

Lecture 17: Stability of linear equations

Thursday, October 22, 2015 9:30 AM

Admin: Next homework out

Recall: Singular-value decomposition

$$A = \sum_i \lambda_i \vec{u}_i \vec{v}_i^T$$

SVD applications...

* Solving linear equations
 $Ax = b$

What is the **sensitivity**,
eg., to numerical errors?

Find the **shortest solution**

When there is no solution,
find x to **minimize** $\|Ax - b\|$

Least-squares regression analysis

* Rank minimization

Principal Component Analysis (PCA)

Data mining, clustering, recommendation systems, ...

NUMERICAL STABILITY OF $Ax = b$ systems of linear equations

Example:

$$\text{Let } A = \begin{pmatrix} 10^{-4} & 1 \\ 1 & 1 \end{pmatrix}, \quad b = \begin{pmatrix} b_1 \\ b_2 \end{pmatrix}$$

How sensitive is the solution $x = A^{-1}b$
to small errors (eg., numerical errors) in b ?

$$\text{Let } B = \begin{pmatrix} 1+10^{-4} & 1 \\ 1 & 1 \end{pmatrix}.$$

How sensitive is $B^{-1}b$ to small errors in b ?

Why not experiment?

Reading:
Meyer 1.6
5.12.1-2
Strang 7.2

```
>> A = [10^-4 1; 1 1];
B = [1+10^-4 1; 1 1];
b = [5;5];
berror = 10^(-5) * randn(2,1);
>> (A\b) - (A\b+berror)
```

```
ans =
    1.0304e-05
   -7.2699e-06
```

← pretty small solution not too sensitive — good!

```
>> (B\b) - (B\b+berror)
```

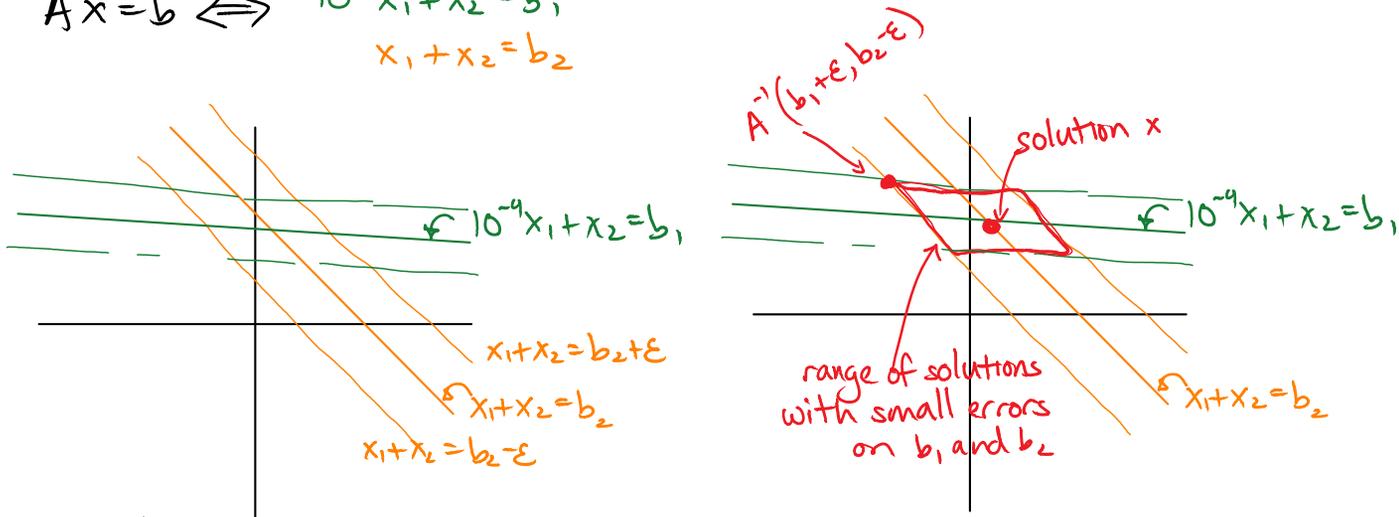
```
ans =
   -1.0303e-01
    1.0304e-01
```

← huge error!

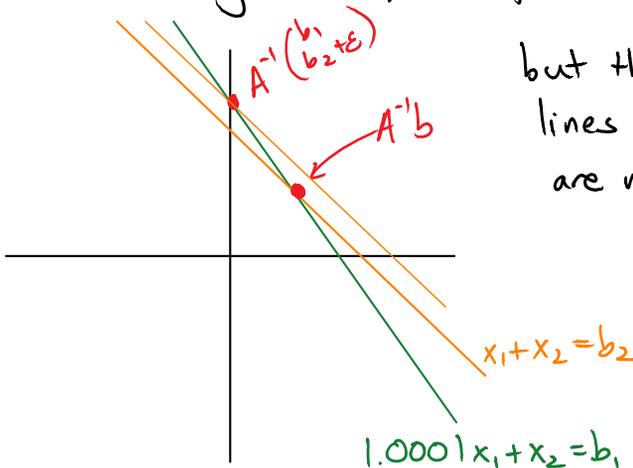
Why?

① Geometrically

$$Ax = b \iff \begin{cases} 10^{-4}x_1 + x_2 = b_1 \\ x_1 + x_2 = b_2 \end{cases}$$



small changes to b_1, b_2 shift the lines by small amounts



but the **intersection** of the lines can shift a lot, if they are nearly parallel

② Algebraic explanation for stability/instability

$$\|A^{-1}b - A^{-1}(b + \epsilon)\| = \|A^{-1}\epsilon\|$$

Algebraic explanation of sensitivity/stability

$$(*) \quad \|A^{-1}b - A^{-1}(b+\delta)\| = \|A^{-1}\delta\| \\ \leq \|A^{-1}\| \cdot \|\delta\|$$

\Rightarrow If $\|A^{-1}\|$ is large, then a small error in b can be amplified to a large error in $x = A^{-1}b$

Continuing the previous example:

```
octave:7> norm(A^-1)
ans = 1.6182
octave:8> norm(B^-1)
ans = 2.0001e+04
```

In the worst case, the upper bound in $(*)$ is reached!

$$A = \lambda_1 u_1 v_1^T + \lambda_2 u_2 v_2^T + \dots + \lambda_n u_n v_n^T$$

$$A^{-1} = \frac{1}{\lambda_1} v_1 u_1^T + \frac{1}{\lambda_2} v_2 u_2^T + \dots + \frac{1}{\lambda_n} v_n u_n^T$$

SVD, with $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n$

$$\Rightarrow \|A^{-1}\| = \lambda_n^{-1}$$

and if error vector $\vec{\delta} = \epsilon \vec{u}_n$,

$$\text{then } A^{-1}\vec{\delta} = \frac{\epsilon}{\lambda_n} v_n$$

$$\Rightarrow \|A^{-1}\delta\| = \|A^{-1}\| \cdot \|\delta\|$$

Easy solution?

Instead of solving $Ax = b$,

$$\text{solve } (1000A)y = b \Rightarrow y = \frac{1}{1000} A^{-1}b$$

and then set $x = 1000y$

$$\|(1000A)^{-1}\| = \frac{1}{1000} \|A^{-1}\| \Rightarrow \text{errors in } b \text{ now} \\ \text{lead to much smaller} \\ \text{errors in } y$$

\rightsquigarrow But this is cheating; the error in x will be the same

\Rightarrow Definition: The **condition number** of a matrix A is

$$\|A\| \cdot \|A^{-1}\|$$

equivalently, = largest singular value of A / smallest singular value of A

smallest singular value

Observe: This is invariant under scaling:
 $\|(1000A)\| \cdot \|(1000A)^{-1}\| = \|A\| \cdot \|A^{-1}\|$.

High CONDITION NUMBER $\longleftrightarrow Ax = b$ UNSTABLE
 $\frac{\max}{\min}$ sing. value

Example:

$$\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \leftarrow \text{condition \#} = 1 \text{ (as small as it can get)}$$
$$\begin{pmatrix} 10^{-4} & 0 \\ 0 & 10^{-4} \end{pmatrix}$$
$$\begin{pmatrix} 1 & 0 \\ 0 & 10^{-4} \end{pmatrix} \leftarrow \text{condition \#} = 10^4 \Rightarrow \text{unstable!}$$

What matters is not the scale of the matrix, but the range of scales in the matrix.

Formally: We have

$$\|A^{-1}f\| \leq \|A^{-1}\| \cdot \|f\|$$

$$\text{and } \|b\| = \|A \cdot A^{-1}b\| \leq \|A\| \cdot \|A^{-1}b\|$$

$$\Rightarrow \frac{\|A^{-1}f\|}{\|A^{-1}b\|} \leq \underbrace{\|A\| \cdot \|A^{-1}\|}_{\text{condition \# of } A} \cdot \frac{\|f\|}{\|b\|}$$

relative error in the solution ($\|A^{-1}f\|/\|x\|$) relative error in b

\Rightarrow condition number bounds the relative errors

Three sources of error in solving $Ax=b$

① Errors in the data \vec{b}

② Errors in the linear solve routine

Even if b is known exactly, errors such as

floating point errors might arise and accumulate while solving the equations.

Example:

$$B = \begin{pmatrix} 1+10^{-4} & 1 \\ 1 & 1 \end{pmatrix}, \quad b = \begin{pmatrix} 5 \\ 5 \end{pmatrix} \quad \text{as above}$$

Say you solve these and got $x = \begin{pmatrix} .5 \\ 4.5 \end{pmatrix}$

```
>> B * [.5;4.5]
```

```
ans =
```

```
5.0000e+00  
5.0000e+00
```

← Looks good!

But is it?

No! The right answer is $\begin{pmatrix} 0 \\ 5 \end{pmatrix}$!

```
>> format long e
```

```
B * [.5;4.5]
```

```
ans =
```

```
5.0000500000000000e+00  
5.0000000000000000e+00
```

when you check the answer,
it looks great
— but it is very far from
the correct answer!

This is actually the same problem we already considered. Even though the error in b is small, the corresponding error in x is huge. This is only a good check for well-conditioned systems.

Moral: With an ill-conditioned system of linear equations

- use high-precision arithmetic
- find another way to check the answer?
- be careful! possibly try to reformulate the problem
→ use a preconditioner

Solve $P^{-1}Ax = P^{-1}b$, where P is such that $P^{-1}A$ is better conditioned, but P^{-1} can be applied efficiently.

This is an important, but more advanced, topic, that we might get to later.

③ Errors in the matrix A

Comparing

$$Ax = b \quad \text{and} \quad (A + \delta A)(x + \delta x) = b$$

some error in the matrix

$$\Rightarrow A \cdot \delta x + (\delta A) \cdot (x + \delta x) = 0$$

the resulting error in the solution

$$\Rightarrow \delta x = -A^{-1} \cdot (\delta A) \cdot (x + \delta x)$$

$$\frac{\|\delta x\|}{\|x + \delta x\|} \leq \|A^{-1}\| \cdot \|A\| \cdot \frac{\|\delta A\|}{\|A\|}$$

relative error in the solution condition # of A relative error in A

Example:

```
octave:11> Aerror = 10^(-5) * randn(2,2);  
octave:12> (A\b) - ((A+Aerror)\b)  
ans =
```

```
-1.1004e-04  
7.6036e-05
```

```
octave:13> (B\b) - ((B+Aerror)\b)  
ans =
```

```
-210.22  
210.26
```

Moral: In all three cases, the (relative) effect of the error is bounded by the condition number

$$= \|A\| \cdot \|A^{-1}\| = \frac{\text{largest singular value}}{\text{smallest singular value}}$$

Note: In practice, computing the condition number can be a lot of work.

It is enough usually to approximate it, or give upper bounds.

Certain classes of matrices reliably have large condition numbers; with experience you learn when to be careful.

Example: What is the condition number of

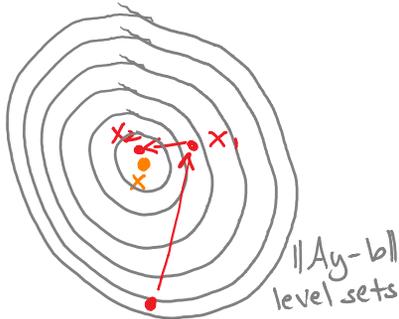
$$\begin{pmatrix} -2 & 1 \\ 1 & 1 \end{pmatrix}$$

POORLY CONDITIONED SYSTEMS OF EQUATIONS TAKE LONGER TO SOLVE

Rough intuition: Iterative solvers for $Ax=b$ start with a guess, and try to improve it.

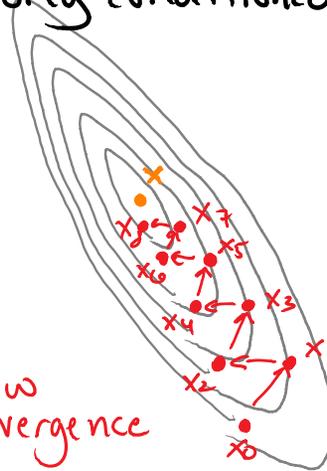
$$x_0 \rightarrow x_1 \rightarrow x_2 \rightarrow x_3 \rightarrow \dots$$

Well-conditioned A



Fast convergence

Poorly conditioned A



Slow convergence

Theorem:

Conjugate-gradient method on a positive semi-definite matrix A , with condition number κ , takes

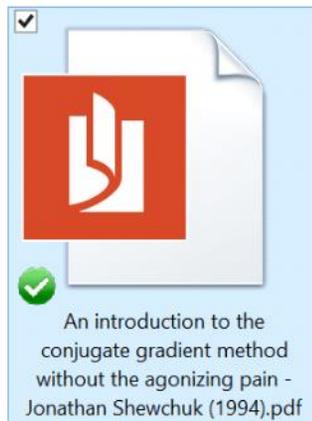
$$O(\sqrt{\kappa} \cdot \log \frac{1}{\epsilon})$$

steps to get within ϵ of the true solution x .
(Or $O(\kappa \log \frac{1}{\epsilon})$ steps for a general matrix A .)

Each step uses a matrix-vector multiplication
= $O(\# \text{ of nonzero entries in } A)$ steps.

Time permitting, we'll discuss this later.

Or, see \rightarrow



Example: GRADIENT DESCENT

Goal: Solve $A\vec{x} = \vec{b}$.

Define the cost function

$$\begin{aligned} C(\vec{x}) &= \|Ax - b\|^2 \\ &= \|A(x - x^*)\|^2 \quad \text{where } Ax^* = b \\ &= (x - x^*)^T \underbrace{A^T A}_{\text{positive semi-definite}} (x - x^*) \end{aligned}$$

If A has SVD

$$\begin{aligned} A &= \sum_i \lambda_i u_i v_i^T \\ A^T A &= \sum_i \lambda_i^2 v_i v_i^T \end{aligned}$$

$$\Rightarrow C(\vec{x}) = \sum_i \lambda_i^2 |v_i \cdot (x - x^*)|^2$$

\Rightarrow level set $\{ \vec{x} \mid C(\vec{x}) = b \}$ is an ellipse!
(like $\lambda_1^2 x_1^2 + \lambda_2^2 x_2^2 + \dots + \lambda_n^2 x_n^2 = b$)

Algorithm: Repeat until convergence:

$x^t = x^{t-1} - \alpha \left(\frac{\partial C}{\partial x_1}(x^{t-1}), \dots, \frac{\partial C}{\partial x_n}(x^{t-1}) \right)$
i.e., step in the direction opposite the largest increase in C .

$$\begin{aligned} \frac{\partial C}{\partial x_j} &= \frac{\partial}{\partial x_j} [(Ax - b)^T (Ax - b)] \\ &= \frac{\partial}{\partial x_j} \sum_i (Ax - b)_i^2 \\ &= 2 \sum_i (Ax - b)_i \cdot \frac{\partial}{\partial x_j} [(Ax - b)_i] \\ &= 2 \sum_i (Ax - b)_i \cdot a_{ij} \\ \Rightarrow \left(\frac{\partial C}{\partial x_1}, \dots, \frac{\partial C}{\partial x_n} \right) &= 2 \sum_i \underbrace{(Ax - b)_i}_{(Ax - b)^T e_i} \cdot \underbrace{(a_{i1}, \dots, a_{in})}_{e_i^T A} \\ &= 2(Ax - b)^T \left(\sum_i e_i e_i^T \right) A \\ &= 2(Ax - b)^T A \end{aligned}$$

(As you might guess!)

Example:

$$A = \begin{pmatrix} 2 & -1.5 \\ -1.5 & 4 \end{pmatrix};$$

$$b = A \cdot \{3, 3\};$$

Eigenvalues[A]

$$\alpha = .01;$$

$$x = \{0, -4\};$$

$$\text{list} = \{x\};$$

For[k = 1, k ≤ 1000, k++,

$$x = x - 2\alpha (A \cdot x - b) \cdot A;$$

AppendTo[list, x];

];

"The solution is:"

x

$$\{4.80278, 1.19722\}$$

The solution is:

$$\{3., 3.\}$$

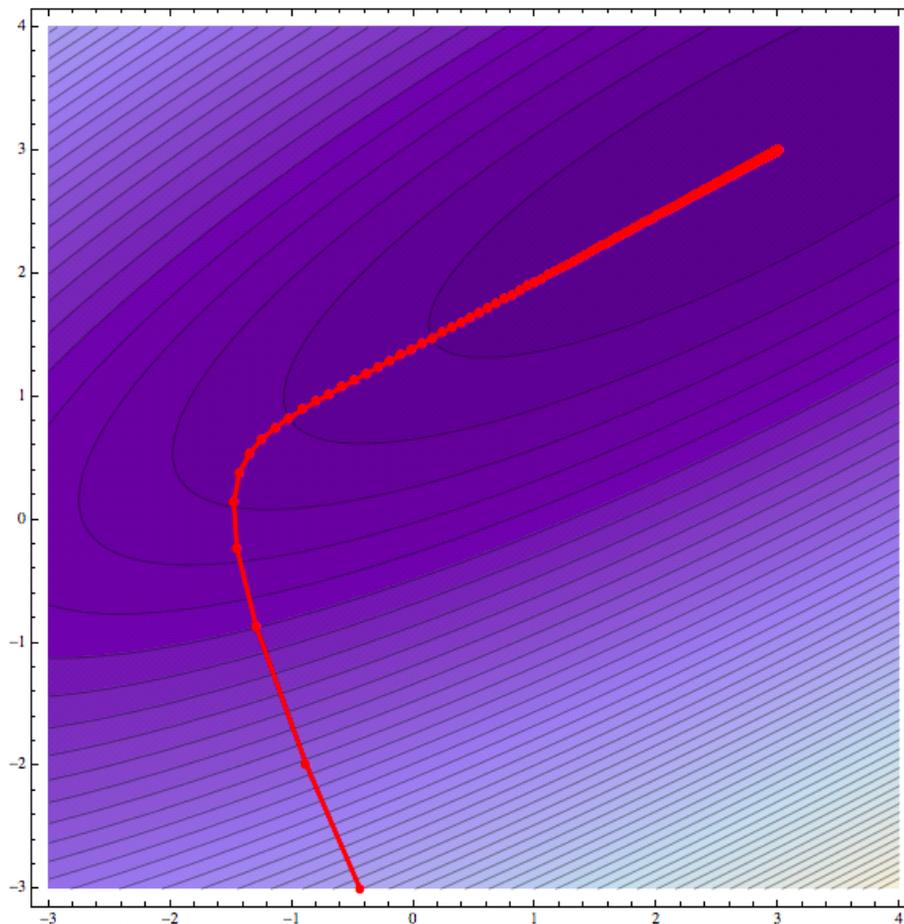
"Display the results:";

iter = ListPlot[list, PlotRange → 3 {{-1, 1}, {-1, 1}}, Joined → True,

PlotStyle → {Red, Thickness[.005]}, PlotMarkers → Automatic];

contour = ContourPlot[Norm[A · {x, y} - b]^2, {x, -3, 4}, {y, -3, 4}, Contours → 50];

Show[contour, iter]



For small enough α , this has to converge, since there is a unique local and global minimum.

Observe: Smaller condition number

⇒ Rounder ellipse ⇒ Faster convergence

Remark: Although not generally the fastest way of solving a set of linear equations, the gradient descent approach is very robust, and many variants are used in many applications.

Caution! Generalized inverses are useful in formulating theoretical statements such as those above, but, just as in the case of the ordinary inverse, generalized inverses are not practical computational tools. In addition to being computationally inefficient, serious numerical problems result from the fact that A^\dagger need

* The analysis is the same for **errors in your linear solve procedure**. Maybe A and b are exact, but your solver returned $\hat{x} = x + \delta x$ instead of x . To check the answer, multiply by A , and see how far the result is from b .

$$A\hat{x} = Ax + A(\delta x) = b + \delta b$$

$$\Rightarrow \delta x = -A^{-1} \cdot (\delta b)$$

the same as before.

$$\Rightarrow \frac{\|\delta x\|}{\|x\|} \leq \|A\| \cdot \|A^{-1}\| \cdot \frac{\|\delta b\|}{\|b\|}$$

\Rightarrow If the condition number is small, then computing $\|\delta b\|/\|b\|$ gives a good estimate for the relative error in your solution $\|\hat{x} - x\|/\|x\|$.

But not otherwise!