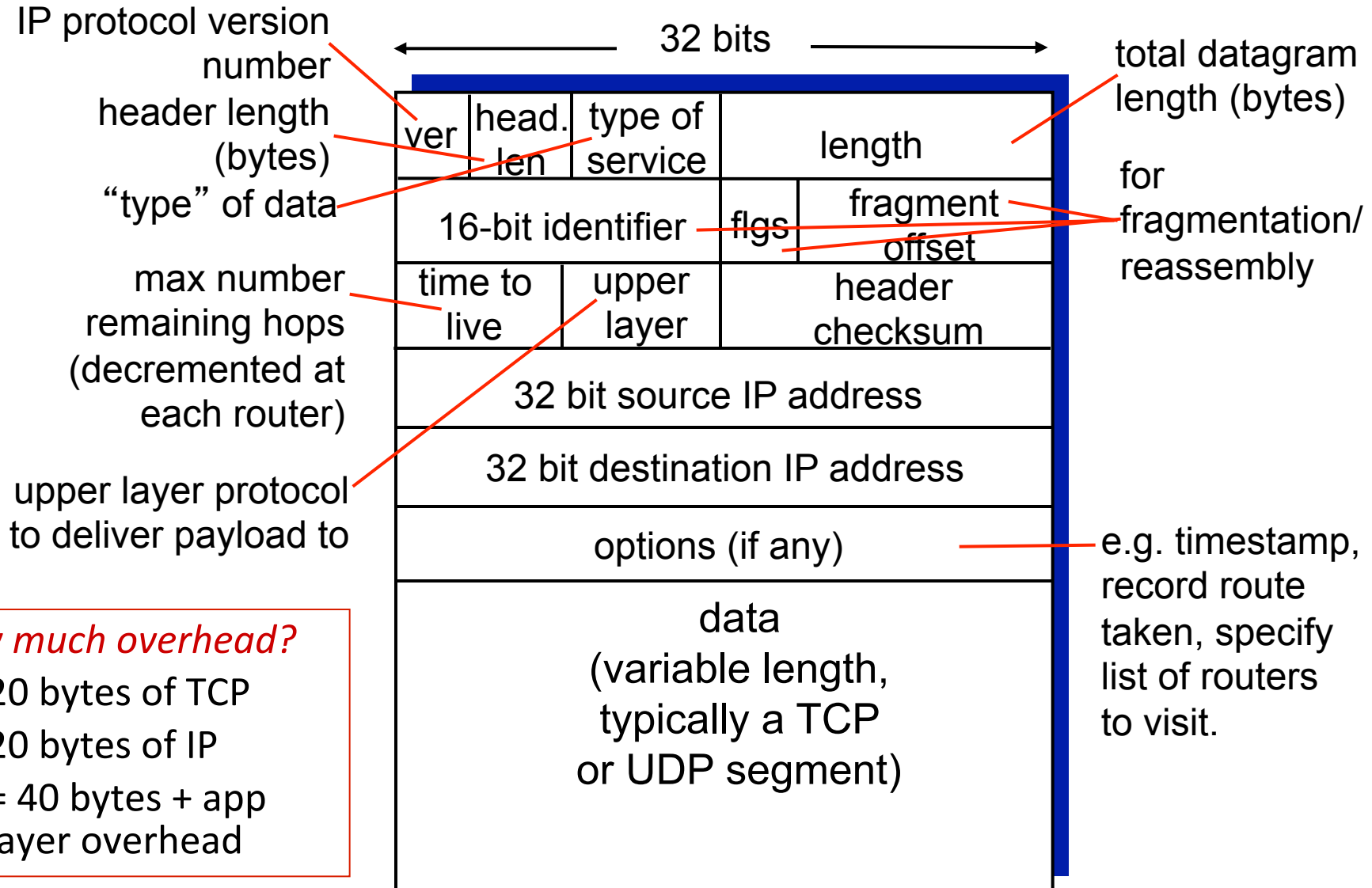
one packet time later:
green packet
experiences HOL
blocking

The Internet network layer

host, router network layer functions:

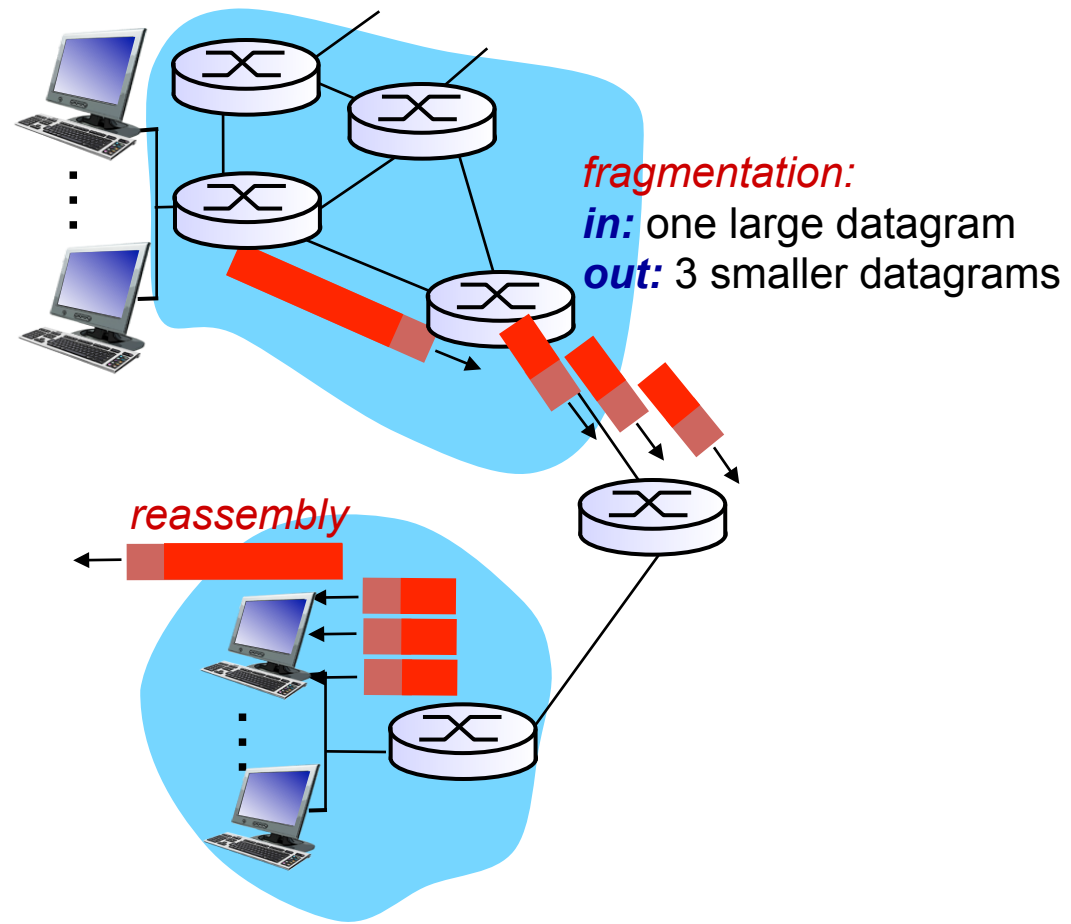


IP datagram format



IP fragmentation, reassembly

- network links have MTU (max transmission unit) - largest possible link-level frame
 - different link types, different MTUs
- large IP datagram divided (“fragmented”) within net
 - one datagram becomes several datagrams
 - “reassembled” only at final destination
 - IP header bits used to identify, order related fragments



IP fragmentation, reassembly

example:

- ❖ 4000 byte datagram
- ❖ MTU = 1500 bytes

	length =4000	ID =x	fragflag =0	offset =0	
--	-----------------	----------	----------------	--------------	--

*one large datagram becomes
several smaller datagrams*

1480 bytes in
data field

offset =
 $1480/8$

	length =1500	ID =x	fragflag =1	offset =0	
--	-----------------	----------	----------------	--------------	--

	length =1500	ID =x	fragflag =1	offset =185	
--	-----------------	----------	----------------	----------------	--

	length =1040	ID =x	fragflag =0	offset =370	
--	-----------------	----------	----------------	----------------	--

How easiness of the midterm?

- A. Very difficult
- B. Difficult
- C. It's okay
- D. Easy
- E. Very easy

Midterm: estimate how many points you will get?

- A. ≥ 25
- B. $21 \leq \text{your grade} < 25$
- C. $18 \leq \text{your grade} < 21$
- D. $14 \leq \text{your grade} < 18$
- E. < 14