

Lecture 2 : Solving linear systems of equations

Thursday, August 27, 2015 9:30 AM

Reading: Chapter 2 of Meyer

Outline : Linear equations
Geometry
Solution stability
Gaussian elimination
Experiments in Matlab/Mathematica
timing
sparse matrices
1D and 2D lattices
LU decomposition
Homogeneous and non-homogeneous systems

Linear equations

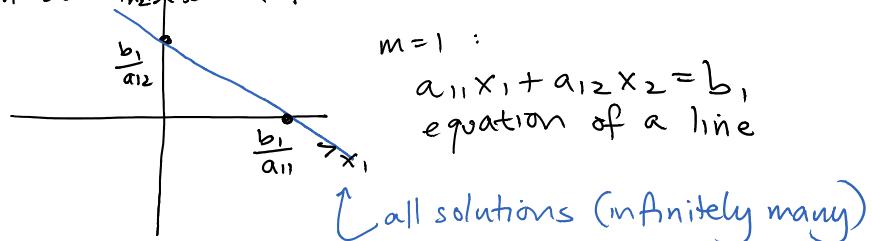
Recall the problem :

$$\underbrace{A \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix}}_{\substack{\text{m} \times n \text{ matrix} \\ n \text{ unknowns}}} = \underbrace{b}_{\substack{\text{m-dim vector}}} \iff \sum_{j=1}^n a_{ij} x_j = b_i \quad \text{for } i=1, \dots, n$$

Geometric picture

$n = 1$ one variable x_1 (trivial)

$n=2$ intersection of lines in \mathbb{R}^2



$m=2$: two lines

unique solution

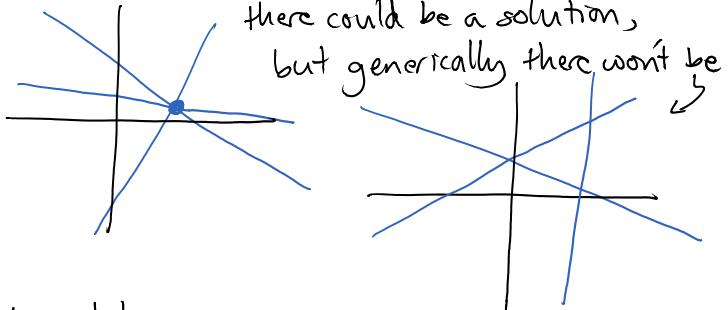
no solution
(parallel lines)

 this situation
is generic

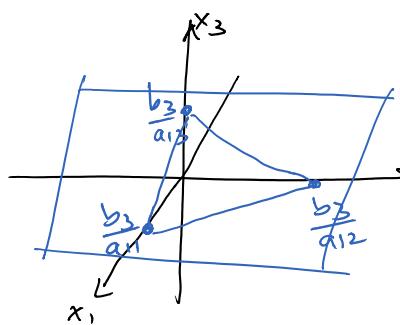
one line
∞ many
solutions

$m=3$: three lines

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$n=3$ variables:



$m=1$ equation

$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 = b_3$$

determines a 2D plane!

special case: $a_{11} = 0$

$\Rightarrow x_1$ doesn't matter

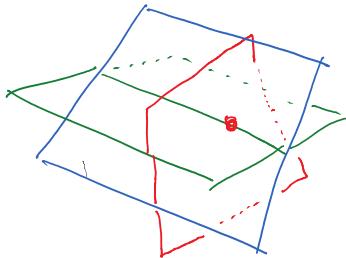
\Rightarrow plane is parallel to x_1 axis

$m=2$ equations:

generically, two planes intersect in a line

but they could also be parallel,
intersecting nowhere (or everywhere
if they are the same)

$m=3$ equations: a plane and a line generically
intersect in a single point



n variables x_1, x_2, \dots, x_n

equation \longleftrightarrow $(n-1)$ -dimensional hyperplane

$$\text{e.g., } \{(x_1, \dots, x_n) : x_n = 0\}$$

generically, two intersect in an $(n-2)$ -dim hyperplane,
and so on ...

n should intersect in a single point

(I'm not proving this, but just relying on your intuition)
(and from analogy to the $n=2, 3$ cases)

We'll say much more about this geometry later...

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Geometrically, though, what does Gaussian elimination correspond to?

- This is more easily seen when we introduce linear transformations: adding and multiplying rows leaves the range unchanged.

It is important!

① Stability — very important

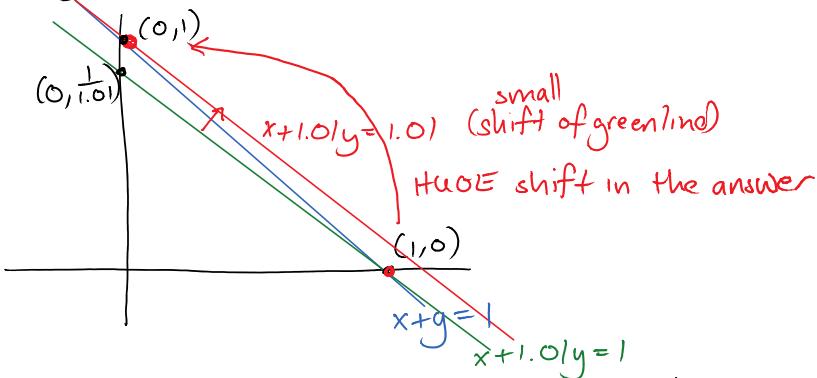
Example: Consider the equations

$$\begin{cases} x + y = 1 \\ x + 1.01y = 1 \end{cases} \Rightarrow (x, y) = (1, 0)$$

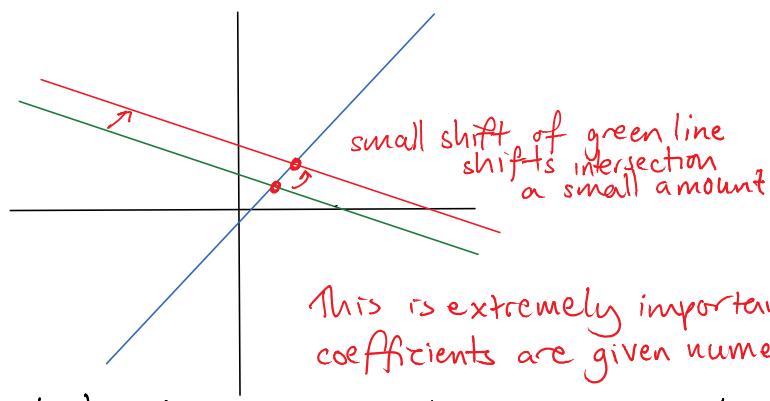
$$\begin{cases} x + y = 1 \\ x + 1.01y = 1.01 \end{cases} \Rightarrow (x, y) = (0, 1)$$

A small change to the equations led to a **HUGE** change in the solutions!

Why? Observe the geometry:



- because the lines are almost parallel
this wouldn't have happened if the lines were at a greater \angle



This is extremely important when coefficients are given numerically.

Note: In higher dimensions, you have the same problem with planes that are nearly parallel $\nearrow \searrow$

Note: In higher dimensions, you have the same problem with planes that are nearly parallel. But they don't have to be nearly parallel; it is more complicated than that. (Eg, think of above green and blue lines as intersections of planes in 3D with $\Sigma z = 0$ plane. They needn't be nearly parallel.)

We'll study this more later.

② Linear programming

Problem: Solve a system of linear inequalities.

$$a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n \leq b_1$$

$$a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n \leq b_2$$

⋮

* Very useful in practice — inequalities often capture resource constraints

Eg., say you have a \$100 budget for rice & beans

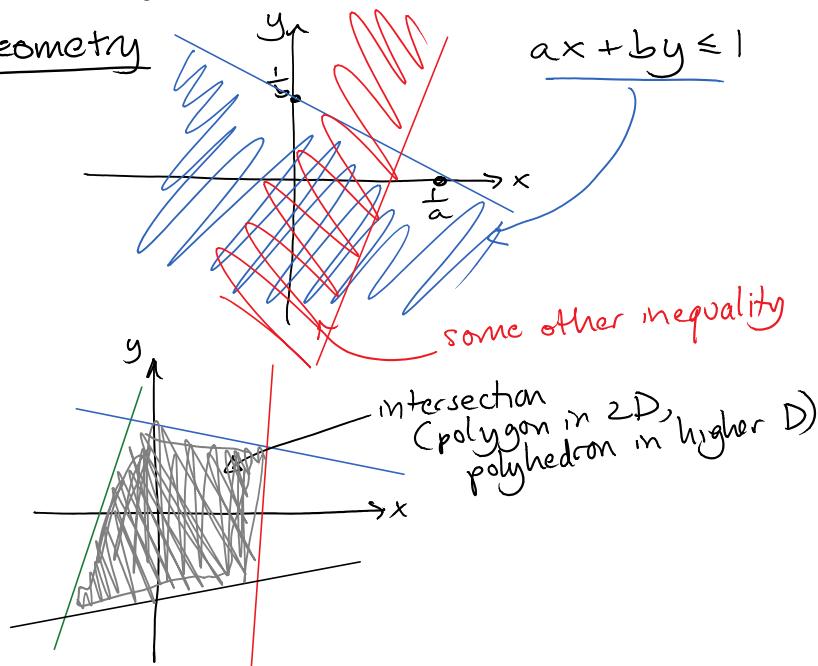
$$3r + 5b \leq 100$$

\$3/lb \$5/lb

* Generalize linear systems of equations

$$\sum_j a_{ij}x_j \leq b_i \quad \text{or} \quad \sum_j a_{ij}x_j \geq b_i \iff \sum_j a_{ij}x_j = b_i$$

Geometry



How To SOLVE SYSTEMS OF LINEAR EQUATIONS

- Gaussian elimination
 - LU decomposition
- Inverting the matrix A (if it is invertible)
- Iterative methods
 - preconditioners

- conjugate gradient (for a positive semi-definite matrix)
- multigrid methods ...

- Methods for solving SDD systems

Recall:

Gaussian elimination:

Observe : 1. Multiplying an equation by a nonzero number doesn't change the solution set.
2. With two equations, adding one to the other doesn't change the solution set.

$$(x \text{ solves } E_1, E_2) \iff x \text{ solves } E_1, E_1 + E_2$$

Therefore: Add multiples of one equation to the others to eliminate one variable from them.

Then repeat. Solve for the last variable, back substitute.
That's it!

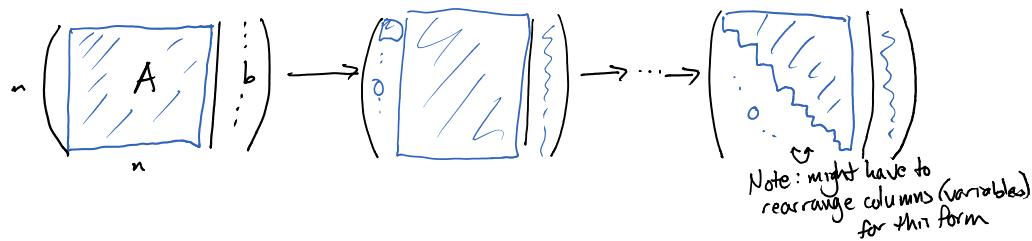
$O(n^3)$ complexity: n variables

for an $n \times n$ system add $O(n)$ pairs of equations to eliminate a variable

$O(n)$ coefficients to update in an equation.

∴ doubling $n \Rightarrow 4 \times$ the memory (to store A)

fix the running time!



Example: this works with any number of equations in

$$\left\{ \begin{array}{l} x_1 + 2x_2 + 2x_3 + 3x_4 = 0 \\ 2x_1 + 4x_2 + x_3 + 3x_4 = 0 \\ 3x_1 + 6x_2 + x_3 + 4x_4 = 0 \end{array} \right.$$

Note: $(0, 0, 0, 0)$ is one solution
—are there more?

$$\left(\begin{array}{cccc|c} 1 & 2 & 2 & 3 & 0 \\ 2 & 4 & 1 & 3 & 0 \\ 3 & 6 & 1 & 4 & 0 \end{array} \right) \xrightarrow{-2} \left(\begin{array}{cccc|c} 1 & 2 & 2 & 3 & 0 \\ 0 & 0 & -3 & -3 & 0 \\ 0 & 0 & -5 & -5 & 0 \end{array} \right) \xrightarrow{\frac{1}{3}} \left(\begin{array}{cccc|c} 1 & 2 & 2 & 3 & 0 \\ 0 & 0 & -3 & -3 & 0 \\ 0 & 0 & -5 & -5 & 0 \end{array} \right)$$

(Note : To minimize accumulation of floating point errors, it is often best to pivot on the largest magnitude entry in the column (s))

$$\left(\begin{array}{cccc|c} 1 & 2 & 2 & 3 & 0 \\ 0 & 0 & -3 & -3 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{array} \right) \xrightarrow{\text{R}_1 + \text{R}_2} \left(\begin{array}{cccc|c} 1 & 2 & 2 & 3 & 0 \\ 0 & 0 & -3 & -3 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{array} \right) \xrightarrow{O=0}$$

$$\xrightarrow{\quad} \left(\begin{array}{cccc|c} 1 & 2 & -1 & 0 & 0 \\ 0 & 0 & -3 & -3 & 0 \end{array} \right)$$

$$x_1 + 2x_2 - x_3 = x_3 + x_4 = 0$$

$$\Rightarrow x_3 = -x_4$$

$$x_1 = -2x_2 - x_4$$

$$\Rightarrow \text{solutions} = \{(x_1, x_2, x_3, x_4) \in \mathbb{R}^4 : \begin{cases} x_1 = -2x_2 - x_4 \\ x_3 = -x_4 \end{cases}\}$$

x_2 & x_4 are free — can be arbitrary

$$\left(\begin{array}{c} x_1 \\ x_2 \\ x_3 \\ x_4 \end{array} \right) = \left(\begin{array}{c} -2x_2 - x_4 \\ x_2 \\ -x_4 \\ x_4 \end{array} \right) = x_2 \left(\begin{array}{c} -2 \\ 1 \\ 0 \\ 0 \end{array} \right) + x_4 \left(\begin{array}{c} -1 \\ 0 \\ 1 \\ 1 \end{array} \right)$$

— the solution set is a 2D plane in \mathbb{R}^4

e.g. $\begin{array}{c|c} (x_2, x_4) & (x_1, x_2, x_3, x_4) \\ \hline (0,0) & (0, 0, 0, 0) \\ (1,0) & (-2, 1, 0, 0) \\ (0,1) & (-1, 0, -1, 1) \\ (1,1) & (-3, 1, 5, 1) \end{array}$

everything expressed in the free variables

Stepping back...

We started with a "homogeneous" system of equations (right-hand sides all 0):

$\Rightarrow \vec{0} = (0, 0, 0, 0)$ is a solution

\Rightarrow any multiple of a solution is a solution

$\vec{x} = (x_1, x_2, x_3, x_4)$ solution $\Rightarrow \lambda \vec{x} = (\lambda x_1, \lambda x_2, \lambda x_3, \lambda x_4)$ is, too

$$\text{b/c } \sum_{j=1}^n a_{ij}(\lambda x_j) = \lambda \cdot \sum_{j=1}^n a_{ij} x_j = \lambda \cdot 0 = 0 \quad -$$

\Rightarrow the sum of two solutions is a solution

\vec{x}, \vec{x}' solutions \Rightarrow so is $\vec{x} + \vec{x}'$

$$\text{b/c } \sum_{j=1}^n a_{ij}(x_j + x'_j) = \left(\sum_j a_{ij} x_j \right) + \left(\sum_j a_{ij} x'_j \right) = 0 + 0 = 0 \quad -$$

— this is the formal definition of a subspace, in this case

a plane parameterized by x_2 and x_4

Definition: A system of equations $Ax = b$ is

- homogeneous if $b = (0, 0, \dots, 0)$

- non-homogeneous if $b \neq 0$

Gaussian elimination works just as well for non-homogeneous systems...

Example:

$$x_1 + 2x_2 + 2x_3 + 3x_4 = 1$$

$$2x_1 + 4x_2 + x_3 + 3x_4 = -1$$

$$3x_1 + 6x_2 + x_3 + 4x_4 = 0 \quad -2$$

same coefficients as before

different $b \neq 0$

$$\left(\begin{array}{cccc|c} 1 & 2 & 2 & 3 & 1 \\ 2 & 4 & 1 & 3 & -1 \\ 3 & 6 & 1 & 4 & 0 \end{array} \right) \xrightarrow{-2} \left(\begin{array}{cccc|c} 1 & 2 & 2 & 3 & 1 \\ 0 & 0 & -3 & -3 & -3 \\ 0 & 0 & -5 & -5 & -3 \end{array} \right) \xrightarrow{-3} \left(\begin{array}{cccc|c} 1 & 2 & 2 & 3 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 \end{array} \right)$$

\nwarrow contradiction
 \Rightarrow no solution!

$$x_1 + 2x_2 - x_3 = -2$$

$$x_2 + x_4 = 1$$

\Rightarrow again, x_2 and x_4 are free

$$\begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{pmatrix} = \begin{pmatrix} -2 \\ -2x_2 + (1-x_4) \\ x_2 \\ 1-x_4 \end{pmatrix} = \underbrace{\begin{pmatrix} -1 \\ 0 \\ 0 \\ 0 \end{pmatrix}}_{\text{this part is the same as for the homogeneous system}} + x_2 \begin{pmatrix} -2 \\ 1 \\ 0 \\ 0 \end{pmatrix} + x_4 \begin{pmatrix} -1 \\ 0 \\ 0 \\ 1 \end{pmatrix}$$

(*) \Rightarrow The solution set for the nonhomogeneous equations is just a translation of the subspace (2D plane) of solutions to the corresponding homogeneous equations!

This makes sense: If $Ax = b$ $\overset{n-h \text{ sol.}}{n \text{ sol.}}$

$$Ay = 0 \Rightarrow A(x+y) = b + 0 = b \quad n-h \text{ sol.}$$

Example: Solving differential equations numerically

(§ 1.4)

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Problem :

Find a function f on $[0, 1]$ with

$$f''(t) = \sqrt{t}$$

$$f(0) = f(1) = 0$$

Answer: $f''(t) = \sqrt{t}$

$$\Rightarrow f'(t) = \frac{2}{3} t^{3/2} + C$$

$$\Rightarrow f(t) = \frac{4}{15} t^{5/2} + Ct + d$$

for some constants c, d

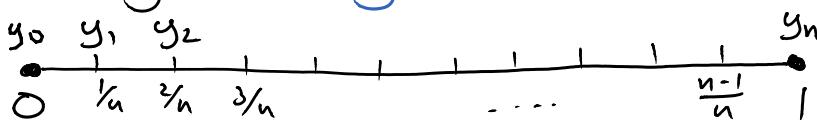
$$f(0) = 0 \Rightarrow d = 0$$

$$f(1) = \frac{4}{15} + c = 0 \Rightarrow c = -\frac{4}{15}$$

$$\therefore \boxed{f(t) = \frac{4}{15}(t^{5/2} - t)}$$

Now let's try solving this numerically.

by discretizing the interval $[0, 1]$



Goal: Solve for f only at these discrete points

variables $y_j = f(\frac{j}{n})$ for $j = 1, \dots, n-1$

$$f'(t) = \frac{d}{dt} f(t) \approx \frac{f(t+\frac{1}{n}) - f(t)}{\frac{1}{n}}$$

$f''(t)$ is the change in the derivative

$$\approx \frac{f'(t) - f'(t-\frac{1}{n})}{\frac{1}{n}} \approx \frac{1}{(\frac{1}{n})^2} ((f(t+\frac{1}{n}) - f(t)) - (f(t) - f(t-\frac{1}{n}))) \\ = n^2 (f(t+\frac{1}{n}) - 2f(t) + f(t-\frac{1}{n}))$$

equations

$$y_0 = 0$$

$$y_n = 0$$

$$y_{j-1} - 2y_j + y_{j+1} = \frac{1}{n^2} \sqrt{t}$$

Solving in Mathematica:

$n = 50$;

SparseArray

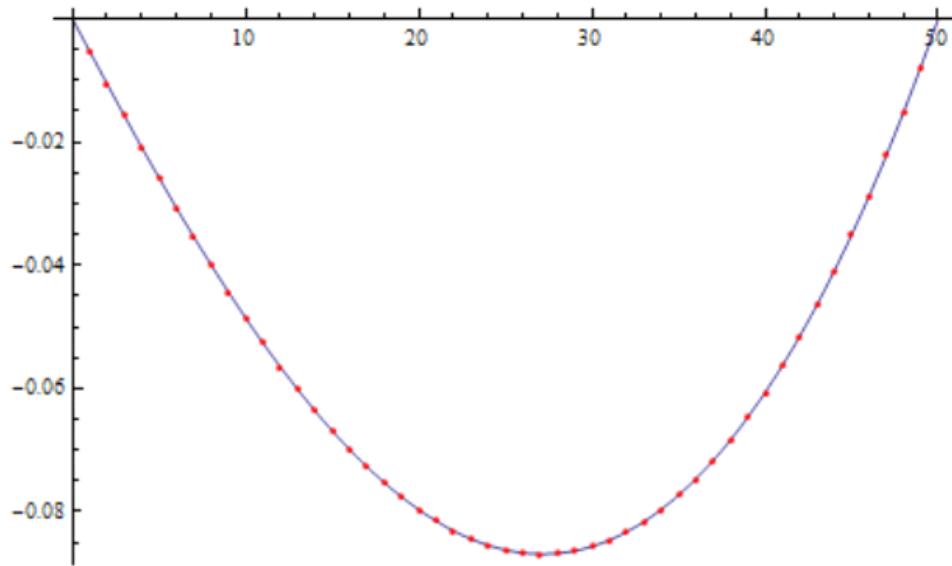
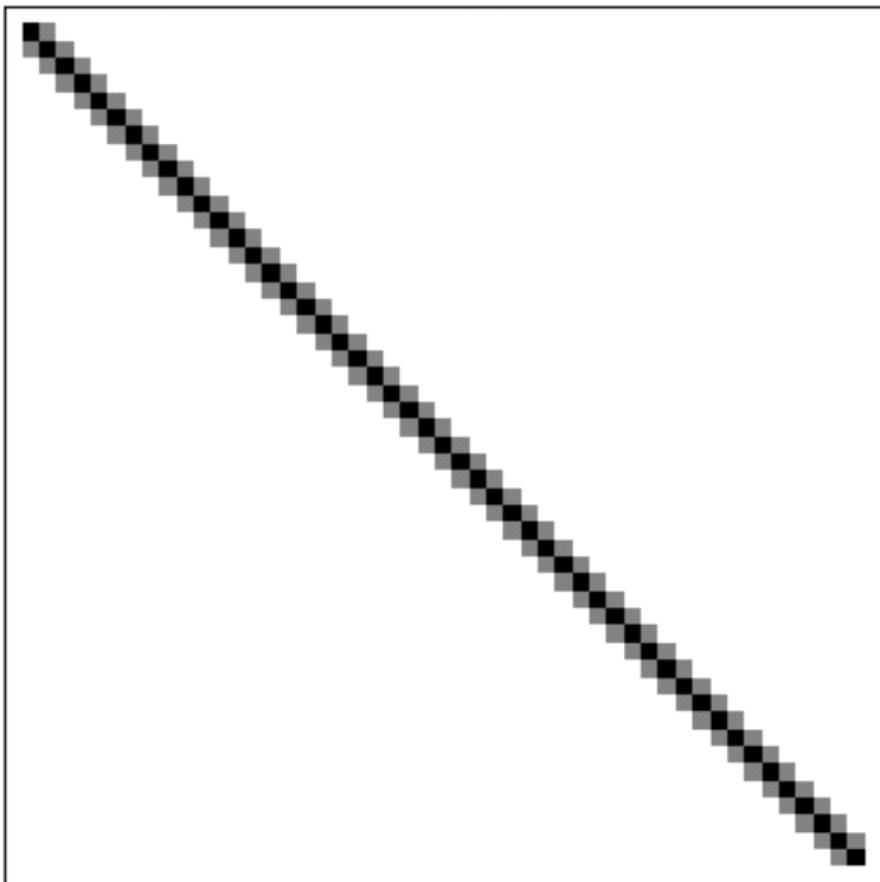
WVU ... + ... = 0

```
n = 50;  
A = SparseArray[Join[  
    Table[{j, j} → -2, {j, 1, n - 1}],  
    Table[{j, j - 1} → 1, {j, 2, n - 1}],  
    Table[{j, j + 1} → 1, {j, 1, n - 2}]  
]];  
  
b = Table[ $\frac{1}{n^2} \sqrt{\frac{j}{n}}$ , {j, 1, n - 1}] // N;  
  
Dimensions[A]  
Dimensions[b]  
  
ArrayPlot[A]  
  
plot1 = LinearSolve[A, b] // ListPlot[#, PlotStyle → RGBColor[1, 0, 0]] &;  
plot2 = Plot[ $\frac{4}{15} \left( \left(\frac{t}{n}\right)^{5/2} - \frac{t}{n} \right)$ , {t, 0, n}];  
Show[plot2, plot1]
```

Sparse Array
means that we store
only the nonzero
entries of the matrix!

compare the effects of changing n

Also, use `Normal[A]` to change it to a dense matrix
and compare the timing.



Extension: What if we had a differential equation in two variables, x and y ?
(on, e.g., the square $[0, 1] \times [0, 1]$)
What would the matrix A look like?