

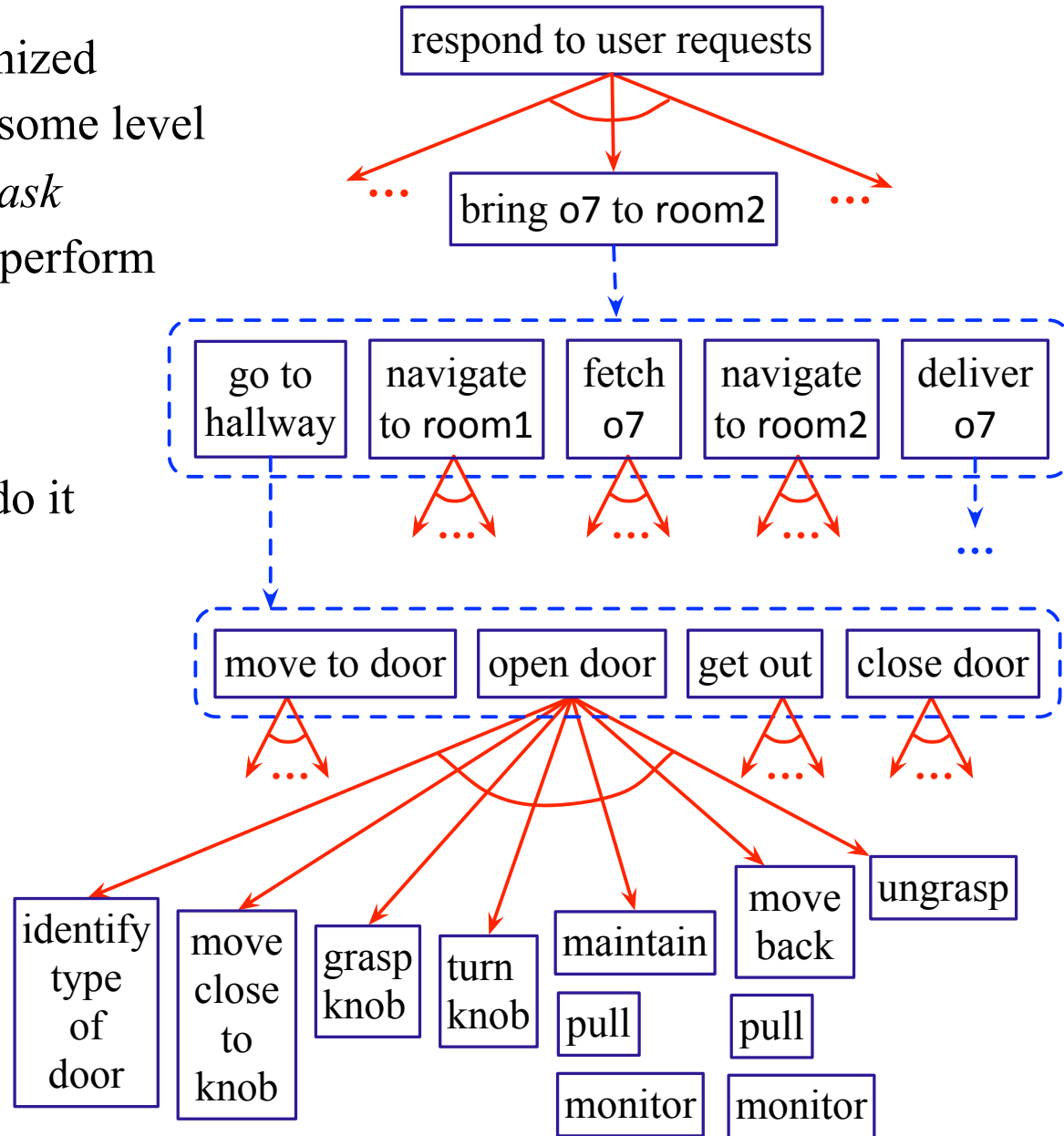
Chapter 3

Deliberation with Refinement Methods

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Motivation

- Actor is hierarchically organized
 - Consider an action a at some level
 - One level down, a is a *task*
 - complex activity to perform
- Recursively refine tasks into subtasks
 - Dashed blue lines: state-space planner can do it
 - Solid red lines: need other techniques
- Refinement method:* operational model
 - Program to compute a refinement
 - Actor can use it online
 - Can also use it for planning



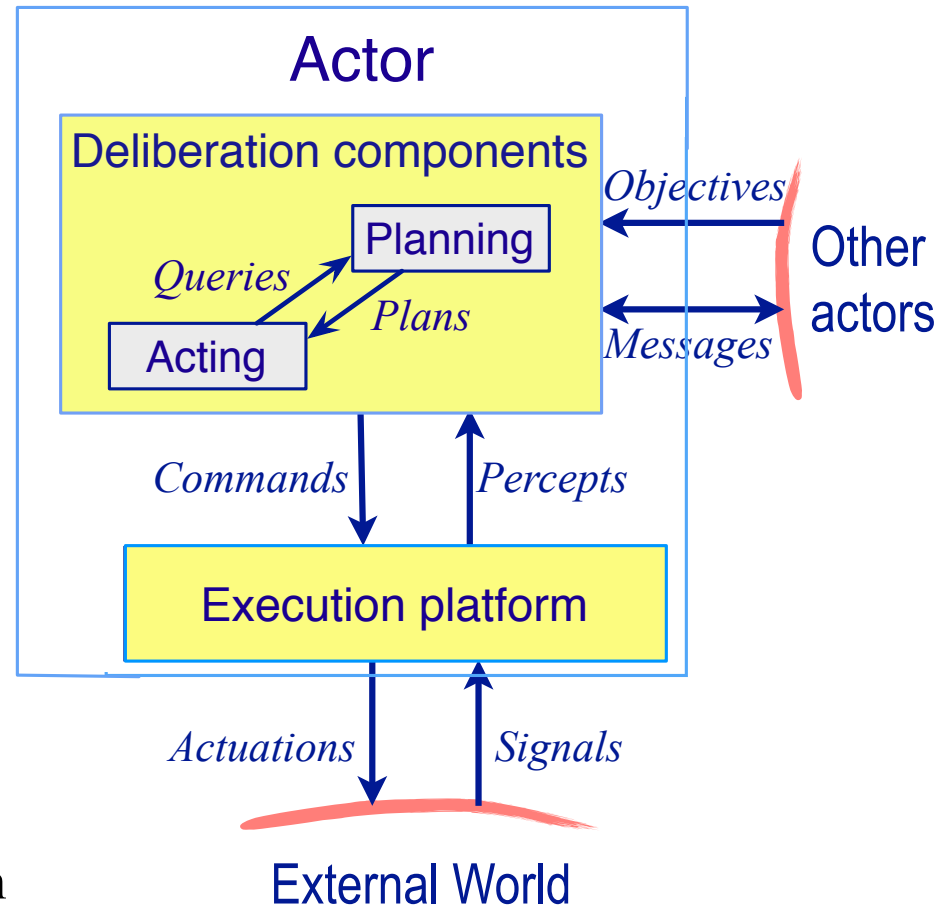
Outline

- 3.1 Operational models
- 3.2 Refinement Acting Engine (RAE)
- 3.3 Refinement planning (SeRPE)
- 3.4 Acting and refinement planning

3.1 Operational Models

Notation and terminology:

- ξ is the actor's current state
- Ξ = set of all possible ξ
- Execution platform keeps ξ up-to-date
 - may differ from the state s the actor gives the planner
- *Fact*: ground atom that's true in ξ
- *Task*: activity that the actor needs to carry out
- *Command*: primitive function that the execution platform can perform
- *Event*: occurrence detected by the execution platform
 - exogenous change in the environment



State-Variable Representation

Extensions:

- $\text{Range}(x)$
 - can be finite, infinite, continuous, discontinuous, vectors, matrices, other data structures
- Assignment statement $x \leftarrow \text{expr}$
 - expression that returns a ground value in $\text{Range}(x)$ and has no side-effects on the current state
- Tests (e.g., preconditions)
 - *Simple*: $x = v$, $x \neq v$, $x > v$, $x < v$
 - *Compound*: conjunction, disjunction, or negation of simple tests

Example

- Objects
 - $Robots = \{rbt\}$, $Containers = \{c1, c2, \dots\}$, $Locations = \{loc1, loc2, \dots\}$
- State variables
 - $loc(r) \in Locations$
 - $load(r) \in Containers \cup \{nil\}$
 - $pos(c) \in Locations \cup Robots \cup \{unknown\}$
 - $view(l) \in \{T, F\}$
 - Whether the robot has looked at location l or not
 - If $view(l) = T$ then for every container c at location l , $pos(c) = l$
- Commands to the execution platform:
 - $take(r, o, l)$: r takes object o at location l
 - $put(r, o, l)$: r puts o at location l
 - $perceive(r, l)$: robot r perceives what objects are at location l

Refinement Methods

- Refinement method: a triple (*task or event*, *precondition*, *body*)
 - a task or event that the method can be used for
 - precondition that needs to be true for the method to be applicable
 - body: a program
- Write in the following format:

method-name($arg_1, \dots, arg_k$)

task: *task-identifier*

pre: *test*

body: *program*

method-name($arg_1, \dots, arg_k$)

event: *event-identifier*

pre: *test*

body: *program*

- *body*: sequence of steps
 - assignments, commands, tasks
 - control constructs: if-then-else, while, loop, etc.

Methods for Tasks

- Example task: fetch container c

➤ If you know c 's location, go get it; else search

$m\text{-fetch}(r, c)$

task: $\text{fetch}(r, c)$

pre:

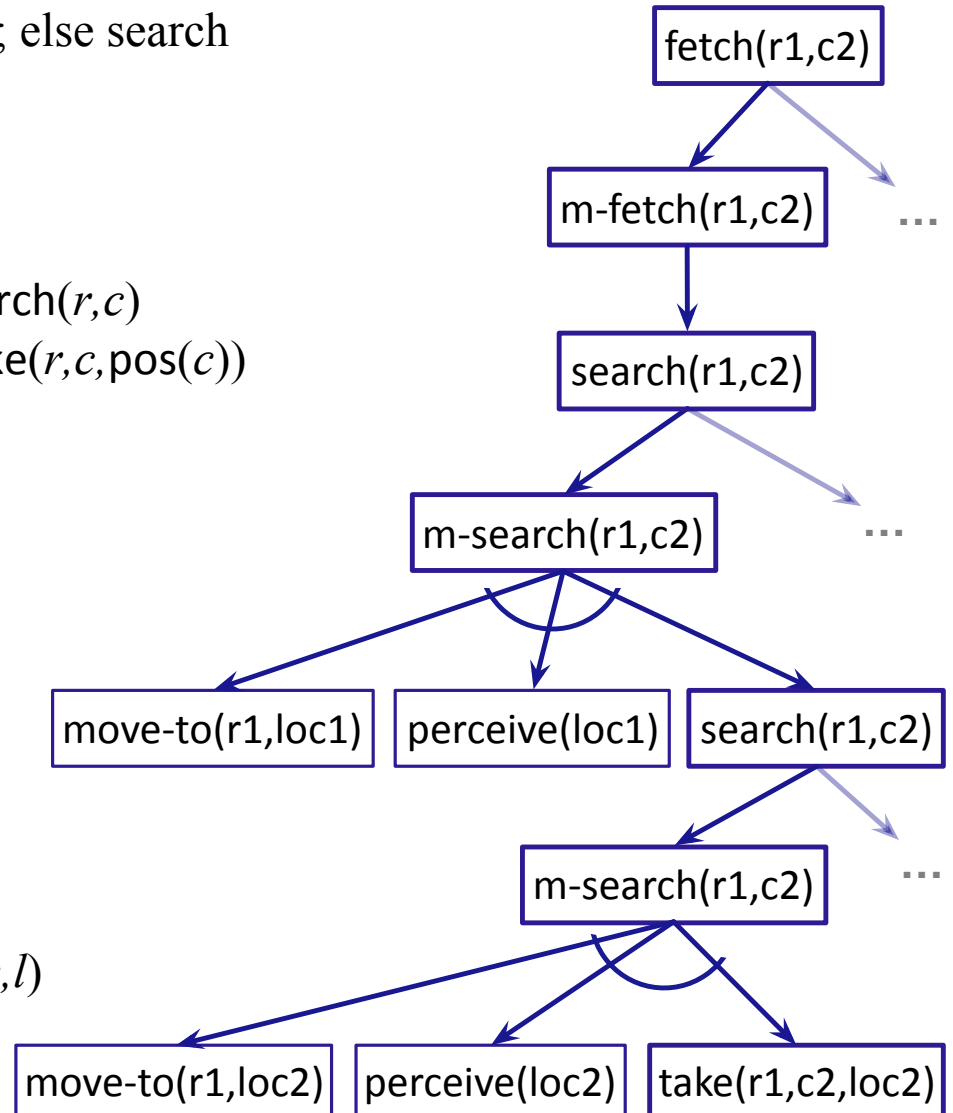
body: if $\text{pos}(c) = \text{unknown}$ then $\text{search}(r, c)$
 else if $\text{loc}(r) = \text{pos}(c)$ then $\text{take}(r, c, \text{pos}(c))$
 else do
 $\text{move-to}(r, \text{pos}(c))$
 $\text{take}(r, c, \text{pos}(c))$

$m\text{-search}(r, c)$

task: $\text{search}(r, c)$

pre: $\text{pos}(c) = \text{unknown}$

body: if $\exists l (\text{view}(l) = \text{F})$ then
 $\text{move-to}(r, l)$
 $\text{perceive}(l)$
 if $\text{pos}(c) = l$ then $\text{take}(r, c, l)$
 else $\text{search}(r, c)$
 else fail



Methods for Events

- Example event: emergency
 - If you aren't already handling another emergency, then
 - stop what you're doing, go to where emergency is, do something about it
- Another state variable
 - $\text{emergency-handling}(r) \in \{T, F\}$
 - whether robot r is handling an emergency

m-emergency(r, l, i)

i is the ID number of the event

event: emergency(l, i)

pre: emergency-handling(r) = F

body: emergency-handling(r) $\leftarrow$ T
if load(r) $\neq$ nil then put(r , load(r))
move-to(l)
address-emergency(l, i)

Example: Opening a Door

- State variables:
 - $\text{reachable}(r,o) \in \{\text{T}, \text{F}\}$
 - is o close enough to grasp?
 - $\text{door-status}(d) \in \{\text{closed}, \text{cracked}, \text{open}, \text{unknown}\}$
- Rigid relations:
 - $\text{adjacent}(l,d)$ location l is adjacent to door d
 - $\text{toward-side}(l,d)$ location l is on the side where d 's door hinges are
 - $\text{away-side}(l,d)$ location l is on the opposite side of d 's door hinges
 - $\text{handle}(d,o)$ o is the handle of door d
 - $\text{type}(d,\text{rotates})$ door d rotates
 - $\text{type}(d,\text{slide})$ door d slides
 - $\text{side}(d,\text{left})$ door d slides left, or rotates left when you pull it
 - $\text{side}(d,\text{right})$ door d slides right, or rotates right when you pull it

Example: Opening a Door

- Commands

- $\text{move-close}(r,o)$: r moves to a position where $\text{reachable}(r,o) = \text{T}$
- $\text{move-by}(r,\lambda)$: r moves with magnitude and direction given by vector λ
- $\text{grasp}(r,o)$: r grasps object o
- $\text{ungrasp}(r,o)$: r release grasp on o
- $\text{turn}(r,o,\alpha)$: r turns a grasped object o by angle $\alpha \in [-\pi, +\pi]$
- $\text{pull}(r,\lambda)$: r pulls its arm by vector λ
- $\text{push}(r,\lambda)$: r pushes its arm by λ
- $\text{monitor-status}(r,d)$: r focuses its perception to keep door-status updated
(starts a process)
- $\text{end-monitor-status}(r,d)$: r stops monitoring d
(terminates a process)

Example: Opening a Door

- Move close to door handle, unlatch door, pull it open

m-opendoor(r, d, l, h)

task: opendoor(r, d)

pre: $\text{loc}(r) = l \wedge \text{adjacent}(l, d)$
 $\wedge \text{handle}(d, h)$

body: while $\neg \text{reachable}(r, h)$ do
 move-close(r, h)
 monitor-status(r, d)
 if door-status(d)=closed then
 unlatch(r, d)
 throw-wide(r, d)
 end-monitor-status(r, d)

- Methods for hinged door that opens to the left, toward the robot
 - Need other methods for doors that slide, rotate to the right, or rotate away from the robot

m1-unlatch(r, d, l, o)

task: unlatch(r, d)

pre: $\text{loc}(r, l) \wedge \text{toward-side}(l, d) \wedge \text{side}(d, \text{left})$
 $\wedge \text{type}(d, \text{rotate}) \wedge \text{handle}(d, o)$

body: grasp(r, o)
 turn($r, o, \text{alpha1}$)
 pull($r, \text{val1}$)
 if door-status(d)=cracked then ungrasp(r, o)
 else fail

m1-throw-wide(r, d, l, o)

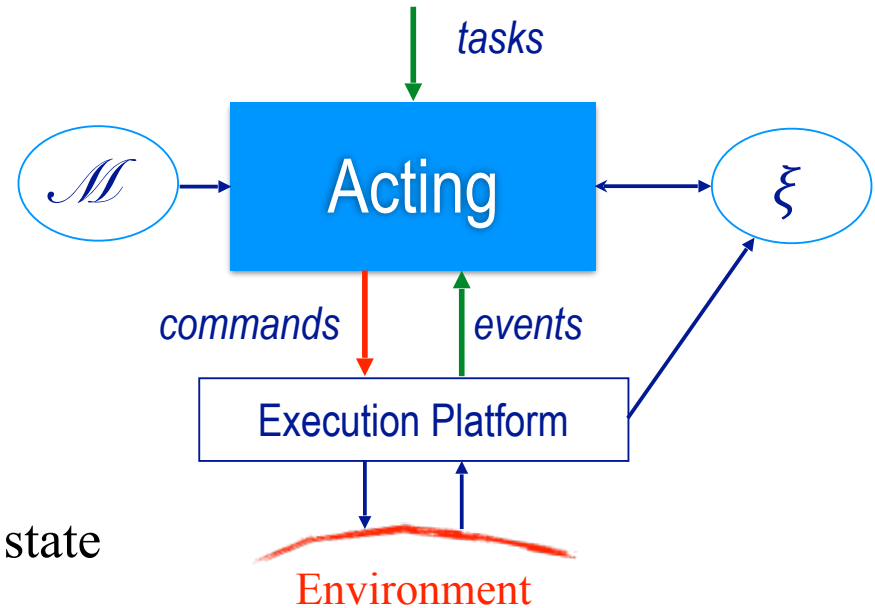
task: throw-wide(r, d)

pre: $\text{loc}(r, l) \wedge \text{toward-side}(l, d)$
 $\wedge \text{side}(d, \text{left}) \wedge \text{type}(d, \text{rotate})$
 $\wedge \text{handle}(d, o) \wedge \text{door-status}(d)=\text{cracked}$

body: grasp(r, o)
 pull($r, \text{val1}$)
 move-by($r, \text{val2}$)

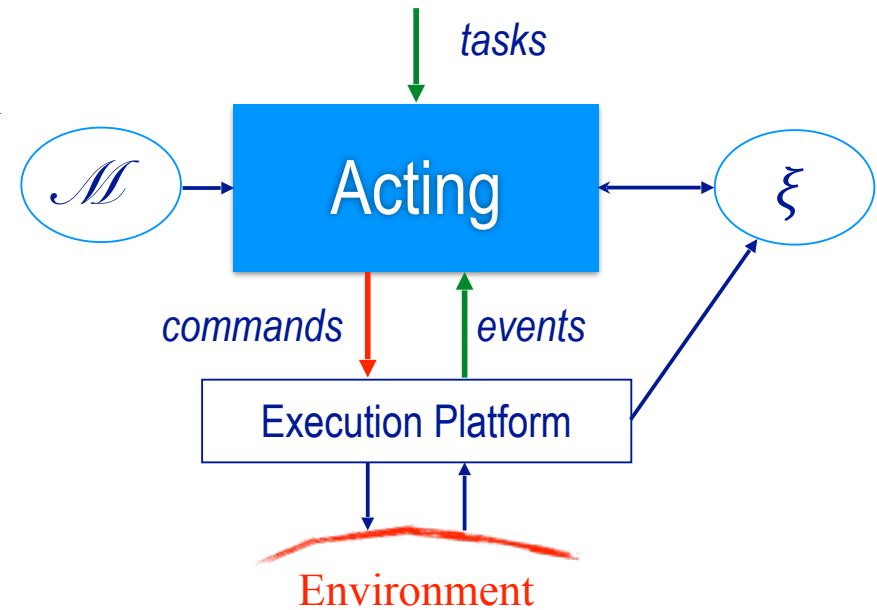
3.2 Refinement Acting Engine (RAE)

- Based on OpenPRS
 - Programming language, open-source robotics software
- Input:
 - facts, external tasks, events, current state
- Output:
 - commands to execution platform



Refinement Acting Engine (RAE)

- Perform external tasks/events in parallel
 - For each, a *refinement stack*
- $Agenda = \{\text{current stacks}\}$



Rae($\mathcal{M}$)

$Agenda \leftarrow \emptyset$

loop

until the input stream of external tasks and events is empty do

read τ in the input stream

$Candidates \leftarrow \text{Instances}(\mathcal{M}, \tau, \xi)$

if $Candidates = \emptyset$ then output("failed to address" τ)

else do

arbitrarily choose $m \in Candidates$

$Agenda \leftarrow Agenda \cup \{ \langle (\tau, m, nil, \emptyset) \rangle \}$

for each $stack \in Agenda$ do

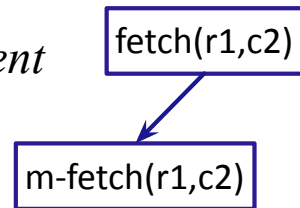
Progress($stack$)

if $stack = \emptyset$ then $Agenda \leftarrow Agenda \setminus \{stack\}$

Example

Call stack:
RAE

Refinement
tree:



Rae($\mathcal{M}$)

$Agenda \leftarrow \emptyset$

loop

until the input stream of external tasks and events is empty do

read τ in the input stream

$Candidates \leftarrow \text{Instances}(\mathcal{M}, \tau, \xi)$

if $Candidates = \emptyset$ then output("failed to address" τ)

else do

arbitrarily choose $m \in Candidates$

$Agenda \leftarrow Agenda \cup \{\langle(\tau, m, \text{nil}, \emptyset)\rangle\}$

for each $stack \in Agenda$ do

Progress($stack$)

if $stack = \emptyset$ then $Agenda \leftarrow Agenda \setminus \{stack\}$

- $\tau = \text{fetch}(r1, c2)$
- $Candidates = \{\text{m-fetch}(r1, c2)\}$
- $m = \text{m-fetch}(r1, c2)$
- $Agenda = \{\langle(\text{fetch}(r1, c2), \text{m-fetch}(r1, c2), \text{nil}, \emptyset)\rangle\}$
- *nothing else in the stream*
- $stack = \langle(\text{fetch}(r1, c2), \text{m-fetch}(r1, c2), \text{nil}, \emptyset)\rangle$
- Call Progress($stack$)

fetch(r1,c2)

Refinement
tree:

m-fetch(r1,c2)

search(r1,c2)

Call stack:
RAE
Progress¹

```

Progress(stack)
  (τ, m, i, tried) ← top(stack)
  if i ≠ nil and m[i] is a command then do
    case status(m[i])
      running: return
      failure: Retry(stack); return
      done: continue
  if i is the last step of m then
    pop(stack) // remove stack's top element
  else do

```

```

  i ← nextstep(m, i)
  case type(m[i])
    assignment: update ξ according to m[i]; return
    command: trigger command m[i]; return
    task or goal: continue
  τ' ← m[i]

```

```

Candidates ← Instances(M, τ', ξ)
if Candidates = ∅ then Retry(stack)
else do
  arbitrarily choose m' ∈ Candidates
  stack ← push((τ', m', nil, ∅), stack)

```

- $stack = \langle (fetch(r1,c2), m\text{-}fetch(r1,c2), nil, \emptyset) \rangle$
- $(\tau, m, i, tried) = (fetch(r1,c2), m\text{-}fetch(r1,c2), \text{nil}, \emptyset)$
- $i \leftarrow nextstep(m\text{-}fetch(r1,c2), nil) = i1$
- $\tau' = m[i1] = search(r1,c2)$

m-fetch(r,c)

task: fetch(r,c)

pre:

body: if pos(c) = unknown then search(r,c)
 else if loc(r) = pos(c) then take(r,c,pos(c))
 else do move-to(r,pos(c); take(r,c,pos(c))

```

Progress(stack)
  ( $\tau, m, i, \text{tried}$ )  $\leftarrow$  top(stack)
  if  $i \neq \text{nil}$  and  $m[i]$  is a command then do
    case status( $m[i]$ )
      running: return
      failure: Retry(stack); return
      done: continue
  if  $i$  is the last step of  $m$  then
    pop(stack) // remove stack's top element
  else do
     $i \leftarrow \text{nextstep}(m, i)$ 
    case type( $m[i]$ )
      assignment: update  $\xi$  according to  $m[i]$ ; return
      command: trigger command  $m[i]$ ; return
      task or goal: continue

```

```

 $\tau' \leftarrow m[i]$ 
Candidates  $\leftarrow \text{Instances}(\mathcal{M}, \tau', \xi)$ 
if Candidates =  $\emptyset$  then Retry(stack)
else do
  arbitrarily choose  $m' \in \text{Candidates}$ 
  stack  $\leftarrow \text{push}((\tau', m', \text{nil}, \emptyset), \text{stack})$ 

```

Call stack:
RAE
Progress¹

Refinement
tree:

fetch($r1, c2$)

m-fetch($r1, c2$)

search($r1, c2$)

m-search($r1, c2$)

- $\tau' = \text{search}(r1, c2)$
- Candidates = {m-search($r1, c2$)}
- $m' = \text{m-search}(r1, c2)$
- *stack* =
 $\langle (\text{search}(r1, c2), \text{m-search}(r1, c2), \text{nil}, \emptyset),$
 $(\text{fetch}(r1, c2), \text{m-fetch}(r1, c2), i1, \emptyset) \rangle$
- Return to RAE

m-fetch(r, c)

task: fetch(r, c)

pre:

body: if pos(c) = unknown then

search(r, c)

else if loc(r) = pos(c) then
take($r, c, \text{pos}(c)$)

else do move-to($r, \text{pos}(c)$)
take($r, c, \text{pos}(c)$)

Rae($\mathcal{M}$)

Agenda $\leftarrow \emptyset$

loop

 until the input stream of external tasks and events is empty do

 read τ in the input stream

 Candidates $\leftarrow$ Instances($\mathcal{M}, \tau, \xi$)

 if Candidates = $\emptyset$ then output(“failed to address” τ)

 else do

 arbitrarily choose $m \in$ Candidates

 Agenda $\leftarrow$ Agenda $\cup \{ \langle (\tau, m, \text{nil}, \emptyset) \rangle \}$

 for each stack $\in$ Agenda do

Progress(stack)

 if stack = $\emptyset$ then Agenda $\leftarrow$ Agenda $\setminus \{ \text{stack} \}$

Call stack:

RAE

Refinement tree:

```

graph TD
    A[fetch(r1,c2)] --> B[m-fetch(r1,c2)]
    B --> C[search(r1,c2)]
    C --> D[m-search(r1,c2)]
          
```

- Returned from Progress(stack)
- stack =

$\langle (\text{search}(r1,c2), \text{m-search}(r1,c2), \text{nil}, \emptyset),$
 $(\text{fetch}(r1,c2), \text{m-fetch}(r1,c2), i1, \emptyset) \rangle$
- finish “for”, iterate outer loop
- input stream empty
- Call Progress(stack)

m-fetch(r,c)

task: fetch(r,c)

pre:

body: if pos(c) = unknown then

search(r,c)

else if loc(r) = pos(c) then

 take(r,c ,pos(c))

else do move-to(r ,pos(c))

 take(r,c ,pos(c))

Nau: Lecture slides for Automated Planning and Acting

Updated 3/7/16

18

```

Progress(stack)
  ( $\tau, m, i, tried$ )  $\leftarrow$  top(stack)
  if  $i \neq \text{nil}$  and  $m[i]$  is a command then do
    case status( $m[i]$ )
      running: return
      failure: Retry(stack); return
      done: continue
  if  $i$  is the last step of  $m$  then
    pop(stack) // remove stack's top element

```

```

else do
   $i \leftarrow \text{nextstep}(m, i)$ 
  case type( $m[i]$ )
    assignment: update  $\xi$  according to  $m[i]$ ; return
    command: trigger command  $m[i]$ ; return
    task or goal: continue

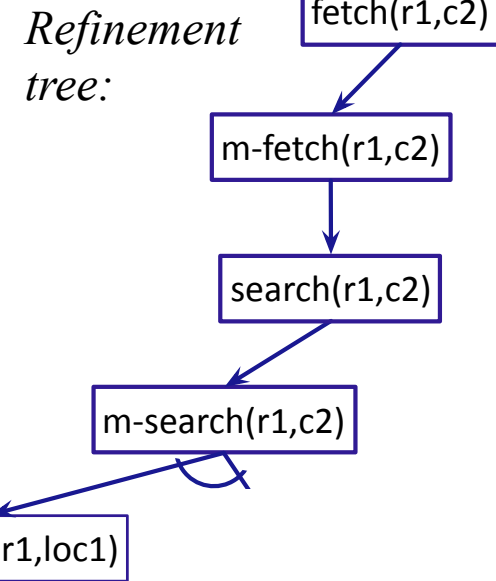
```

```

 $\tau' \leftarrow m[i]$ 
Candidates  $\leftarrow$  Instances( $\mathcal{M}, \tau', \xi$ )
if Candidates =  $\emptyset$  then Retry(stack)
else do
  arbitrarily choose  $m' \in$  Candidates
  stack  $\leftarrow$  push( $(\tau', m', \text{nil}, \emptyset), \text{stack}$ )

```

Call stack:
RAE
Progress²



- *stack* = $\langle (\text{search}(r1,c2), \text{m-search}(r1,c2), \text{nil}, \emptyset), (\text{fetch}(r1,c2), \text{m-fetch}(r1,c2), i1, \emptyset)) \rangle$
- $l \leftarrow \text{loc1}; i \leftarrow i2$
- trigger $m[i2] = \text{move-to}(r1, \text{loc1})$
- return to RAE; RAE calls Progress again

$r=r1; c=c2$

```

m-search( $r, c$ )
task: search( $r, c$ )
pre: pos( $c$ ) = unknown
body: if  $\exists l (\text{view}(l)=F)$  then
  move-to( $r, l$ )
  perceive( $l$ )

```

```

    if pos( $c$ )= $l$  then take( $r, c, l$ ) else search( $r, c$ )
  else fail

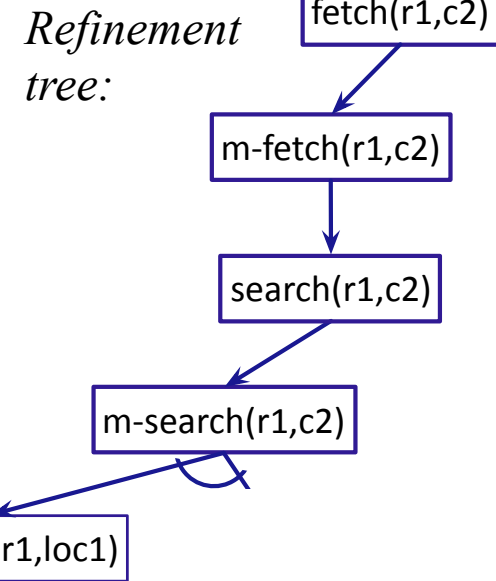
```

```

Progress(stack)
  ( $\tau, m, i, tried$ )  $\leftarrow$  top(stack)
  if  $i \neq \text{nil}$  and  $m[i]$  is a command then do
    case status( $m[i]$ )
      running: return
      failure: Retry(stack); return
      done: continue
  if  $i$  is the last step of  $m$  then
    pop(stack) // remove stack's top element
  else do
     $i \leftarrow \text{nextstep}(m, i)$ 
    case type( $m[i]$ )
      assignment: update  $\xi$  according to  $m[i]$ ; return
      command: trigger command  $m[i]$ ; return
      task or goal: continue
     $\tau' \leftarrow m[i]$ 
    Candidates  $\leftarrow$  Instances( $\mathcal{M}, \tau', \xi$ )
    if Candidates =  $\emptyset$  then Retry(stack)
    else do
      arbitrarily choose  $m' \in \text{Candidates}$ 
      stack  $\leftarrow$  push( $(\tau', m', \text{nil}, \emptyset), \text{stack}$ )

```

Call stack:
RAE
Progress³



- $stack = \langle (\text{search}(r1, c2), \text{m-search}(r1, c2), i2, \emptyset), (\text{fetch}(r1, c2), \text{m-fetch}(r1, c2), i1, \emptyset) \rangle$
- status($m[i2]$) = running
- return to RAE; RAE calls Progress again

$r=r1; c=c2$

```

m-search( $r, c$ )
  task: search( $r, c$ )
  pre: pos( $c$ ) = unknown
  body: if  $\exists l (\text{view}(l)=F)$  then
    move-to( $r, l$ )
    perceive( $l$ )
    if pos( $c$ )= $l$  then take( $r, c, l$ ) else search( $r, c$ )
  else fail

```

```

Progress(stack)
  ( $\tau, m, i, \text{tried}$ )  $\leftarrow$  top(stack)
  if  $i \neq \text{nil}$  and  $m[i]$  is a command then do
    case status( $m[i]$ )
      running: return
      failure: Retry(stack); return
      done: continue
  if  $i$  is the last step of  $m$  then
    pop(stack) // remove stack's top element
  else do

```

```

   $i \leftarrow \text{nextstep}(m, i)$ 
  case type( $m[i]$ )
    assignment: update  $\xi$  according to  $m[i]$ ; return
    command: trigger command  $m[i]$ ; return
    task or goal: continue

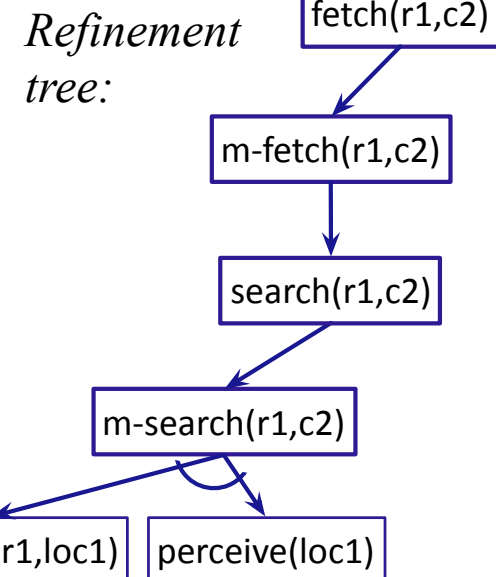
```

```

   $\tau' \leftarrow m[i]$ 
  Candidates  $\leftarrow$  Instances( $\mathcal{M}, \tau', \xi$ )
  if Candidates =  $\emptyset$  then Retry(stack)
  else do
    arbitrarily choose  $m' \in \text{Candidates}$ 
    stack  $\leftarrow$  push( $(\tau', m', \text{nil}, \emptyset), \text{stack}$ )

```

Call stack:
RAE
Progress⁴



- *stack*
= $\langle (\text{search}(r1,c2), \text{m-search}(r1,c2), i2, \emptyset), (\text{fetch}(r1,c2), \text{m-fetch}(r1,c2), i1, \emptyset)) \rangle$
- status($m[i2]$) = done; $i \leftarrow i3$
- trigger perceive(loc1)
- return to RAE; RAE calls Progress again

$r=r1; c=c2$

```

m-search( $r, c$ )
  task: search( $r, c$ )
  pre: pos( $c$ ) = unknown
  body: if  $\exists l (\text{view}(l)=F)$  then
    move-to( $r, l$ )
    perceive( $l$ )

```

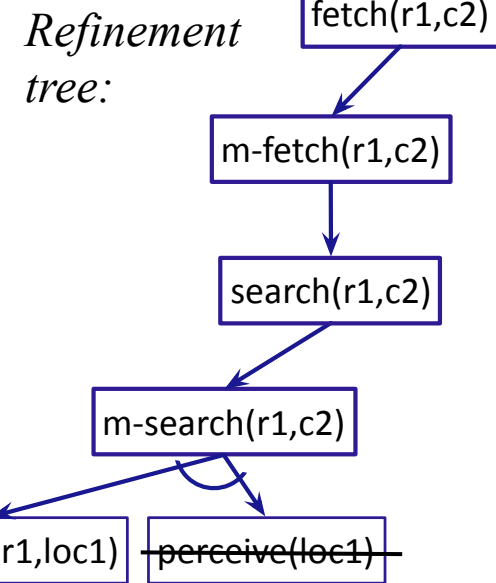
if pos(c)= l then take(r, c, l) else search(r, c)
else fail

```

Progress(stack)
  ( $\tau, m, i, tried$ )  $\leftarrow$  top(stack)
  if  $i \neq \text{nil}$  and  $m[i]$  is a command then do
    case status( $m[i]$ )
      running: return
      failure: Retry(stack); return
      done: continue
  if  $i$  is the last step of  $m$  then
    pop(stack) // remove stack's top element
  else do
     $i \leftarrow \text{nextstep}(m, i)$ 
    case type( $m[i]$ )
      assignment: update  $\xi$  according to  $m[i]$ ; return
      command: trigger command  $m[i]$ ; return
      task or goal: continue
     $\tau' \leftarrow m[i]$ 
    Candidates  $\leftarrow$  Instances( $\mathcal{M}, \tau', \xi$ )
    if Candidates =  $\emptyset$  then Retry(stack)
    else do
      arbitrarily choose  $m' \in$  Candidates
      stack  $\leftarrow$  push( $(\tau', m', \text{nil}, \emptyset), \text{stack}$ )

```

Call stack:
RAE
Progress⁵



- *stack*
= $\langle (\text{search}(r1,c2), \text{m-search}(r1,c2), i3, \emptyset), (\text{fetch}(r1,c2), \text{m-fetch}(r1,c2), i1, \emptyset) \rangle$
- $m[i3] = \text{perceive}(\text{loc1})$
- status($m[i3]$) = failure (*sensor failure*)
- call Retry

$r=r1; c=c2$

```

m-search( $r, c$ )
  task: search( $r, c$ )
  pre: pos( $c$ ) = unknown
  body: if  $\exists l$  (view( $l$ )=F) then
    move-to( $r, l$ )
    perceive( $l$ )
    if pos( $c$ )= $l$  then take( $r, c, l$ ) else search( $r, c$ )
  else fail

```

Retry(*stack*)

```

( $\tau, m, i, \text{tried}$ )  $\leftarrow$  pop(stack)
 $\text{tried} \leftarrow \text{tried} \cup \{m\}$ 
Candidates  $\leftarrow$  Instances( $\mathcal{M}, \tau, \xi$ ) \ tried
if Candidates  $\neq \emptyset$  then do
    arbitrarily choose  $m' \in$  Candidates
    stack  $\leftarrow$  push( $(\tau, m', \text{nil}, \text{tried})$ , stack)
else do
    if stack  $\neq \emptyset$  then Retry(stack)
    else do
        output("failed to accomplish"  $\tau$ )
        Agenda  $\leftarrow$  Agenda \ stack

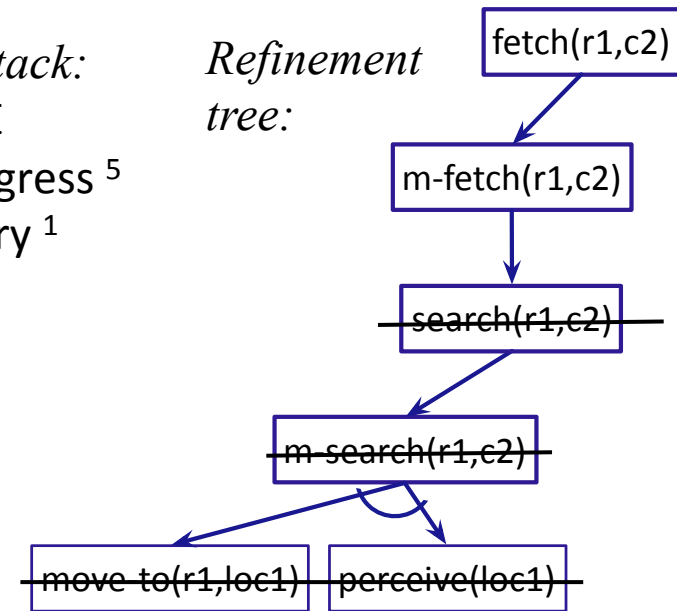
```

- $\text{stack} = \langle (\text{search}(r1, c2), \text{m-search}(r1, c2), i3, \emptyset), (\text{fetch}(r1, c2), \text{m-fetch}(r1, c2), i1, \emptyset) \rangle$
- $(\tau, m, i, \text{tried}) \leftarrow (\text{search}(r1, c2), \text{m-search}(r1, c2), i3, \emptyset)$
- $\text{tried} \leftarrow \{\text{m-search}(r1, c2)\}$
- $\text{Instances}(\mathcal{M}, \tau, \xi) = \{\text{m-search}(r1, c2)\}$
- Candidates $\leftarrow$
 $\text{Instances}(\mathcal{M}, \tau, \xi) \setminus \text{tried}$
 $= \emptyset$
- Call Retry(*stack*)

Call stack:

RAE
Progress⁵
Retry¹

Refinement
tree:



m-search(*r, c*)

task: search(*r, c*)

pre: pos(*c*) = unknown

body: if $\exists l (\text{view}(l) = F)$ then

 move-to(*r, l*)

 perceive(*l*)

 if pos(*c*) = *l* then take(*r, c, l*) else search(*r, c*)

else fail

Retry(*stack*)

```

( $\tau, m, i, \text{tried}$ )  $\leftarrow$  pop(stack)
 $\text{tried} \leftarrow \text{tried} \cup \{m\}$ 
 $\text{Candidates} \leftarrow \text{Instances}(\mathcal{M}, \tau, \xi) \setminus \text{tried}$ 
if  $\text{Candidates} \neq \emptyset$  then do
    arbitrarily choose  $m' \in \text{Candidates}$ 
     $\text{stack} \leftarrow \text{push}((\tau, m', \text{nil}, \text{tried}), \text{stack})$ 
else do
    if  $\text{stack} \neq \emptyset$  then Retry(stack)
    else do
        output("failed to accomplish"  $\tau$ )
         $\text{Agenda} \leftarrow \text{Agenda} \setminus \text{stack}$ 

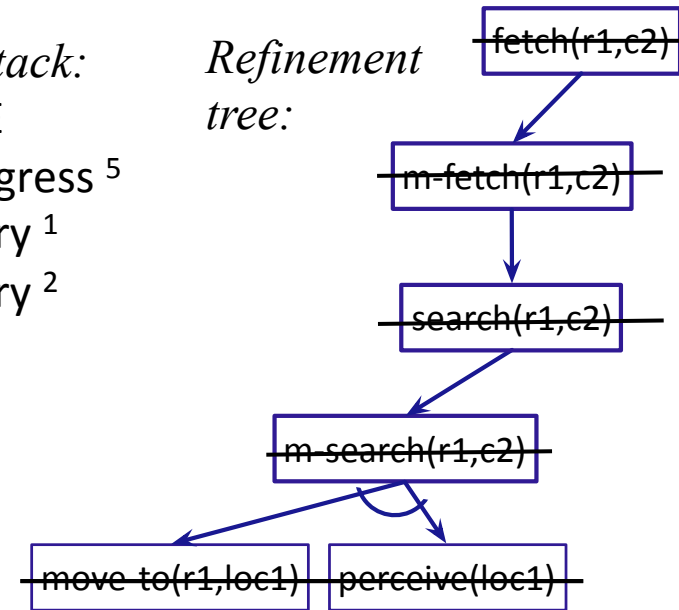
```

- $\text{stack} = \langle (\text{fetch}(r1, c2), m\text{-fetch}(r1, c2), i1, \emptyset) \rangle$
- $(\tau, m, i, \text{tried})$
 $\leftarrow (\text{fetch}(r1, c2), m\text{-fetch}(r1, c2), i1, \emptyset)$
- $\text{tried} \leftarrow \{m\text{-fetch}(r1, c2)\}$
- $\text{Instances}(\mathcal{M}, \tau, \xi) = \{m\text{-fetch}(r1, c2)\}$
- $\text{Candidates} \leftarrow \text{Instances}(\mathcal{M}, \tau, \xi) \setminus \text{tried} = \emptyset$
- output("failed to accomplish" $\text{fetch}(r1, c2)$)
- Return to first Retry
 - Return to Progress
 - Return to RAE

Call stack:

RAE
 Progress⁵
 Retry¹
 Retry²

Refinement
 tree:



$m\text{-fetch}(r, c)$

task: $\text{fetch}(r, c)$

pre:

body: if $\text{pos}(c) = \text{unknown}$ then

$\text{search}(r, c)$

else if $\text{loc}(r) = \text{pos}(c)$ then

take($r, c, \text{pos}(c)$)

else do move-to($r, \text{pos}(c)$)

take($r, c, \text{pos}(c)$)

Example

Call stack:
RAE

Rae($\mathcal{M}$)

$Agenda \leftarrow \emptyset$

loop

until the input stream of external tasks and events is empty do

read τ in the input stream

$Candidates \leftarrow \text{Instances}(\mathcal{M}, \tau, \xi)$

if $Candidates = \emptyset$ then output("failed to address" τ)

else do

arbitrarily choose $m \in Candidates$

$Agenda \leftarrow Agenda \cup \{ \langle (\tau, m, \text{nil}, \emptyset) \rangle \}$

for each $stack \in Agenda$ do

Progress($stack$)

if $stack = \emptyset$ then $Agenda \leftarrow Agenda \setminus \{stack\}$

- $Agenda = \emptyset$
- repeatedly execute loop, waiting for input

Goals in RAE

- Special kind of task: `achieve(condition)`
 - Like other tasks, but includes monitoring
- Can easily modify Progress to do this
 - Stop executing `body(m)` early if *condition* becomes true
 - fail if *condition* isn't true after executing `body(m)`
 - triggers a Retry
 - `m-fetch(r,c)`
 - task: `fetch(r,c)`
 - pre:
 - body: `achieve(pos(c) ≠ unknown)`
 - `goto(pos(c))`
 - `take(r,c)`
 - `m-find-where(r,c)`
 - goal: `achieve(pos(c) ≠ unknown)`
 - pre:
 - body: while $\exists l \text{ (view}(l)=F)$ do
 - `goto(l)`
 - `perceive(l)`

- What we had before:

```
m-fetch(r,c)
  task: fetch(r,c)
  pre:
  body: if pos(c) = unknown then
        search(r,c)
      else if loc(r) = pos(c) then
        take(r,c,pos(c))
      else do
        move-to(r,pos(c))
        take(r,c,pos(c))
```

```
m-search(r,c)
  task: search(r,c)
  pre: pos(c) = unknown
  body: if  $\exists l \text{ (view}(l)=F)$  then
        move-to(r,l); perceive(l)
      if pos(c)=l then take(r,c,l)
      else search(r,c)
  else fail
```

Extensions to RAE

- Section 3.2.4 discusses additional extensions to RAE:
 - Controlling the progress of tasks, e.g., suspend a task for a while
 - Concurrent subtasks
 - $m\text{-foo}(x,y,z)$
task: $\text{foo}(x,y,z)$
pre:
body: ...
 $\{\text{concurrent: } \langle v_{1,1}, \dots, v_{1,n} \rangle \langle v_{2,1}, \dots, v_{2,m} \rangle \dots \langle v_{k,1}, \dots, v_{k,l} \rangle\}$
 - Each $\langle v_{i,1}, \dots, v_{i,j} \rangle$ is a sequence of steps, like the body of a method
 - Split the current stack into k sub-stacks, one for each $\langle v_{i,1}, \dots, v_{i,j} \rangle$
 - If there are multiple stacks, which ones get higher priority?
 - Application-specific heuristics
 - For a task t , which method to try first?
 - Refinement planning (Section 3.3)

3.3 Refinement planning (SeRPE)

- When dealing with an event or task, RAE may need to make either/or choices
 - $Agenda = \{\langle \tau_1, \dots \rangle, \langle \tau_2, \dots \rangle, \dots, \langle \tau_n, \dots \rangle\}$
 - Several tasks/events, which one to address first?
 - $Candidates = \{m_1, m_2, \dots, \}$
 - Several candidate methods or commands, which one to try first?
- RAE immediately executes commands
 - Bad choices may be costly or irreversible
- Look ahead to predict the consequences
 - Instead of executing commands immediately, simulate effects using descriptive action models (like in Chap. 2)

Sequential Refinement Planning Engine (SeRPE)

- Refines tasks similarly to RAE, except:
 - Where RAE executes a command that changes the world state, SeRPE instead uses a descriptive action model to predict the next world state

Restrictive assumptions:

- All action models need to have deterministic outcomes (as in Chapter 2)
 - e.g., predict exactly what grasp, turn, pull will do
- Don't allow concurrent tasks
 - {concurrent:
 $\langle v_{1,1}, \dots, v_{1,n} \rangle \langle v_{2,1}, \dots, v_{2,m} \rangle \dots \langle v_{k,1}, \dots, v_{k,l} \rangle \}$
 - We'll relax this restriction later

```
m1-unlatch(r,d,l,o)  
task: unlatch(r,d)  
pre: ...  
body: grasp(r,o)  
      turn(r,o,alpha1)  
      pull(r,val1)  
      if door-status(d)=cracked  
      then ungrasp(r,o)  
      else fail
```

Planning Problem

- Planning problem

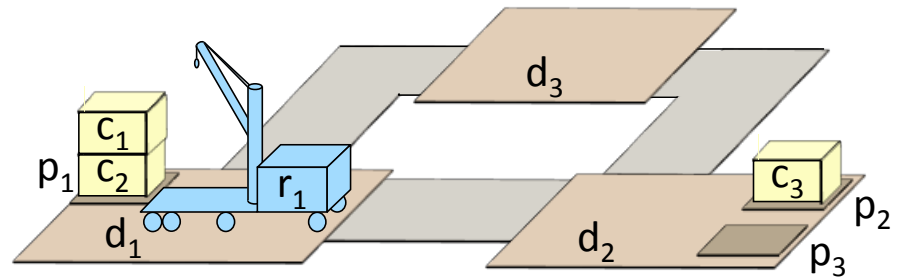
$$P = (\mathcal{M}, \mathcal{A}, s_0, \tau)$$

- $\mathcal{M}$: refinement methods
- $\mathcal{A}$: action templates
- s_0 : initial state
- τ : task or goal
 - $\text{put-in-pile}(c_1, p_2)$

Example:

Rigid relations

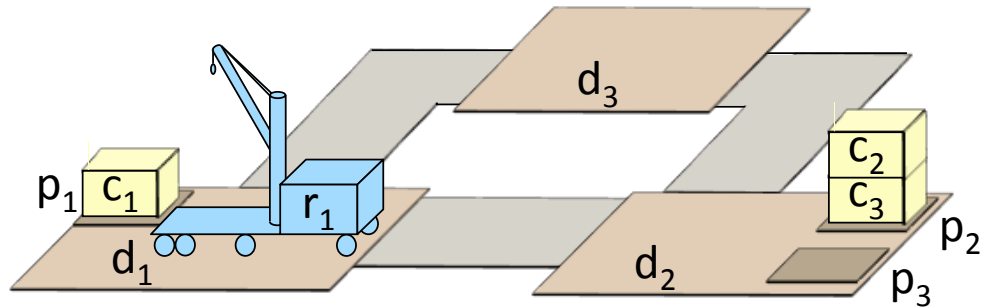
- $\text{at}(p_1, d_1), \text{at}(p_2, d_2), \text{at}(p_3, d_2)$
- $\text{adjacent}(d_i, d_j)$ for every i, j



$$s_0 = \{\text{loc}(r_1)=d_1, \text{cargo}(r_1)=\text{nil}, \\ \text{occupied}(d_1)=\text{T}, \text{occupied}(d_2)=\text{F}, \text{occupied}(d_3)=\text{F}, \\ \text{pos}(c_1)=\text{nil}, \text{pos}(c_2)=c_3, \text{pos}(c_3)=\text{nil}, \\ \text{pile}(c_1)=p_1, \text{pile}(c_2)=p_2, \text{pile}(c_3)=p_2, \\ \text{top}(p_1)=c_1, \text{top}(p_2)=c_2, \text{top}(p_3)=\text{nil}\}$$

Action Templates

- From Example 2.12:



$\text{load}(r, c, c', p, d)$

pre: $\text{at}(p, d)$, $\text{cargo}(r) = \text{nil}$, $\text{loc}(r) = d$, $\text{pos}(c) = c'$, $\text{top}(p) = c$

eff: $\text{cargo}(r) \leftarrow c$, $\text{pile}(c) \leftarrow \text{nil}$, $\text{pos}(c) \leftarrow r$, $\text{top}(p) \leftarrow c'$

$\text{unload}(r, c, c', p, d)$

pre: $\text{at}(p, d)$, $\text{pos}(c) = r$, $\text{loc}(r) = d$, $\text{top}(p) = c'$

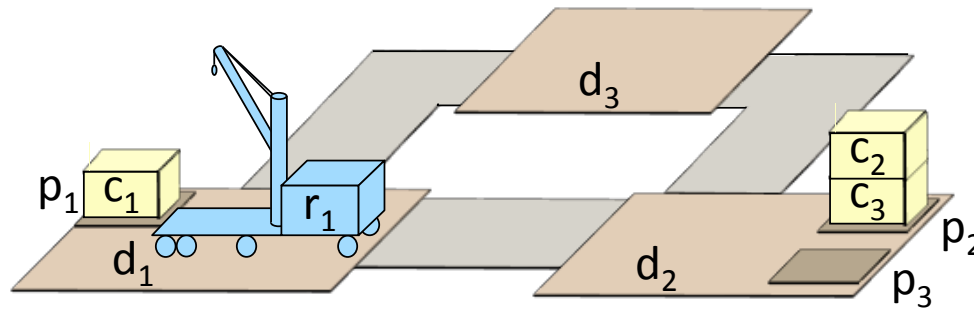
eff: $\text{cargo}(r) \leftarrow \text{nil}$, $\text{pile}(c) \leftarrow p$, $\text{pos}(c) \leftarrow c'$, $\text{top}(p) \leftarrow c$

$\text{move}(r, d, d')$

pre: $\text{adjacent}(d, d')$, $\text{loc}(r) = d$, $\text{occupied}(d') = \text{F}$

eff: $\text{loc}(r) = d'$, $\text{occupied}(d) = \text{F}$, $\text{occupied}(d') = \text{T}$

Tasks and Methods



$\text{put-in-pile}(c, p')$ – put c into pile p' if it isn't already there

$\text{m1-put-in-pile}(c, p')$

task: $\text{put-in-pile}(c, p')$

pre: $\text{pile}(c) = p'$

body: *// empty*

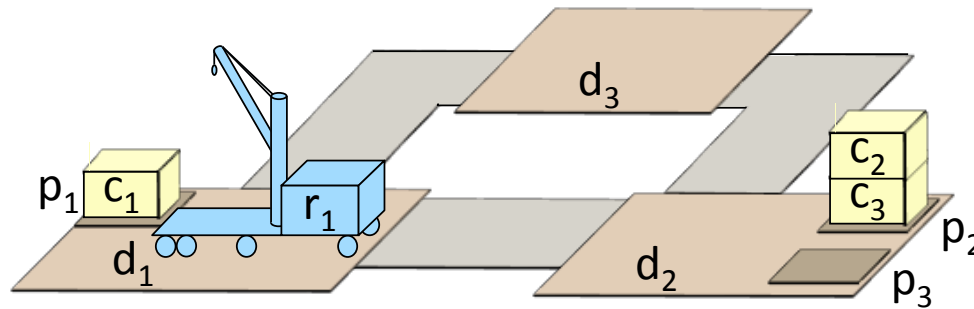
$\text{m2-put-in-pile}(r, c, p, d, p', d')$

task: $\text{put-in-pile}(c, p')$

pre: $\text{pile}(c) = p \wedge \text{at}(p, d) \wedge \text{at}(p', d') \wedge p \neq p' \wedge \text{cargo}(r) = \text{nil}$

body: if $\text{loc}(r) \neq d$ then $\text{navigate}(r, d)$
 $\text{uncover}(c)$
 $\text{load}(r, c, \text{pos}(c), p, d)$
 if $\text{loc}(r) \neq d'$ then $\text{navigate}(r, d')$
 $\text{unload}(r, c, \text{top}(p'), p', d)$

Tasks and Methods



$\text{uncover}(c)$ – ensure no containers are on c

$\text{m1-uncover}(c)$

task: $\text{uncover}(c)$

pre: $\text{top}(\text{pile}(c))=c$

body: $// \text{ empty}$

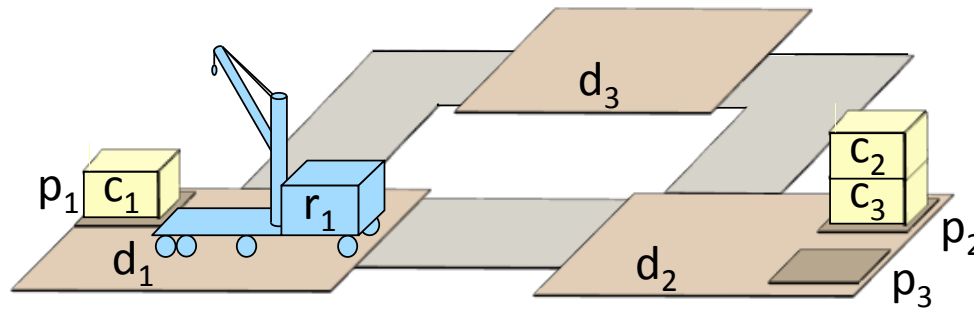
$\text{m2-uncover}(r, c, c, p', d)$

task: $\text{uncover}(c)$

pre: $\text{pile}(c)=p \wedge \text{top}(p) \neq c$
 $\wedge \text{at}(p, d) \wedge \text{at}(p', d) \wedge p' \neq p$
 $\wedge \text{loc}(r)=d \wedge \text{cargo}(r)=\text{nil}$

body: while $\text{top}(p) \neq c$ do
 $c' \leftarrow \text{top}(p)$
 $\text{load}(r, c', \text{pos}(c'), p, d)$
 $\text{unload}(r, c', \text{top}(p'), p', d)$

Tasks and Methods



$\text{navigate}(r, d')$ – move r along some route that ends at loading dock d'

$\text{m1-navigate}(r, d')$

task: $\text{navigate}(r, d')$

pre: $\text{loc}(r) = d'$

body: *// empty*

$\text{m2-navigate}(r, d')$

task: $\text{navigate}(r, d')$

pre: $\text{loc}(r) \neq d' \wedge \text{adjacent}(\text{loc}(r), d')$

body: $\text{move}(r, \text{loc}(r), d')$

$\text{m3-navigate}(r, d, d')$

task: $\text{navigate}(r, d')$

pre: $\text{loc}(r) \neq d' \wedge d \neq d' \wedge$
 $\text{adjacent}(\text{loc}(r), d)$

body: $\text{move}(r, \text{loc}(r), d)$
 $\text{navigate}(r, d')$

SeRPE Algorithm

```

Progress-to-finish( $\mathcal{M}, \mathcal{A}, s, \tau, m$ )
   $i \leftarrow \text{nil}$       // instruction pointer for body( $m$ )
   $\pi \leftarrow \langle \rangle$     // plan produced from body( $m$ )
  loop
    if  $\tau$  is a goal and  $s \models \tau$  then return  $\pi$ 
    if  $i$  is the last step of  $m$  then
      if  $\tau$  is a goal and  $s \not\models \tau$  then return failure
      return  $\pi$ 
     $i \leftarrow \text{nextstep}(m, i)$ 
  case type( $m[i]$ )
    assignment: update  $s$  according to  $m[i]$ 
    command:
       $a \leftarrow$  the descriptive model of  $m[i]$  in  $A$ 
      if  $s \models \text{pre}(a)$  then
         $s \leftarrow \gamma(s, a)$ ;  $\pi \leftarrow \pi.a$ 
      else return failure
    task or goal:
       $\pi' \leftarrow \text{SeRPE}(\mathcal{M}, \mathcal{A}, s, m[i])$ 
      if  $\pi' = \text{failure}$  then return failure
       $s \leftarrow \gamma(s, \pi')$ ;  $\pi \leftarrow \pi.\pi'$ 

SeRPE( $\mathcal{M}, \mathcal{A}, s, \tau$ )
   $Candidates \leftarrow \text{Instances}(\mathcal{M}, \tau, s)$ 
  if  $Candidates = \emptyset$  then return failure
  nondeterministically choose  $m \in Candidates$ 
  return Progress-to-finish( $\mathcal{M}, \mathcal{A}, s, \tau, m$ )
  
```

Example

$\tau = \text{put-in-pile}(c_1, p_2)$

Candidates = {~~m1-put-in-pile(c_1, p_2)~~,
m2-put-in-pile($r_1, c_1, p_1, d_1, p_2, d_2$), ...}

Here, implementation would use a lifted expression m2-put-in-pile(r, c_1, p_1, d, p', d'). It would instantiate variables here:

SeRPE($\mathcal{M}, \mathcal{A}, s, \tau$)

Candidates $\leftarrow$ Instances($\mathcal{M}, \tau, s$)

if *Candidates* = $\emptyset$ then return failure

nondeterministically choose $m \in \text{Candidates}$

return Progress-to-finish($\mathcal{M}, \mathcal{A}, s, \tau, m$)

$s_0 = \{\text{loc}(r_1)=d_1, \text{cargo}(r_1)=\text{nil}, \text{occupied}(d_1)=T,$
 $\text{occupied}(d_2)=F, \text{occupied}(d_3)=F,$
 $\text{pos}(c_1)=\text{nil}, \text{pos}(c_2)=c_3, \text{pos}(c_3)=\text{nil},$
 $\text{pile}(c_1)=p_1, \text{pile}(c_2)=p_2, \text{pile}(c_3)=p_2,$
 $\text{top}(p_1)=c_1, \text{top}(p_2)=c_2, \text{top}(p_3)=\text{nil}\}$

m1-put-in-pile(c, p')

task: put-in-pile(c, p')

pre: pile(c)= p'

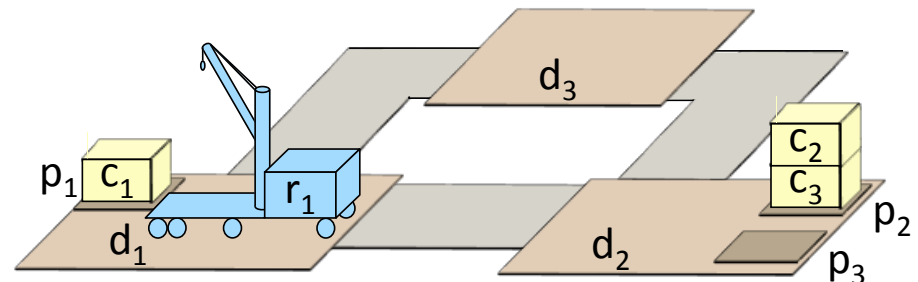
body: // empty

m2-put-in-pile(r, c, p, d, p', d')

task: put-in-pile(c, p')

pre: pile(c)= $p \wedge \text{at}(p, d) \wedge \text{at}(p', d)$
 $\wedge p \neq p' \wedge \text{cargo}(r)=\text{nil}$

body: if loc(r) $\neq d$ then navigate(r, d)
 uncover(c)
 load($r, c, \text{pos}(c), p, d$)
 if loc(r) $\neq d'$ then
 navigate(r, d')
 unload($r, c, \text{top}(p'), p', d$)



Example

task
put-in-pile(c_1, p_2)
|
method
m2-put-in-pile($r_1, c_1, p_1, d_1, p_2, d_2$)

Refinement tree

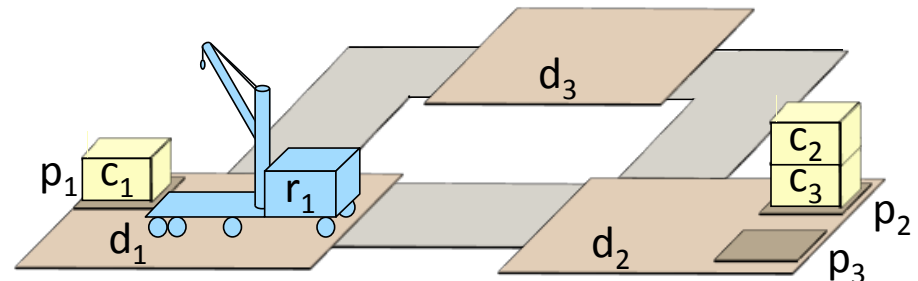
SeRPE($\mathcal{M}, \mathcal{A}, s, \tau$)

$Candidates \leftarrow Instances(\mathcal{M}, \tau, s)$
if $Candidates = \emptyset$ then return failure
nondeterministically choose $m \in Candidates$
return Progress-to-finish($\mathcal{M}, \mathcal{A}, s, \tau, m$)

$s_0 = \{loc(r_1)=d_1, cargo(r_1)=nil, occupied(d_1)=T,$
 $occupied(d_2)=F, occupied(d_3)=F,$
 $pos(c_1)=nil, pos(c_2)=c_3, pos(c_3)=nil,$
 $pile(c_1)=p_1, pile(c_2)=p_2, pile(c_3)=p_2,$
 $top(p_1)=c_1, top(p_2)=c_2, top(p_3)=nil\}$

$r_1, c_1, p_1, d_1, p_2, d_2$

m2-put-in-pile(r, c, p, d, p', d')
task: put-in-pile(c, p')
pre: $pile(c)=p \wedge at(p, d) \wedge at(p', d)$
 $\wedge p \neq p' \wedge cargo(r)=nil$
body: if $loc(r) \neq d$ then navigate(r, d)
uncover(c)
load($r, c, pos(c), p, d$)
if $loc(r) \neq d'$ then
navigate(r, d')
unload($r, c, top(p'), p', d$)



Example

Progress-to-finish($\mathcal{M}, \mathcal{A}, s, \tau, m$)

$i \leftarrow \text{nil}$ // instruction pointer for body(m)

$\pi \leftarrow \langle \rangle$ // plan produced from body(m)

loop

if τ is a goal and $s \models \tau$ then return π

if i is the last step of m then

if τ is a goal and $s \not\models \tau$ then return failure

return π

$i \leftarrow \text{nextstep}(m, i)$

case type($m[i]$)

assignment: update s according to $m[i]$

command:

$a \leftarrow$ the descriptive model of $m[i]$ in A

if $s \models \text{pre}(a)$ then

$s \leftarrow \gamma(s, a); \pi \leftarrow \pi.a$

else return failure

task or goal:

$\pi' \leftarrow \text{SeRPE}(\mathcal{M}, \mathcal{A}, s, m[i])$

if $\pi' = \text{failure}$ then return failure

$s \leftarrow \gamma(s, \pi'); \pi \leftarrow \pi.\pi'$

$r_1, c_1, p_1, d_1, p_2, d_2$

m2-put-in-pile(r, c, p, d, p', d')

task: put-in-pile(c, p')

pre: $\text{pile}(c)=p \wedge \text{at}(p, d) \wedge \text{at}(p', d)$
 $\wedge p \neq p' \wedge \text{cargo}(r)=\text{nil}$

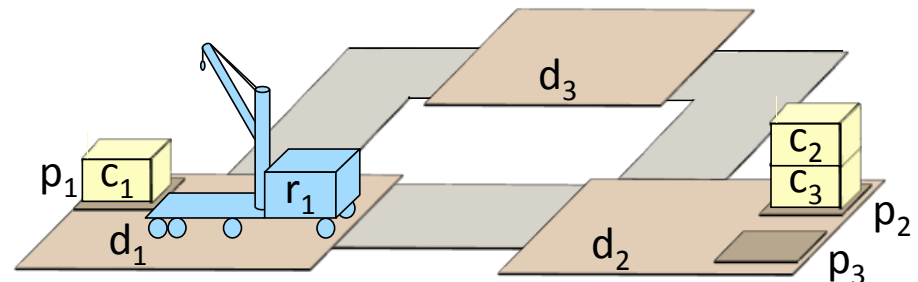
body: if $\text{loc}(r) \neq d$ then navigate(r, d)
 uncover(c)

load($r, c, \text{pos}(c), p, d$)

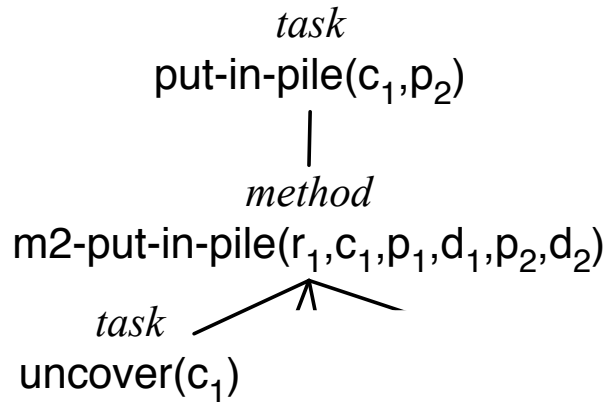
if $\text{loc}(r) \neq d'$ then

navigate(r, d')

unload($r, c, \text{top}(p'), p', d$)



Example



$r_1, c_1, p_1, d_1, p_2, d_2$

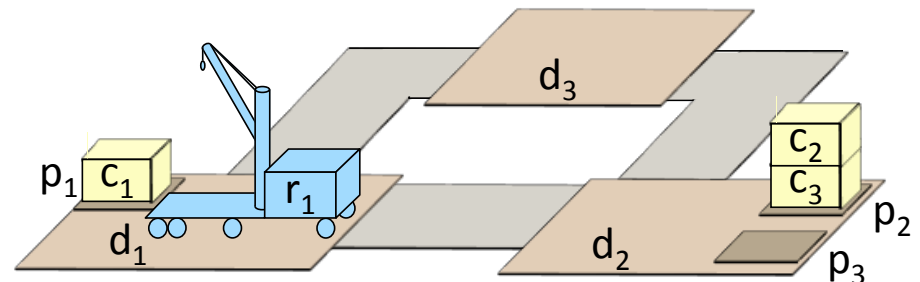
m2-put-in-pile(r, c, p, d, p', d')

task: put-in-pile(c, p')

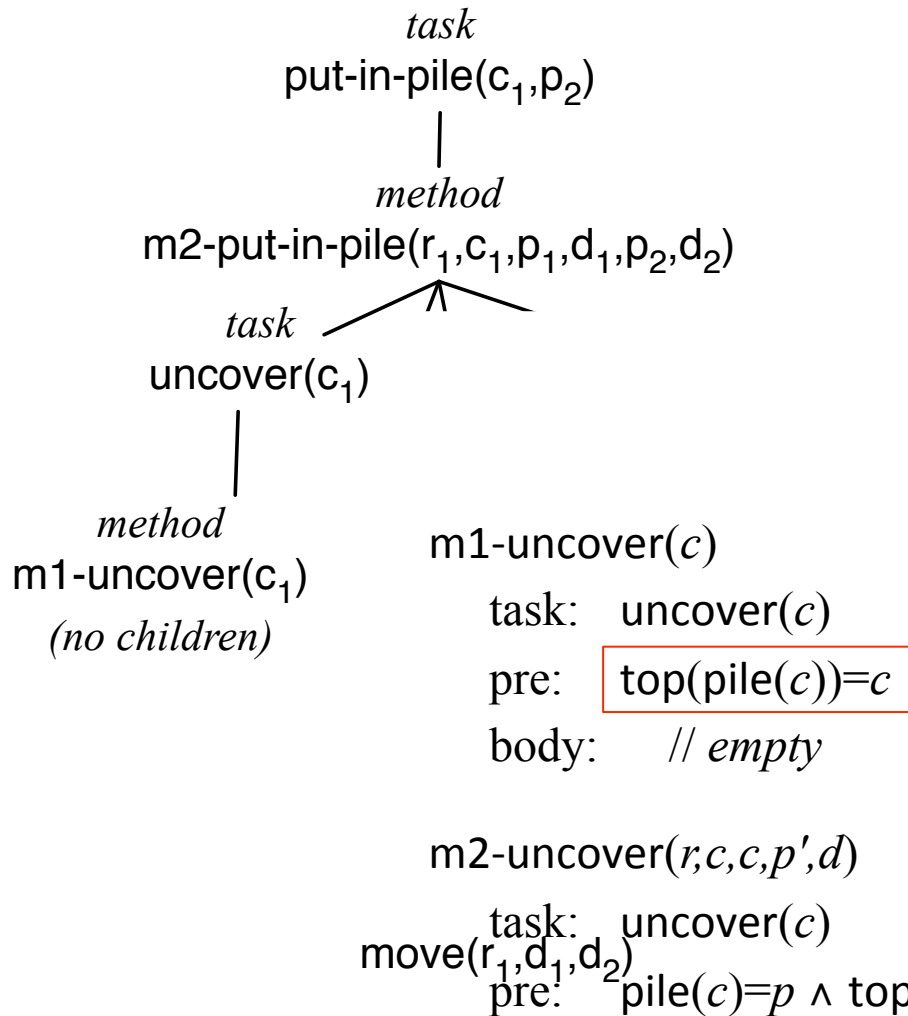
pre: $\text{pile}(c)=p \wedge \text{at}(p, d) \wedge \text{at}(p', d)$
 $\wedge p \neq p' \wedge \text{cargo}(r)=\text{nil}$

body: if $\text{loc}(r) \neq d$ then navigate(r, d)
 uncover(c)
 load($r, c, \text{pos}(c), p, d$)
 if $\text{loc}(r) \neq d'$ then
 navigate(r, d')
 unload($r, c, \text{top}(p'), p', d$)

$s_0 = \{\text{loc}(r_1)=d_1, \text{cargo}(r_1)=\text{nil}, \text{occupied}(d_1)=T,$
 $\text{occupied}(d_2)=F, \text{occupied}(d_3)=F,$
 $\text{pos}(c_1)=\text{nil}, \text{pos}(c_2)=c_3, \text{pos}(c_3)=\text{nil},$
 $\text{pile}(c_1)=p_1, \text{pile}(c_2)=p_2, \text{pile}(c_3)=p_2,$
 $\text{top}(p_1)=c_1, \text{top}(p_2)=c_2, \text{top}(p_3)=\text{nil}\}$



Example



$r_1, c_1, p_1, d_1, p_2, d_2$

m2-put-in-pile(r, c, p, d, p', d')

task: put-in-pile(c, p')

pre: $\text{pile}(c) = p \wedge \text{at}(p, d) \wedge \text{at}(p', d) \wedge p \neq p' \wedge \text{cargo}(r) = \text{nil}$

body: if $\text{loc}(r) \neq d$ then navigate(r, d)

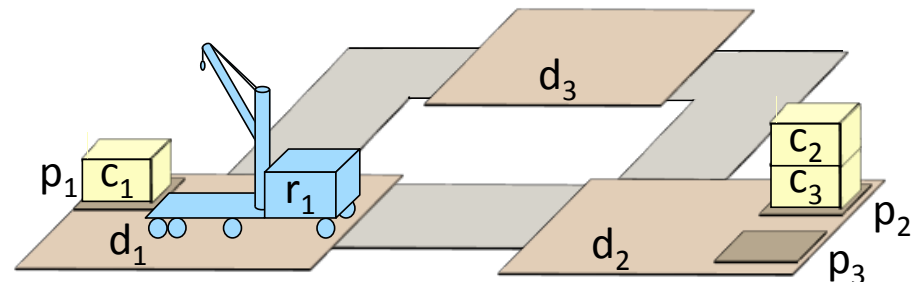
uncover(c)

load($r, c, \text{pos}(c), p, d$)

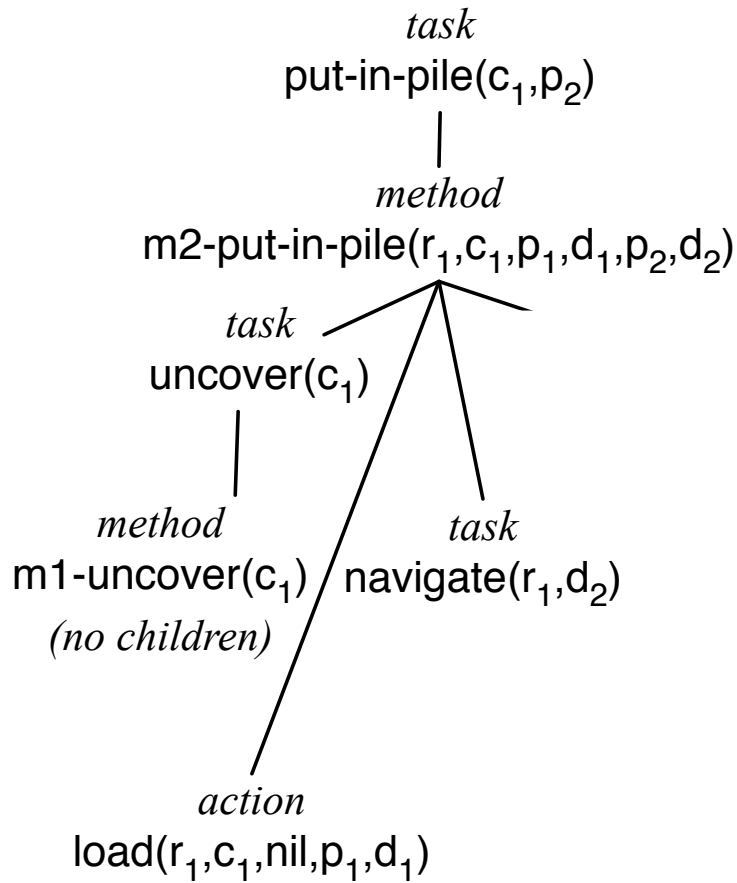
if $\text{loc}(r) \neq d'$ then

navigate(r, d')

unload($r, c, \text{top}(p'), p', d$)



Example



$r_1, c_1, p_1, d_1, p_2, d_2$

m2-put-in-pile(r, c, p, d, p', d')

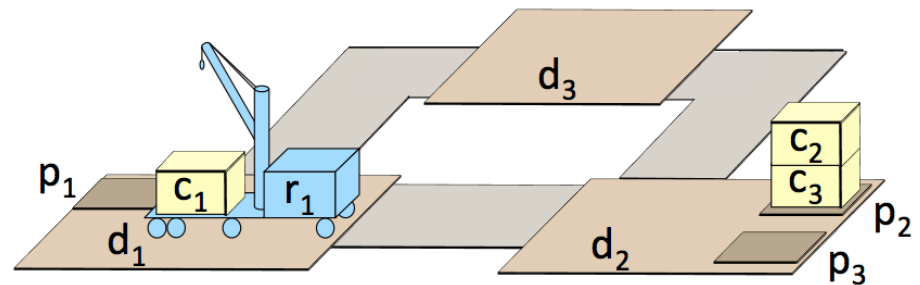
task: put-in-pile(c, p')

pre: $\text{pile}(c)=p \wedge \text{at}(p, d) \wedge \text{at}(p', d)$
 $\wedge p \neq p' \wedge \text{cargo}(r)=\text{nil}$

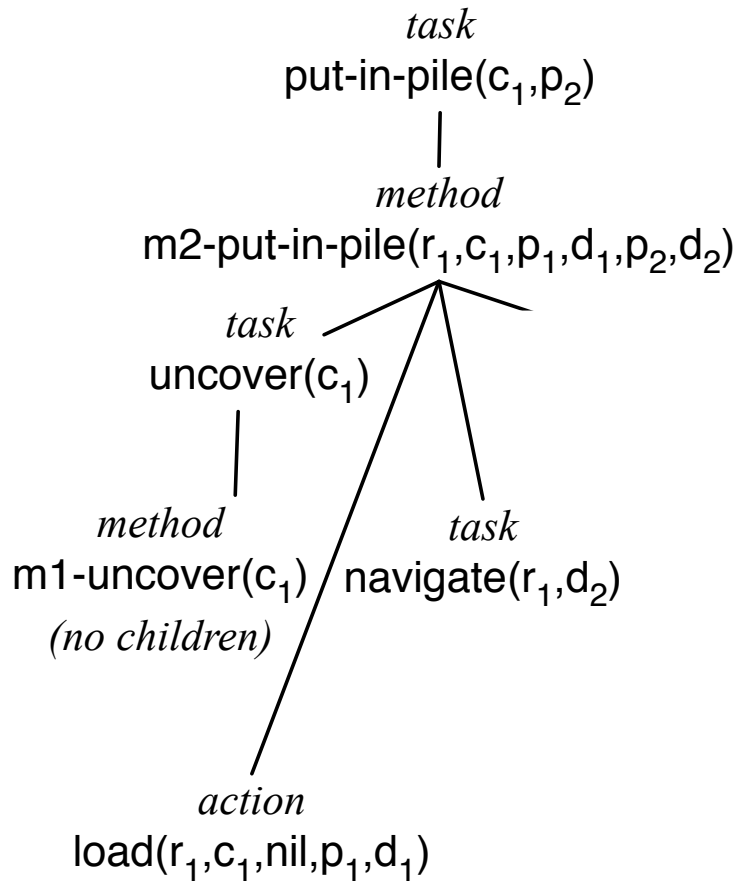
body: if $\text{loc}(r) \neq d$ then navigate(r, d)
 uncover(c)

load($r, c, \text{pos}(c), p, d$)
 if $\text{loc}(r) \neq d'$ then
 navigate(r, d')

unload($r, c, \text{top}(p'), p', d$)



Example

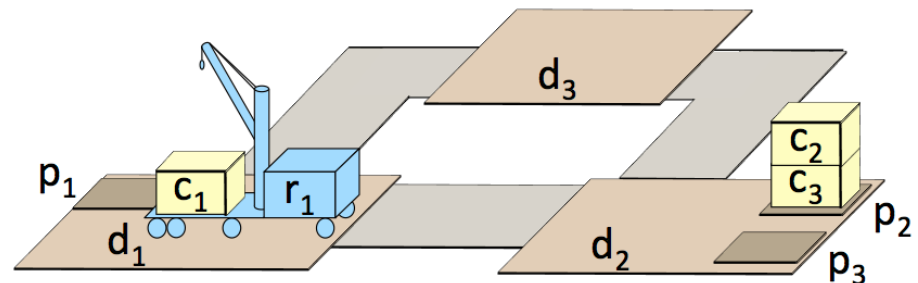


Candidates = { ...,
m2-navigate(r₁, d₂), ...,
 m3-navigate(r₁, d₃, d₂), ... }

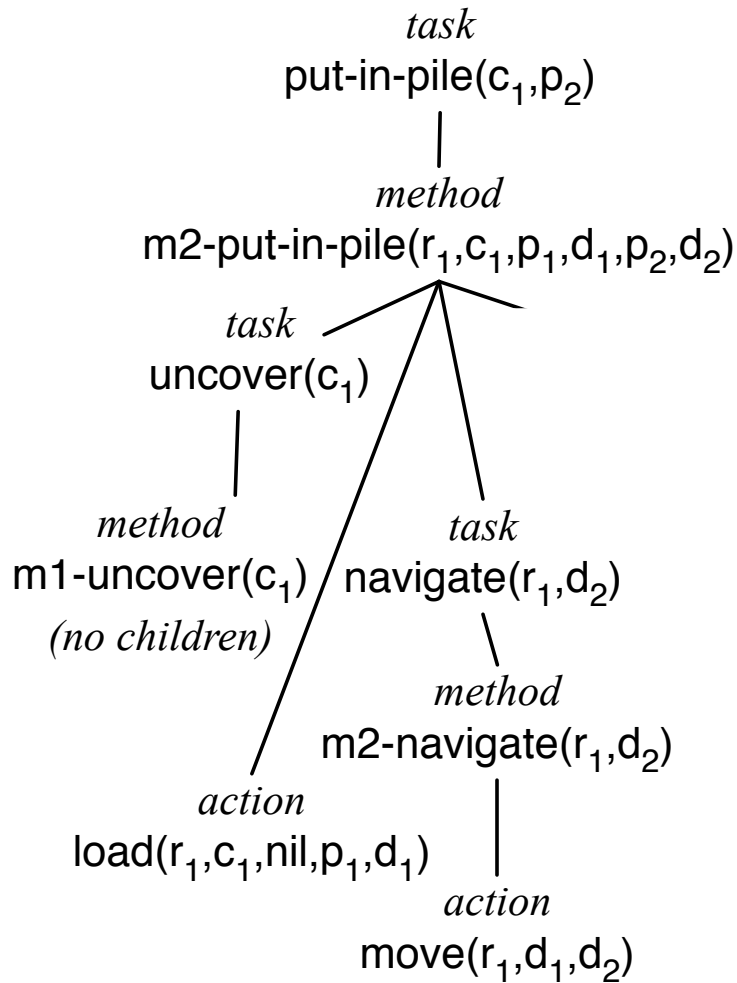
m1-navigate(r, d')
 task: navigate(r, d')
 pre: loc(r) = d'
 body: // empty

m2-navigate(r, d') **r₁, d₂**
 task: navigate(r, d')
 pre: loc(r) ≠ d' ∧ adjacent(loc(r), d')
 body: move(r, loc(r), d')

m3-navigate(r, d, d')
 task: navigate(r, d')
 pre: loc(r) ≠ d' ∧ d ≠ d'
 ∧ adjacent(loc(r), d)
 body: move(r, loc(r), d)
 navigate(r, d')



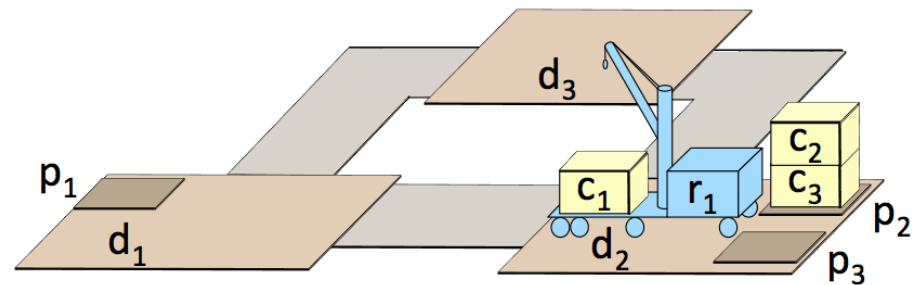
Example



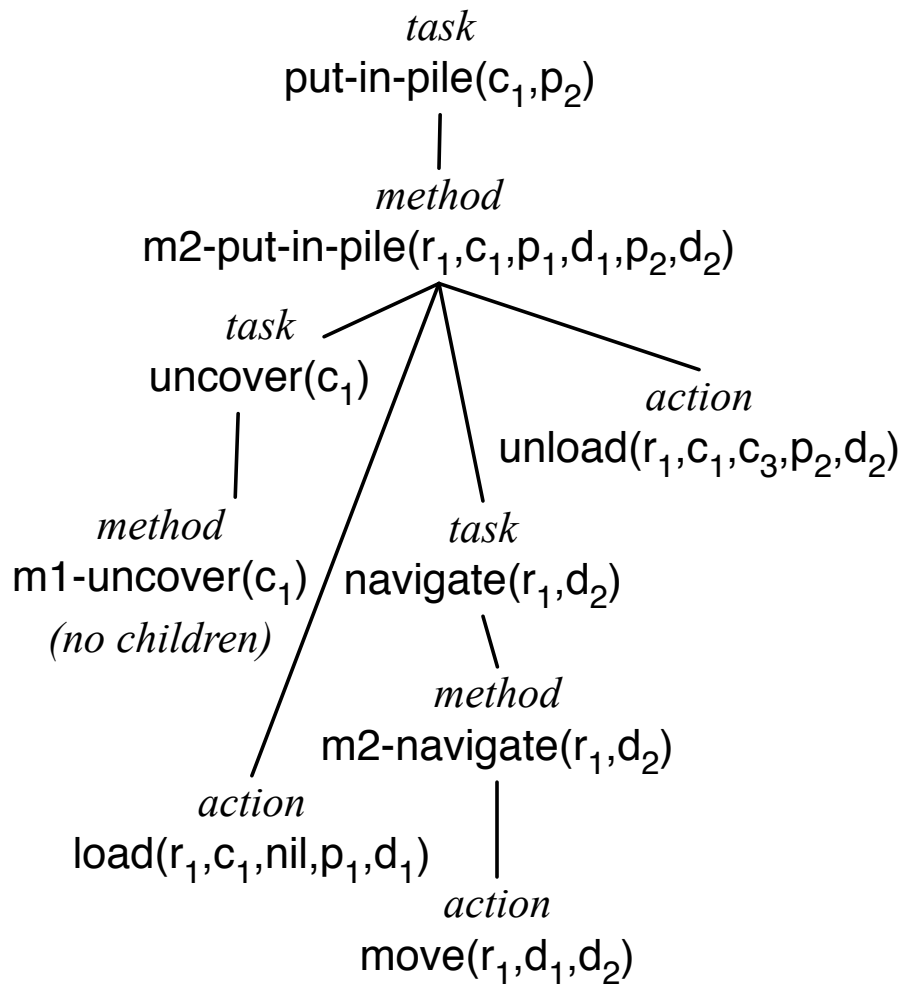
m1-navigate(r, d')
 task: navigate(r, d')
 pre: $\text{loc}(r) = d'$
 body: // empty

m2-navigate(r, d') r_1, d_2
 task: navigate(r, d')
 pre: $\text{loc}(r) \neq d' \wedge \text{adjacent}(\text{loc}(r), d')$
 body: $\text{move}(r, \text{loc}(r), d')$

m3-navigate(r, d, d')
 task: navigate(r, d')
 pre: $\text{loc}(r) \neq d' \wedge d \neq d'$
 $\wedge \text{adjacent}(\text{loc}(r), d)$
 body: $\text{move}(r, \text{loc}(r), d)$
 navigate(r, d')



Example



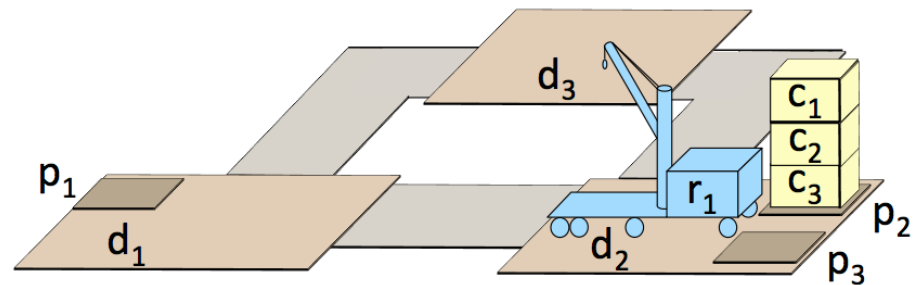
$r_1, c_1, p_1, d_1, p_2, d_2$

m2-put-in-pile(r, c, p, d, p', d')

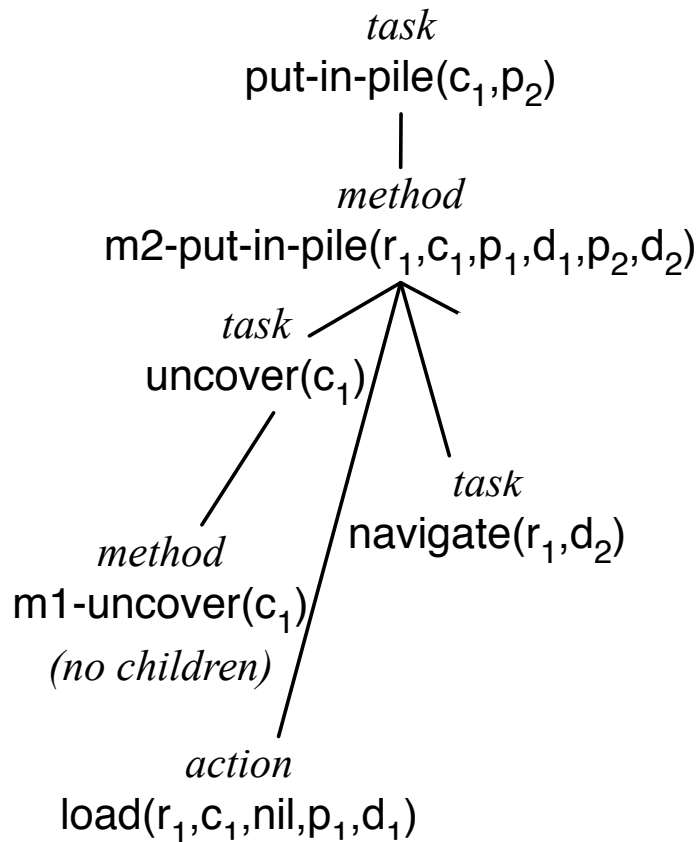
task: put-in-pile(c, p')

pre: $\text{pile}(c)=p \wedge \text{at}(p, d) \wedge \text{at}(p', d)$
 $\wedge p \neq p' \wedge \text{cargo}(r)=\text{nil}$

body: if $\text{loc}(r) \neq d$ then navigate(r, d)
 uncover(c)
 load($r, c, \text{pos}(c), p, d$)
 if $\text{loc}(r) \neq d'$ then
 navigate(r, d')
 unload($r, c, \text{top}(p'), p', d$)



Example

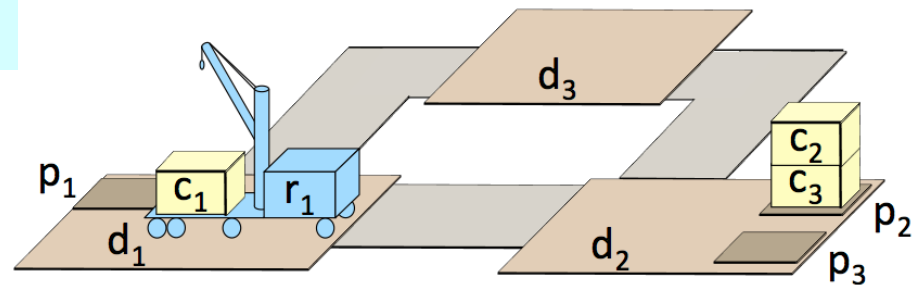


Candidates = { ...,
m2-navigate(r_1, d_2), ...,
m3-navigate(r_1, d_3, d_2), ... }

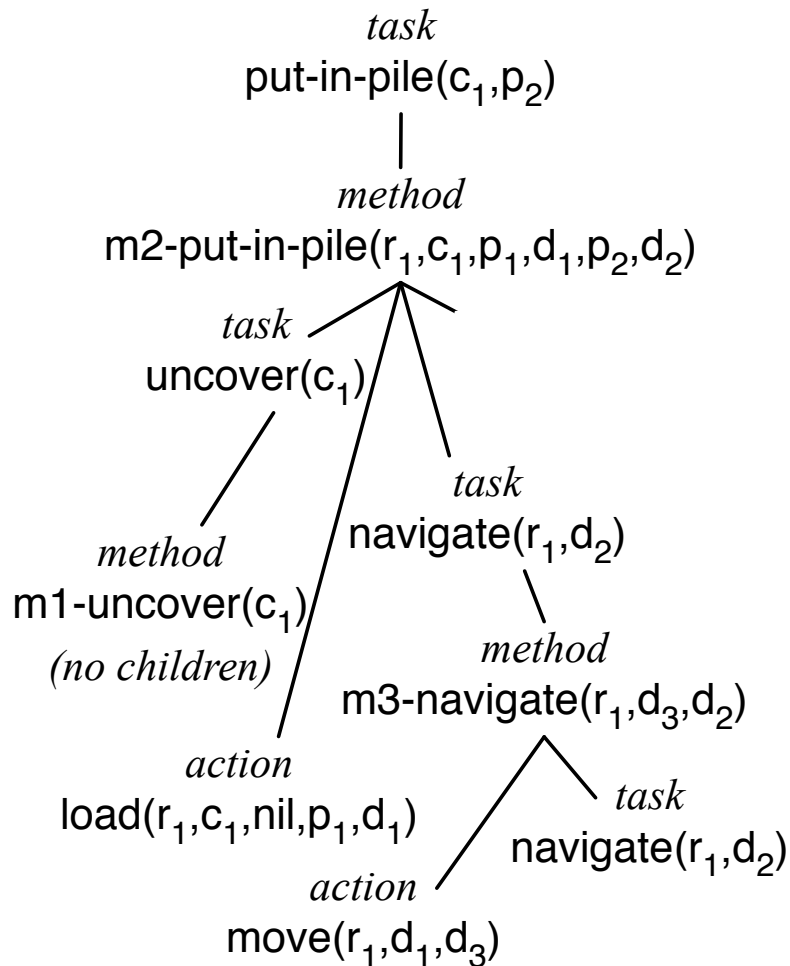
m1-navigate(r, d')
task: navigate(r, d')
pre: $\text{loc}(r) = d'$
body: // empty

m2-navigate(r, d')
task: navigate(r, d')
pre: $\text{loc}(r) \neq d' \wedge \text{adjacent}(\text{loc}(r), d')$
body: move($r, \text{loc}(r), d'$)

m3-navigate(r, d, d') **r_1, d_3, d_2**
task: navigate(r, d')
pre: $\text{loc}(r) \neq d' \wedge d \neq d'$
 $\wedge \text{adjacent}(\text{loc}(r), d)$
body: move($r, \text{loc}(r), d$)
navigate(r, d')



Example

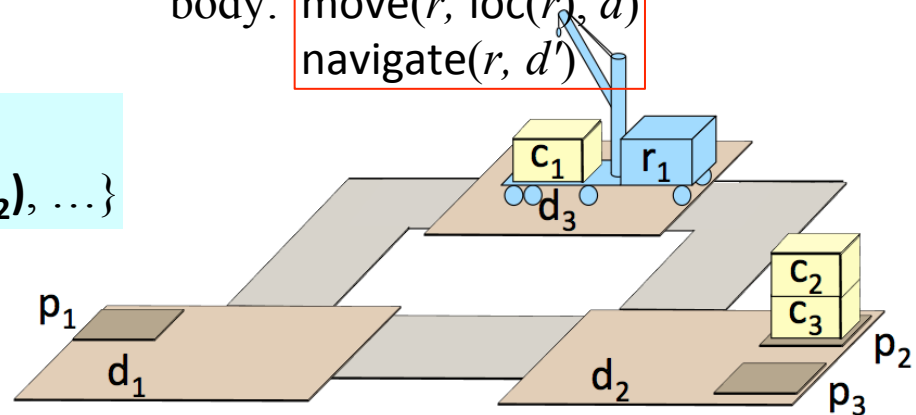


Candidates = { ...,
m2-navigate(r1, d2), ... }

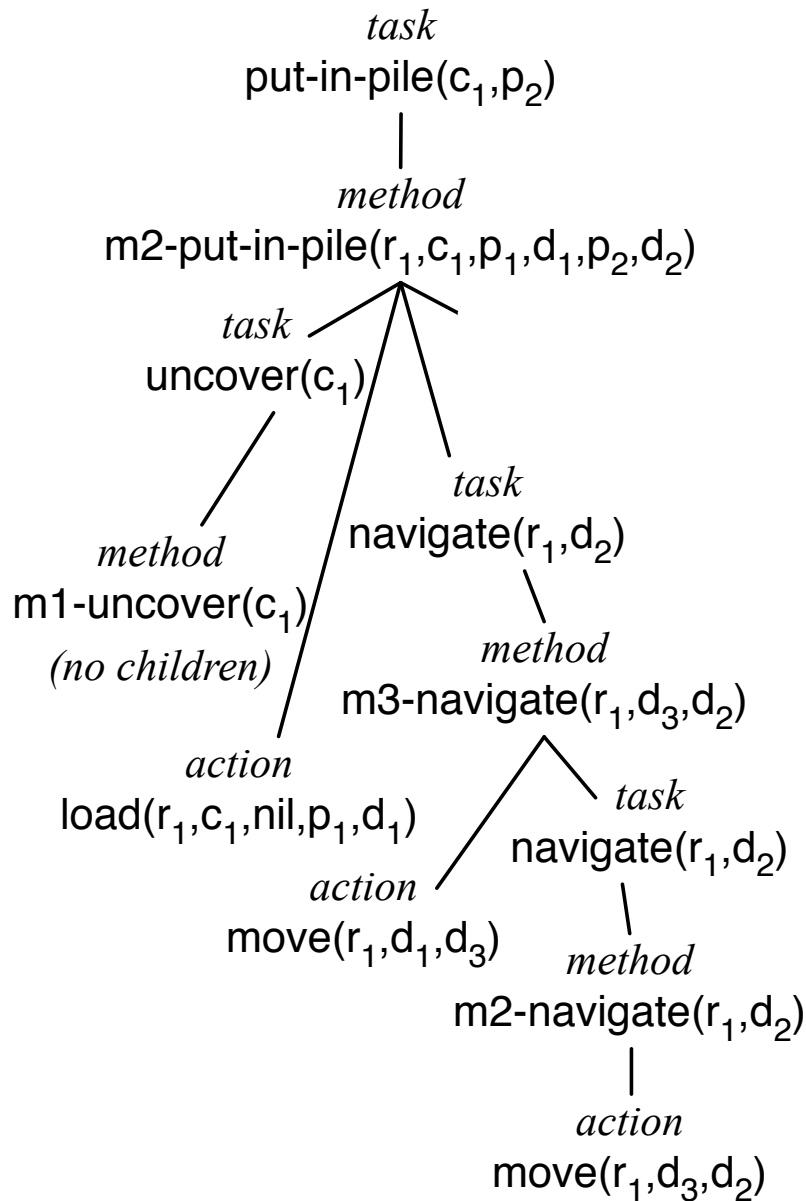
m1-navigate(r, d')
 task: navigate(r, d')
 pre: $\text{loc}(r) = d'$
 body: // empty

m2-navigate(r, d')
 task: navigate(r, d')
 pre: $\text{loc}(r) \neq d' \wedge \text{adjacent}(\text{loc}(r), d')$
 body: move($r, \text{loc}(r), d'$)

m3-navigate(r, d, d') r_1, d_3, d_2
 task: navigate(r, d')
 pre: $\text{loc}(r) \neq d' \wedge d \neq d'$
 $\wedge \text{adjacent}(\text{loc}(r), d)$
 body: move($r, \text{loc}(r), d$)
 navigate(r, d')



Example



m1-navigate(r, d')

task: navigate(r, d')

pre: $\text{loc}(r) = d'$

body: *// empty*

m2-navigate(r, d') r_1, d_2

task: navigate(r, d')

pre: $\text{loc}(r) \neq d' \wedge \text{adjacent}(\text{loc}(r), d')$

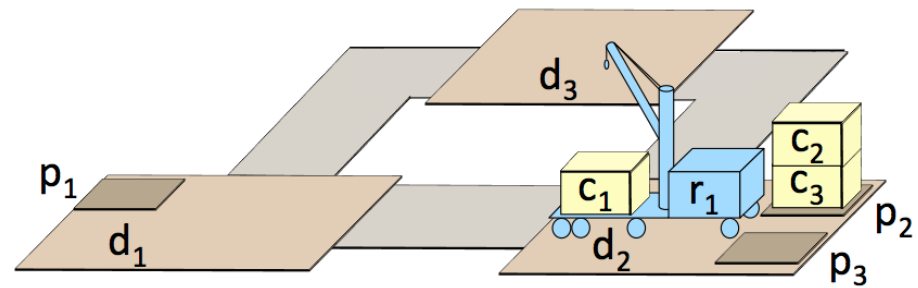
body: $\text{move}(r, \text{loc}(r), d')$

m3-navigate(r, d, d')

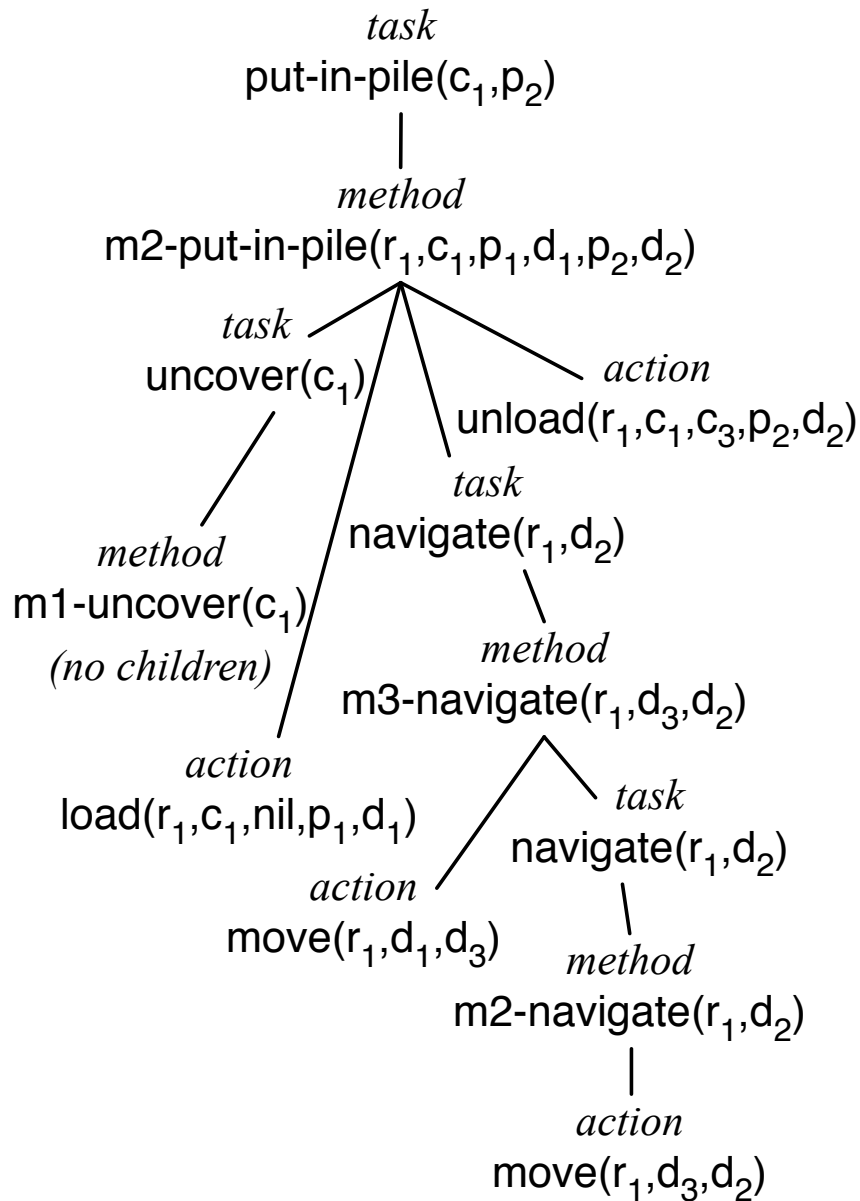
task: navigate(r, d')

pre: $\text{loc}(r) \neq d' \wedge d \neq d'$
 $\wedge \text{adjacent}(\text{loc}(r), d)$

body: $\text{move}(r, \text{loc}(r), d)$
 $\text{navigate}(r, d')$



Example



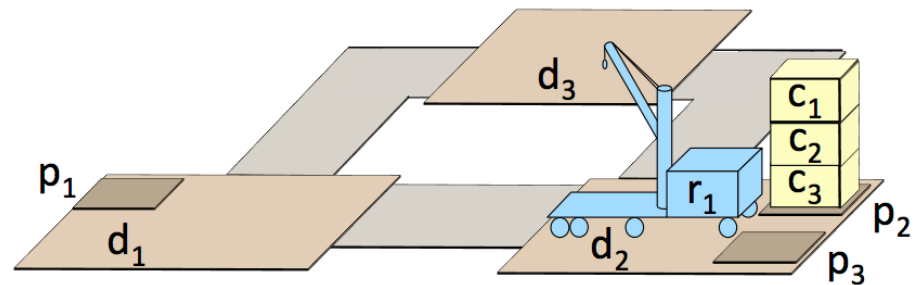
$r_1, c_1, p_1, d_1, p_2, d_2$

m2-put-in-pile(r, c, p, d, p', d')

task: put-in-pile(c, p')

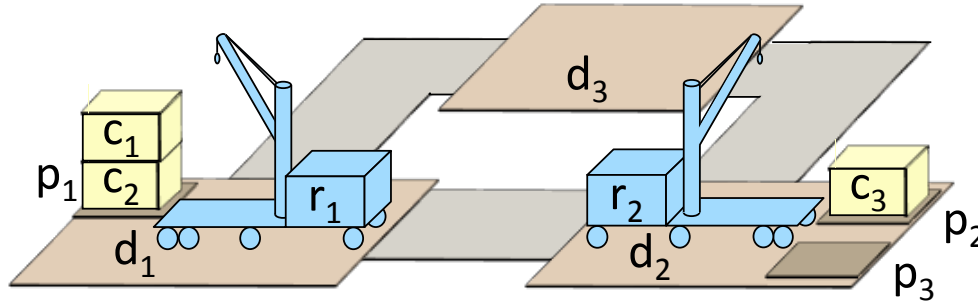
pre: $\text{pile}(c)=p \wedge \text{at}(p, d) \wedge \text{at}(p', d)$
 $\wedge p \neq p' \wedge \text{cargo}(r)=\text{nil}$

body: if $\text{loc}(r) \neq d$ then navigate(r, d)
 uncover(c)
 load($r, c, \text{pos}(c), p, d$)
 if $\text{loc}(r) \neq d'$ then
 navigate(r, d')
 unload($r, c, \text{top}(p'), p', d$)



Interleaved Refinement Tree (IRT) Procedure

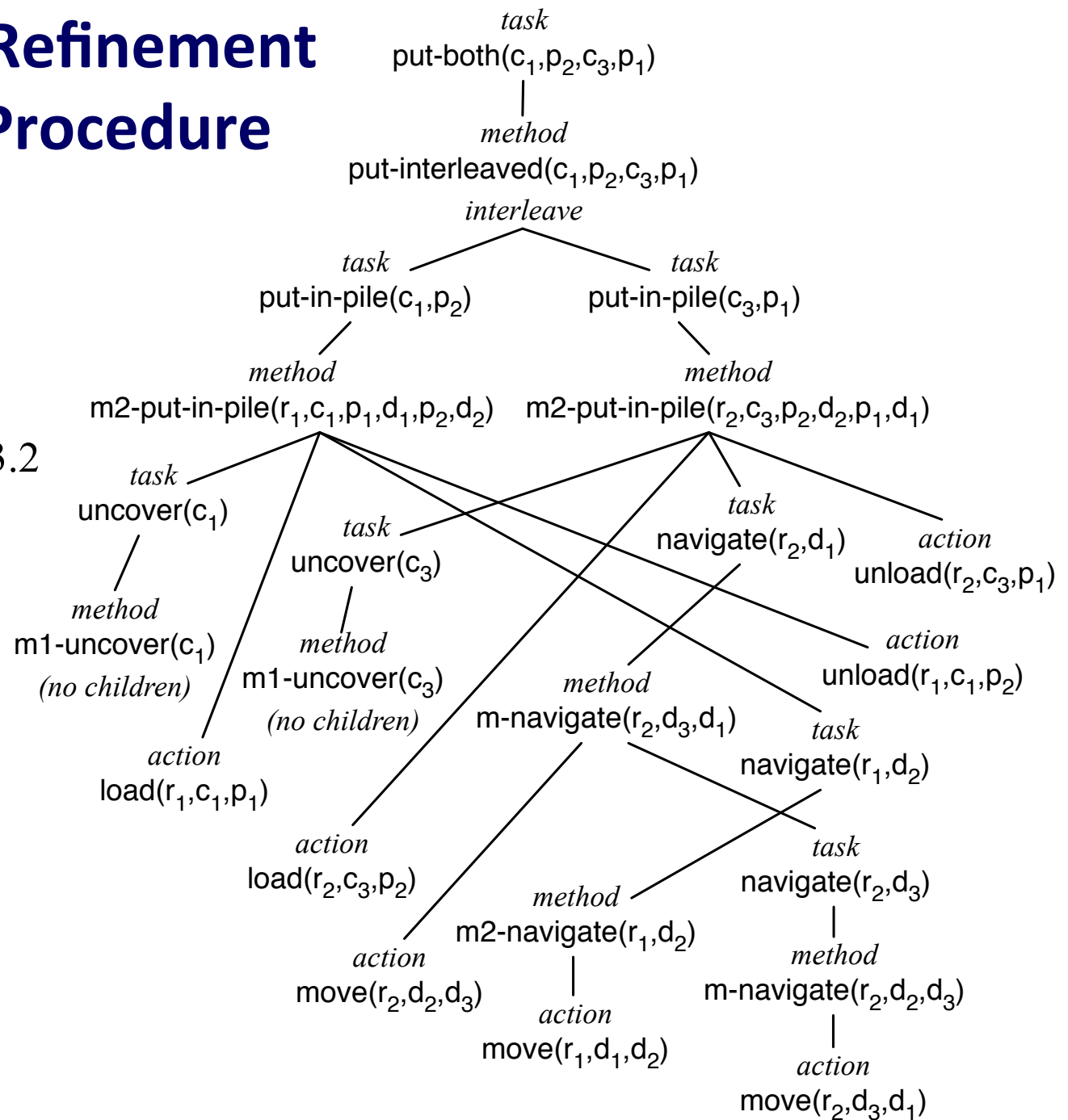
- SeRPE doesn't allow the 'concurrent' programming construct



- Assume same initial state as before
- Want to move c_1 to p_2 , and c_3 to p_1 using these plans:
 $\langle \text{load}(r_1, c_1, c_2, p_1, d_1), \text{move}(r_1, d_1, d_2), \text{unload}(r_1, c_1, p_3, \text{nil}, d_2) \rangle$
 $\langle \text{load}(r_2, c_3, \text{nil}, p_2, d_2), \text{move}(r_2, d_2, d_3), \text{move}(r_2, d_3, d_1), \text{unload}(r_2, c_3, c_2, p_1, d_1) \rangle$
- Need to interleave, otherwise they'll fail

Interleaved Refinement Tree (IRT) Procedure

- IRT extends SeRPE to interleave plans for different tasks
- Details: Section 3.3.2



3.4 Acting and Refinement Planning

- (to be added)