



Example of Naive (1)

$f_0 = \emptyset$

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

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Example of Naïve (2)

$f_1 = I_p(f_0) =$

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

$Local0(r, node1).$
 $Local0(s, node1).$
 $Local0(q, node1).$
 $Local0(u, node2).$
 $Op1(s1, r, node1).$
 $Op1(s2, u, node2).$
 $Op2(j1, s1, u1, node1).$
 $Op2(u1, s, q, node1).$
 $Op2(j2, j1, s2, node1).$

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Example of Naïve (3)

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

```

f2 = Ip(f1) =
  Local(r, node1, 1).
  Local(s, node1, 1).
  Local(q, node1, 1).
  Local(u, node2, 1).
  Local0(r, node1).
  Local0(s, node1).
  Local0(q, node1).
  Local0(u, node2).
  Op1(s1, r, node1).
  Op1(s2, u, node2).
  Op2(j1, s1, u1, node1).
  Op2(u1, s, q, node1).
  Op2(j2, j1, s2, node1).

```



Example of Naïve (4)

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

```

f3 = Ip(f2) =
  Local(r, node1, 1).
  Local(s, node1, 1).
  Local(q, node1, 1).
  Local(u, node2, 1).
  Local(s1, node1, 2).
  Local(u1, node1, 3).
  Local(s2, node2, 2).
  Local0(r, node1).
  Local0(s, node1).
  Local0(q, node1).
  Local0(u, node2).
  Op1(s1, r, node1).
  Op1(s2, u, node2).
  Op2(j1, s1, u1, node1).
  Op2(u1, s, q, node1).
  Op2(j2, j1, s2, node1).

```



Example of Naïve (5)

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

```

f4 = Ip(f3) =
Local(r, node1, 1).
Local(s, node1, 1).
Local(q, node1, 1).
Local(u, node2, 1).
Local(s1, node1, 2).
Local(u1, node1, 3).
Local(j1, node1, 6).
Local(s2, node2, 2).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

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Example of Naïve (6)

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

```

f5 = Ip(f4) = f4
Local(r, node1, 1).
Local(s, node1, 1).
Local(q, node1, 1).
Local(u, node2, 1).
Local(s1, node1, 2).
Local(u1, node1, 3).
Local(j1, node1, 6).
Local(s2, node2, 2).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

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Problem with Computed Predicates

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).

Local0(r, node1).
Op1(s1, r, node1).
Op1(s1, s1, node1).

f1 = Ip(∅)
Local0(r, node1).
Op1(s1, r, node1).
Op1(s1, s1, node1).
-----
f2 = Ip(f1)
Local(r, node1, 1).
Local0(r, node1).
Op1(s1, r, node1).
Op1(s1, s1, node1).
-----
f3 = Ip(f2)
Local(r, node1, 1).
Local(s1, node1, 3).
Local0(r, node1).
Op1(s1, r, node1).
Op1(s1, s1, node1).

```

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Some Refinements of Naive

- Start with $f = \text{facts}(P)$
- Use dependency graph
 - Nodes are the different predicates of P
 - Have an edge $p \rightarrow q$ if there is a rule of the form

$$q(\dots) \text{ :- } \dots, p(), \dots .$$

The paper uses a more detailed version of the dependency graph than the book, where there are nodes for rules as well.

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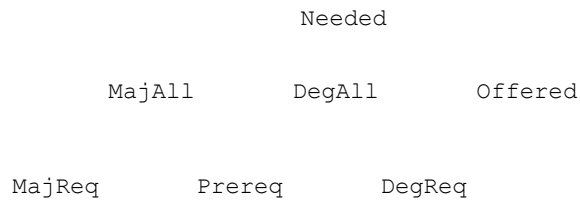


Example Dependency Graph

```
Needed(D, M, C) :- MajAll(M, C), Offered(D, M).
Needed(D, M, C) :- DegAll(D, C), Offered(D, M).
```

```
MajAll(M, C) :- MajReq(M, C).
MajAll(M, C) :- MajAll(M, C1), Prereq(C, C1).
```

```
DegAll(D, C) :- DegReq(D, C).
DegAll(D, C) :- DegAll(D, C1), Prereq(C, C1).
```



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Evaluate in Order of Dependency Graph

- **Compute** DegAll from DegReq and Prereq
- **Compute** MajorAll from MajReq and Prereq
- **Compute** Needed from MajAll, DegAll, Offered

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Semi-Naïve Evaluation

Naïve evaluation discovers the same fact over and over.

E.g., derive `local(r, node1, 1)` repeatedly

Idea: When evaluating a rule, make sure we use at least one newly discovered fact.

Otherwise, we derive a fact we've already seen



Example

Consider the rule

```
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
```

When computing $f_4 = I_p(f_3)$.

If we use facts `Op2(u1, s, q, node1)`,
`Local(s, node1, 1)`, `Local(q, node1, 1)`, we
 derive `Local(u1, node1, 3)`, which we already
 have

If we use `Op2(j1, s1, u1, node1)`,
`Local(s1, node1, 2)`, `Local(u1, node1, 3)`, we
 get



Implementing Semi-Naïve

Notice after f_1 , only new facts are `Local` facts

Suppose at each step in Naïve we determine

`NewLocal(Id, Loc, K)` if `Local(Id, Loc, K)` in $f - f_{old}$.

We can work with a modified set of rules

```
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    NewLocal(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    NewLocal(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), NewLocal(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
```



One Other Modification

If P' is the modified program, then instead of

$$f = I_p(f)$$

we need

$$f = I_{p'}(f) \cup f.$$

Note: Semi-Naïve doesn't guarantee we won't generate a fact at multiple steps, if there are distinct "proofs" of it

```
MajReq(eng, wr310). MajReq(eng, wr408).
Prereq(wr309, wr310).
Prereq(wr401, wr408).
Prereq(wr309, wr401).
```



Directed Methods

Suppose we have the query

```
:- Local(I, node2, K).
```

We could push the constant into the rules:

```
Local(Id, node2, 1) :- Local0(Id, node2).
```

```
Local(Id, node2, K) :- Op1(Id, In, node2),  
    Local(In, node2, M), Sum(M, 1, K).
```

```
Local(Id, node2, K) :- Op2(Id, In1, In2, node2),  
    Local(In1, node2, M), Local(In2, node2, N),  
    Sum(M, N, J), Sum(J, 1, K).
```



Top-Down Evaluation

What Prolog uses

- Start with the first goal in the list
- If you find a matching rule head, replace that goal by the rule body (with appropriate bindings)
- Continue with the next goal.
- Might have to backtrack if you get stuck or want more answers.



Top-Down Example

Initial goal

```
:- Local(u1, L, C).
```

Try different rules

```
Local(Id, Loc, 1) :- Local0(Id, Loc).
```

```
Local(Id, Loc, K) :- Op1(Id, In, Loc),  
    Local(In, Loc, M), Sum(M, 1, K).
```

```
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),  
    Local(In1, Loc, M), Local(In2, Loc, N),  
    Sum(M, N, J), Sum(J, 1, K).
```



Have a New Goal (List)

```
:- Op2(u1, In1, In2, Loc),  
    Local(In1, Loc, M), Local(In2, Loc, N),  
    Sum(M, N, J), Sum(J, 1, C).
```

Try Op2 facts

```
Op2(j1, s1, u1, node1).
```

```
Op2(u1, s, q, node1).
```



Next Goal List

```
:- Local(s, node1, M), Local(q, node1, N),
   Sum(M, N, J), Sum(J, 1, C).
```

Try matching first goal

```
Local(Id, Loc, 1) :- Local0(Id, Loc).
```

Which gives

```
:- Local0(s, node1), Local(q, node1, N),
   Sum(1, N, J), Sum(J, 1, C).
```

Try Local0 facts

```
Local0(r, node1).
```

```
Local0(s, node1).
```



Continuing

```
:- Local(q, node1, N),
   Sum(1, N, J), Sum(J, 1, C).
```

Try matching first goal

```
Local(Id, Loc, 1) :- Local0(Id, Loc).
```

Which gives

```
:- Local0(q, node1), Sum(1, 1, J), Sum(J, 1, C).
```

Try Local0 facts

```
Local0(r, node1).
```

```
Local0(s, node1).
```

```
Local0(q, node1).
```



Finishing up

```
:- Sum(1, 1, J), Sum(J, 1, C).
```

```
Sum(1, 1, 2).
```

```
:- Sum(2, 1, C).
```

```
Sum(2, 1, 3).
```

Success!



Top-Down Methods

Sometimes you can propagate constants statically (at “compile” time), before any computation.

However, not all constants are available at compile time.



Example of Constant Propagation

Consider query

```
:- Needed(ba, eng, C).
```

for our program

```
Needed(D, M, C) :- MajAll(M, C), Offered(D, M).
```

```
Needed(D, M, C) :- DegAll(D, C), Offered(D, M).
```

```
MajAll(M, C) :- MajReq(M, C).
```

```
MajAll(M, C) :- MajAll(M, C1), Prereq(C, C1).
```

```
DegAll(D, C) :- DegReq(D, C).
```

```
DegAll(D, C) :- DegAll(D, C1), Prereq(C, C1).
```



Can Do Some Constant Propagation Right Away

Modified program

```
Needed(ba, eng, C) :- MajAll(eng, C), Offered(ba, eng).
```

```
Needed(ba, eng, C) :- DegAll(ba, C), Offered(ba, eng).
```

```
MajAll(eng, C) :- MajReq(eng, C).
```

```
MajAll(eng, C) :- MajAll(eng, C1), Prereq(C, C1).
```

```
DegAll(ba, C) :- DegReq(ba, C).
```

```
DegAll(ba, C) :- DegAll(ba, C1), Prereq(C, C1).
```



Consider Evaluation of MajAll

Suppose we have the facts

`MajReq(eng, wr310).`

`MajReq(eng, lit481).`

We'll derive (by the first rule)

`MajAll(eng, wr310).`

`MajAll(eng, lit481).`

Specializing the second rule:

`MajAll(eng, C) :- MajAll(eng, wr310), Prereq(C, wr310).`

`MajAll(eng, C) :- MajAll(eng, lit481),
Prereq(C, lit481).`



New Sub-queries

Pose the new queries

`:- Prereq(C, wr310).`

`:- Prereq(C, lit481).`

which can be used to instantiate the rules
together with the MajAll facts.

The Iterative Query-Sub-Query (QSQ) Method

Combine queries with the same *query form* (pattern of free and bound positions) into a *generalized query*:

```
:- Prereq(C, {wr310, lit481}).
```

Use generalized query:

- Evaluate against extensional DB (if appropriate)
- Plug results into rules for new facts
- Use new facts to generate new sub-queries
- Combine sub-queries with existing generalized query, if possible

Example

Suppose we have the facts:

```
Prereq(wr301, wr310).
Prereq(wr301, wr418).
Prereq(wr200, wr301).
Prereq(lit480, lit481).
Prereq(lit217, lit218).
```

Which match the generalized query?

```
:- Prereq(C, {wr310, lit481}).
```



Example (2)

Plug

```
Prereq(wr301, wr310).
Prereq(lit480, lit481).
```

plus

```
MajAll(eng, wr310).
MajAll(eng, lit481).
```

into

```
MajAll(eng, C) :- MajAll(eng, C1), Prereq(C, C1).
```

to get

```
MajAll(eng, wr301).
MajAll(eng, lit480).
```

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Example (3)

Use

```
MajAll(eng, wr301).
MajAll(eng, lit480).
```

with

```
MajAll(eng, C) :- MajAll(eng, C1), Prereq(C, C1).
```

to get new sub-queries

```
:- Prereq(C, wr301).
:- Prereq(C, lit480).
```

Expand generalized query

```
:- Prereq(C, {wr310, lit481, wr301, lit480}).
```

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Example (4)

Evaluate generalized query again

```
Prereq(wr301, wr310).  
Prereq(wr301, wr418).  
Prereq(wr200, wr301).  
Prereq(lit480, lit481).  
Prereg(lit217, lit218).
```

Which match the new generalized query?

```
:- Prereq(C, {wr310, lit481, wr301, lit480})
```



Example (5)

Plug

```
Prereq(wr301, wr310).  
Prereq(wr201, wr301).  
Prereq(lit480, lit481).
```

plus

```
MajAll(eng, wr310).  
MajAll(eng, lit481).  
MajAll(eng, wr301).  
MajAll(eng, lit480).
```

into

```
MajAll(eng, C) :- MajAll(eng, C1), Prereq(C, C1).
```



Example (6)

We get

```
MajAll(eng, wr301).  
MajAll(eng, lit480).  
MajAll(eng, wr200).
```

Plus we expand the generalized query to

```
:- Prereq(C,  
    {wr310, lit481, wr301, lit480, wr200}).
```



Example (7)

Evaluate generalized query again

```
Prereq(wr301, wr310).  
Prereq(wr301, wr418).  
Prereq(wr200, wr301).  
Prereq(lit480, lit481).  
Prereq(lit217, lit218).
```

Which match the new generalized query?

```
:- Prereq(C,  
    {wr310, lit481, wr301, lit480, wr200}).
```



Example (8)

The answers we get are the same for the last version of the generalized query:

```
Prereq(wr301, wr310) .  
Prereq(wr201, wr301) .  
Prereq(lit480, lit481) .
```

So we will get no new facts nor new sub-queries.