



Example of Naive (1)

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

 $f_0 = \emptyset$

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Example of Naïve (2)

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

 $f_1 = I_p(f_0) =$

```

Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

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Example of Naïve (3)

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

f2 = Ip(f1) =
Local(r, node1, 1).
Local(s, node1, 1).
Local(q, node1, 1).
Local(u, node2, 1).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

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Example of Naïve (4)

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

f3 = Ip(f2) =
Local(r, node1, 1).
Local(s, node1, 1).
Local(q, node1, 1).
Local(u, node2, 1).
Local(s1, node1, 2).
Local(u1, node1, 3).
Local(s2, node2, 2).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

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Example of Naïve (5)

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

f4 = Ip(f3) =
Local(r, node1, 1).
Local(s, node1, 1).
Local(q, node1, 1).
Local(u, node2, 1).
Local(s1, node1, 2).
Local(u1, node1, 3).
Local(j1, node1, 6).
Local(s2, node2, 2).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

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Example of Naïve (6)

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

f5 = Ip(f4) = f4
Local(r, node1, 1).
Local(s, node1, 1).
Local(q, node1, 1).
Local(u, node2, 1).
Local(s1, node1, 2).
Local(u1, node1, 3).
Local(j1, node1, 6).
Local(s2, node2, 2).
Local0(r, node1).
Local0(s, node1).
Local0(q, node1).
Local0(u, node2).
Op1(s1, r, node1).
Op1(s2, u, node2).
Op2(j1, s1, u1, node1).
Op2(u1, s, q, node1).
Op2(j2, j1, s2, node1).

```

fixed point

:- Local(j1, L, C)

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Problem with Computed Predicates

```

Local(Id, Loc, 1) :- Local0(Id, Loc).
Local(Id, Loc, K) :- Op1(Id, In, Loc),
    Local(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
    Local(In1, Loc, M), Local(In2, Loc, N),
    Sum(M, N, J), Sum(J, 1, K).

Local0(r, node1).
Op1(s1, r, node1).
Op1(s1, s1, node1).

f1 = Ip(∅)
Local0(r, node1).
Op1(s1, r, node1).
Op1(s1, s1, node1).
-----
f2 = Ip(f1)
Local(r, node1, 1).
Local0(r, node1).
Op1(s1, r, node1).
Op1(s1, s1, node1).
-----
f3 = Ip(f2)
Local(r, node1, 1).
Local(s1, node1, 3).
Local0(r, node1).
Op1(s1, r, node1).
Op1(s1, s1, node1).

```

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Some Refinements of Naive

- Start with $f = \text{facts}(P)$
- Use dependency graph
 - Nodes are the different predicates of P
 - Have an edge $p \rightarrow q$ if there is a rule of the form

$$q(\dots) \text{ :- } \dots, p(), \dots$$

The paper uses a more detailed version of the dependency graph than the book, where there are nodes for rules as well.

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 *degree major course*
Example Dependency Graph

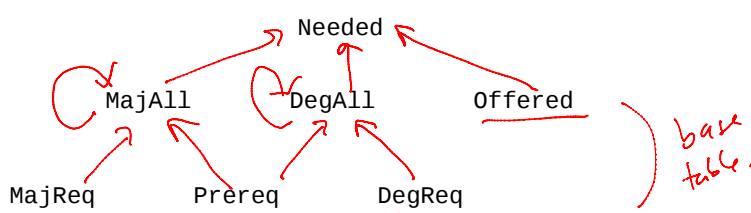
```

Needed(D, M, C) :- MajAll(M, C), Offered(D, M).
Needed(D, M, C) :- DegAll(D, C), Offered(D, M).

MajAll(M, C) :- MajReq(M, C).
MajAll(M, C) :- MajAll(M, C1), Prereq(C, C1).

DegAll(D, C) :- DegReq(D, C).
DegAll(D, C) :- DegAll(D, C1), Prereq(C, C1).

```



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 **Evaluate in Order of Dependency Graph**

- Compute DegAll from DegReq and Prereq
- Compute MajorAll from MajReq and Prereq
- Compute Needed from MajAll, DegAll, Offered

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Semi-Naïve Evaluation

Naïve evaluation discovers the same fact over and over.

E.g., derive `local(r, node1, 1)` repeatedly

Idea: When evaluating a rule, make sure we use at least one newly discovered fact.

Otherwise, we derive a fact we've already seen



Example

Consider the rule

```
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
  Local(In1, Loc, M), Local(In2, Loc, N),
  Sum(M, N, J), Sum(J, 1, K).
```

When computing $f_4 = I_p(f_3)$.

If we use facts `Op2(u1, s, q, node1)`,
`Local(s, node1, 1)`, `Local(q, node1, 1)`, we
 derive `Local(u1, node1, 3)`, which we already
 have *in f3*

If we use `Op2(j1, s1, u1, node1)`,
`Local(s1, node1, 2)`, `Local(u1, node1, 3)`, we
 get *Local(j1, node1, 6)*.

Implementing Semi-Naïve

Notice after f_1 , only new facts are Local facts

Suppose at each step in Naïve we determine

$\text{NewLocal}(\text{Id}, \text{Loc}, K)$ if $\text{Local}(\text{Id}, \text{Loc}, K)$ in $f - f_{\text{old}}$.

We can work with a modified set of rules

2nd Rule

```

Local(Id, Loc, K) :- Op1(Id, In, Loc),
  NewLocal(In, Loc, M), Sum(M, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
  NewLocal(In1, Loc, M), Local(In2, Loc, N),
  Sum(M, N, J), Sum(J, 1, K).
Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
  Local(In1, Loc, M), NewLocal(In2, Loc, N),
  Sum(M, N, J), Sum(J, 1, K).
  
```

Local 2 New Local

3rd Rule

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One Other Modification

If P' is the modified program, then instead of

$f = I_P(f)$

we need

$f = I_{P'}(f) \cup f.$

f might not be contained in $I_{P'}(f)$

Note: Semi-Naïve doesn't guarantee we won't generate a fact at multiple steps, if there are distinct "proofs" of it

$\text{MajReq}(\text{eng}, \text{wr310}). \text{MajReq}(\text{eng}, \text{wr408}).$

$\text{Prereq}(\text{wr309}, \text{wr310}).$

$\text{Prereq}(\text{wr401}, \text{wr408}).$

$\text{Prereq}(\text{wr309}, \text{wr401}).$

MajAll (eng, wr309) f2
MajAll (eng, wr408) f3
MajAll (eng, wr309) f3

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Directed Methods

Suppose we have the query $\text{:- local}(j1, L, K).$
 $\text{:- Local(I, node2, K).}$

We could push the constant into the rules:

$\text{Local(Id, node2, 1) :- Local0(Id, node2).}$

$\text{Local(Id, node2, K) :- Op1(Id, In, node2),}$
 $\text{Local(In, node2, M), Sum(M, 1, K).}$

$\text{Local(Id, node2, K) :- Op2(Id, In1, In2, node2),}$
 $\text{Local(In1, node2, M), Local(In2, node2, N),}$
 $\text{Sum(M, N, J), Sum(J, 1, K).}$

Top-Down Evaluation

What Prolog uses

- Start with the first goal in the list
- If you find a matching rule head, replace that goal by the rule body (with appropriate bindings)
- Continue with the next goal.
- Might have to backtrack if you get stuck or want more answers.

success if no goal goals left
 look at bindings of variables
 in the initial goal to get
 answers.



Top-Down Example

Initial goal

`:- Local(u1, L, C).`

Try different rules

`Local(Id, Loc, 1) :- Local0(Id, Loc).`

`Local(Id, Loc, K) :- Op1(Id, In, Loc),
Local(In, Loc, M), Sum(M, 1, K).`

`Local(Id, Loc, K) :- Op2(Id, In1, In2, Loc),
Local(In1, Loc, M), Local(In2, Loc, N),
Sum(M, N, J), Sum(J, 1, K).`

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Have a New Goal (List)

`:- Op2(u1, In1, In2, Loc),
Local(In1, Loc, M), Local(In2, Loc, N),
Sum(M, N, J), Sum(J, 1, C).`

Try Op2 facts

`Op2(j1, s1, u1, node1).`

`Op2(u1, s, q, node1).`

*In1 = s In2 = q
L = node1*

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Next Goal List

$M=1$
`:- Local(s, node1, M), Local(q, node1, N),
 Sum(M, N, J), Sum(J, 1, C).`

Try matching first goal

s $node1$ $M=1$
`Local( $s$ ,  $node1$ , 1) :- Local0( $s$ ,  $node1$ ).`

Which gives

s $node1$
`:- Local0(s, node1), Local(q, node1, N),
 Sum(1, N, J), Sum(J, 1, C).`

Try Local0 facts

`Local0(r, node1).` ✗

`Local0(s, node1).` ✓

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Continuing

$N=1$
`:- Local(q, node1, N),
 Sum(1, N, J), Sum(J, 1, C).`

Try matching first goal

q $node1$
`Local( $q$ ,  $node1$ , 1) :- Local0( $q$ ,  $node1$ ).`

Which gives

`:- Local0(q, node1), Sum(1, 1, J), Sum(J, 1, C).`

Try Local0 facts

`Local0(r, node1).` ✗

`Local0(s, node1).` ✗

`Local0(q, node1).` ✓

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Finishing up

```

:- Sum(1, 1, J), Sum(J, 1, C).
Sum(1, 1, 2).

:- Sum(2, 1, C).
Sum(2, 1, 3).
    
```

Success!

Handwritten notes:
 $J = 2$
 $C = 3$
 $L = \text{node 1}$
 $\text{local}(\text{node 1}, \text{node 1}, 3)$

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Top-Down Methods

Sometimes you can propagate constants statically (at “compile” time), before any computation.

However, not all constants are available at compile time.

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Example of Constant Propagation

Consider query

```
:- Needed(ba, eng, C).
```

for our program

```
Needed(D, M, C) :- MajAll(M, C), Offered(D, M).
```

```
Needed(D, M, C) :- DegAll(D, C), Offered(D, M).
```

```
MajAll(M, C) :- MajReq(M, C).
```

```
MajAll(M, C) :- MajAll(M, C1), Prereq(C, C1).
```

```
DegAll(D, C) :- DegReq(D, C).
```

```
DegAll(D, C) :- DegAll(D, C1), Prereq(C, C1).
```



Can Do Some Constant Propagation Right Away

Modified program

```
Needed(ba, eng, C) :- MajAll(eng, C), Offered(ba, eng).
```

```
Needed(ba, eng, C) :- DegAll(ba, C), Offered(ba, eng).
```

```
MajAll(eng, C) :- MajReq(eng, C).
```

```
MajAll(eng, C) :- MajAll(eng, C1), Prereq(C, C1).
```

```
DegAll(ba, C) :- DegReq(ba, C).
```

```
DegAll(ba, C) :- DegAll(ba, C1), Prereq(C, C1).
```



Consider Evaluation of MajAll

Suppose we have the facts

`MajReq(eng, wr310).`

`MajReq(eng, lit481).`

We'll derive (by the first rule)

`MajAll(eng, wr310).`

`MajAll(eng, lit481).`

Specializing the second rule:

`MajAll(eng, C) :- MajAll(eng, wr310), Prereq(C, wr310).`

`MajAll(eng, C) :- MajAll(eng, lit481),
Prereq(C, lit481).`



New Sub-queries

Pose the new queries

`:- Prereq(C, wr310).`

`:- Prereq(C, lit481).`

which can be used to instantiate the rules
together with the MajAll facts.

The Iterative Query-Sub-Query (QSQ) Method

Combine queries with the same *query form* (pattern of free and bound positions) into a *generalized query*.

match any of these
`:- Prereq(C, {wr310, lit481}).`

Use generalized query:

- Evaluate against extensional DB (if appropriate)
- Plug results into rules for new facts
- Use new facts to generate new sub-queries
- Combine sub-queries with existing generalized query, if possible

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Example

Suppose we have the facts:

`Prereq(wr301, wr310).`
`Prereq(wr301, wr418).`
`Prereq(wr200, wr301).`
`Prereq(lit480, lit481).`
`Prereq(lit217, lit218).`

Which match the generalized query?

`:- Prereq(C, {wr310, lit481}).`

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Example (2)

Plug

```
Prereq(wr301, wr310).  
Prereq(lit480, lit481).
```

plus

```
MajAll(eng, wr310).  
MajAll(eng, lit481).
```

into

```
MajAll(eng, C) :- MajAll(eng, C1), Prereq(C, C1).
```

to get

```
MajAll(eng, wr301).  
MajAll(eng, lit480).
```

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Example (3)

Use

```
MajAll(eng, wr301).  
MajAll(eng, lit480).
```

with

```
MajAll(eng, C) :- MajAll(eng, C1), Prereq(C, C1).
```

to get new sub-queries

```
:- Prereq(C, wr301).  
:- Prereq(C, lit480).
```

Expand generalized query

```
:- Prereq(C, {wr310, lit481, wr301, lit480}).
```

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Example (4)

Evaluate generalized query again

```
Prereq(wr301, wr310).
Prereq(wr301, wr418).
Prereq(wr200, wr301).
Prereq(lit480, lit481).
Prereq(lit217, lit218).
```

Which match the new generalized query?

```
:- Prereq(C, {wr310, lit481, wr301, lit480})
```



Example (5)

Plug

```
Prereq(wr301, wr310).
Prereq(wr201, wr301).
Prereq(lit480, lit481).
```

plus

```
MajAll(eng, wr310).
MajAll(eng, lit481).
MajAll(eng, wr301).
MajAll(eng, lit480).
```

into

```
MajAll(eng, C) :- MajAll(eng, C1), Prereq(C, C1).
```



Example (6)

We get

```
MajAll(eng, wr301).  
MajAll(eng, lit480).  
MajAll(eng, wr200).
```

Plus we expand the generalized query to

```
:- Prereq(C,  
    {wr310, lit481, wr301, lit480, wr200}).
```



Example (7)

Evaluate generalized query again

```
Prereq(wr301, wr310).  
Prereq(wr301, wr418).  
Prereq(wr200, wr301).  
Prereq(lit480, lit481).  
Prereg(lit217, lit218).
```

Which match the new generalized query?

```
:- Prereq(C,  
    {wr310, lit481, wr301, lit480, wr200}).
```



Example (8)

The answers we get are the same for the last version of the generalized query:

```
Prereq(wr301, wr310).  
Prereq(wr201, wr301).  
Prereq(lit480, lit481).
```

So we will get no new facts nor new sub-queries.