

7-5

Review:

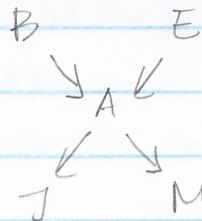
B = burglar

E = earthquake

A = alarm

J = John calls

M = Mary calls



* Belief network (BN)

= directed acyclic graph (DAG)

+ conditional probability tables (CPTs)

* conditional independence

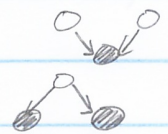
$$\begin{aligned}
 P(X_1, X_2, \dots, X_n) &= P(X_1) P(X_2 | X_1) \dots P(X_n | X_1, \dots, X_{n-1}) \quad \text{product rule} \\
 &= \prod_{i=1}^n P(X_i | X_1, \dots, X_{i-1}) \\
 &= \prod_{i=1}^n P(X_i | \text{pa}(X_i)) \quad \text{where } \text{pa}(X_i) = \text{parents of } X_i \text{ is subset of } \{X_1, X_2, \dots, X_{i-1}\}
 \end{aligned}$$

* Types of reasoning

1. competing explanations of observed event

2. multiple events with common explanation

3. intervening events



* Representing CPTs

$X_1 \quad X_2 \quad \dots \quad X_k$

Assume $X_i \in \{0, 1\}$

$Y \in \{0, 1\}$



How to represent CPT $P(Y=1 | X_1, X_2, \dots, X_k)$?

Option #1: look up table $O(2^k)$ can store arbitrary CPT

	X_1	X_2	$\dots$	X_k	$P(Y=1 X_1, X_2, \dots, X_k)$
2^k rows	0	0	$\dots$	0	0.6
	1	0	$\dots$	0	0.3
	$\vdots$	$\vdots$	$\dots$	$\vdots$	$\vdots$
	1	1	$\dots$	1	0.8

But what if k is very large?

Option #2: "deterministic" node

"AND" $P(Y=1 | X_1, X_2, \dots, X_k) = \prod_{i=1}^k x_i$

"OR" $P(Y=0 | X_1, X_2, \dots, X_k) = \prod_{i=1}^k (1-x_i)$

Option #3: noisy-OR CPT

Use k numbers $p_i \in [0, 1]$ to parameterize $O(2^k)$ elements of CPT:

$$P(Y=0 | X_1, X_2, \dots, X_k) = \prod_{i=1}^k (1-p_i)^{x_i} \quad \text{where } x_i \in \{0, 1\}$$

$$P(Y=1 | X_1, X_2, \dots, X_k) = 1 - \prod_{i=1}^k (1-p_i)^{x_i}$$

Why call "noisy-OR"?

- Look at $\underbrace{\text{all parents are off}}$

$$\begin{aligned} P(Y=1 | X_1=0, X_2=0, \dots, X_k=0) &= 1 - \prod_{i=1}^k (1-p_i)^0 \\ &= 1 - \prod_{i=1}^k (1) \\ &= 0 \end{aligned}$$

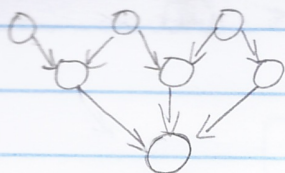
If $x_i=0$ for all i , then $Y=0$, recovers behavior of "OR".

- Look at

$$\begin{aligned} P(Y=1 | X_1=0, X_2=0, \dots, X_{i-1}=0, \underbrace{X_i=1}_{\text{on}}, X_{i+1}=0, \dots, X_k=0) \\ &= 1 - \prod_{j=1}^k (1-p_j)^{x_j} \\ &= 1 - (1-p_i)^1 \prod_{j \neq i}^k (1-p_j)^0 \\ &= 1 - (1-p_i) \\ &= p_i \end{aligned}$$

Intuitively, p_i represents the probability that $x_i=1$ by itself triggers $Y=1$.

Conditional independence



A node is conditionally independent of its non-parent ancestors given its parents.

$$P(x_i | pa(x_i)) = P(x_i | x_1, \dots, x_{i-1})$$

* More generally:

Let X , Y , and E refer to disjoint sets of nodes.

When is X conditionally independent of Y given evidence E ?

When is $P(X|E, Y) = P(X|E)$?

$$P(Y|E, X) = P(Y|E)?$$

$$P(X, Y|E) = P(X|E)P(Y|E)?$$

We've done already the special case:

$$X = \{x_i\}$$

$$E = \{pa(x_i)\}$$

$$Y = \{x_1, x_2, \dots, x_{i-1}\} - \{pa(x_i)\}$$

ancestors - parents

$$P(X|E) = P(X|E, Y)$$

* d-separation

"direction-dependent" separation

Relates conditional independence to graph-theoretic properties.

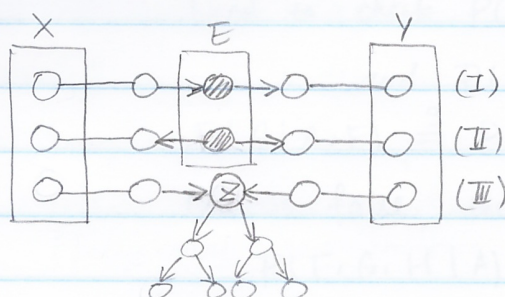
* Thm: $P(X, Y|E) = P(X|E)P(Y|E)$ if and only if every undirected path from a node in X to a node in Y is "d-separated" by E (blocked)

* Def. a path π is d-separated if there exists a node $Z \in \pi$ for which one of 3 conditions hold:

(I) $Z \in E$ with $\rightarrow \textcircled{Z} \rightarrow$ is an "intervening" event in a causal chain.

(I) $z \in E$ with $\longleftrightarrow$ is a "common explanation"

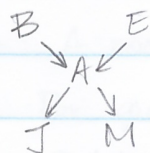
(II) $z \notin E$ descendants(z) $\notin E$ with $\rightarrow \textcircled{z} \leftarrow$.
no observed "common effects"



* Proof that d-separation $\iff$ conditional independence is beyond course

* Efficient algorithms exist for testing d-separation.

* Alarm BN example



True or False

$$(1) P(B|A, M) \stackrel{?}{=} P(B|A)$$

(Condition I)

True - alarm is the intervening event

$$(2) P(J, M|A) \stackrel{?}{=} P(J|A)P(M|A)$$

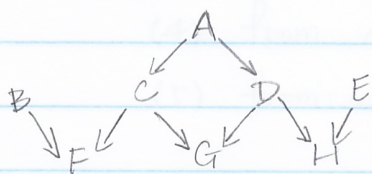
(Condition II)

True - alarm is common explanation

$$(3) P(B) \stackrel{?}{=} P(B|E)$$

True - independent when there is no observed common effect. (Condition III)

* Loopy BN example



True or False

$$P(D|H) \stackrel{?}{=} P(D|E, H) \quad \text{False by (III)}$$

$$P(F, H|A) \stackrel{?}{=} P(F|A)P(H|A)$$

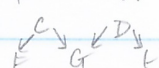
True

check path



true by (II)

check path



true by (III)

$$P(F, G, H | A) \stackrel{?}{=} P(F | A) P(G | A) P(H | A)$$

product rule $P(F | A) P(G | A, F) P(H | A, F, G) \stackrel{?}{=} P(F | A) P(G | A) P(H | A)$

$$\text{Need to check } P(G | A, F) \stackrel{?}{=} P(G | A)$$

$$\text{and } P(H | A, F, G) \stackrel{?}{=} P(H | A)$$

$$P(G | A, F) \stackrel{?}{=} P(G | A) \quad \text{False}$$

One is false, then

$$P(F, G, H | A) \stackrel{?}{=} P(F | A) P(G | A) P(H | A) \text{ False.}$$

* Def = Markov blanket B_x of individual node x consists of parents of x , children of x , and spouses of x

↓ parents of children of x , not including x itself.

* Thm: A node x is conditionally independent of nodes outside B_x given nodes inside B_x

$$P(x | B_x, Y) = P(x | B_x) \text{ when } Y \notin \{B_x, x\}$$

Proof: For any node $Y \notin \{B_x, x\}$, the undirected path from Y to x must pass through B_x . There are five cases to consider:

- | | |
|---|---------------------------|
| (1) from parent of parent of x — case I | } All paths are separated |
| (2) from child of parent of x — case II | |
| (3) from parent of spouse of x — case I | |
| (4) from child of child of x — case I | |
| (5) from child of spouse of x — case II | |

Inference

* Problem

E = set of evidence nodes

Q = set of query nodes

How to compute posterior probabilities $P(Q|E)$?

- * Question: When can we perform this efficiently?
(polynomial time in size of DAG and CPTs)?

Answer: Polytrees.

- * Def: polytree = singly connected network;
at most one undirected path between
any two nodes, no loops.

Goal: Compute $P(X|E)$

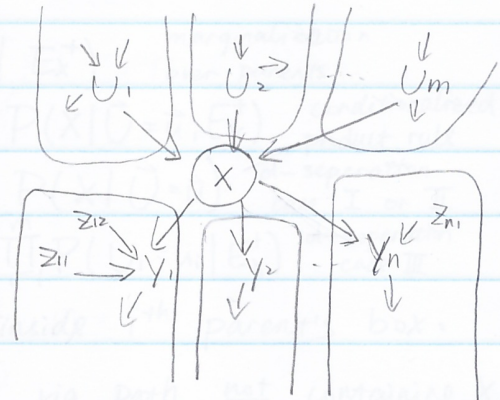
Claim: Boxes don't overlap! No loops!

Polytree: Node X

Parents $U_1, U_2, \dots, U_m$

Children $Y_1, Y_2, \dots, Y_n$

Spouse: Z_{jk} (k^{th} parent of j^{th} child)



Type of evidence:

E_x^+ = evidence connected to X thru its parents.

E_x^- = evidence connected to X thru its children

$$E = E_x^+ \cup E_x^-$$

Assume $X \notin E$, otherwise inference is trivial.

* Inference in polytree

$$\begin{aligned}
 P(X|E) &= P(X|E_x^-, E_x^+) \\
 &= \frac{P(X, E_x^- | E_x^+)}{P(E_x^- | E_x^+)} \quad \text{conditionalized product rule} \\
 &= \frac{P(X, E_x^- | E_x^+)}{\sum_x P(X=x, E_x^- | E_x^+)} \quad \text{conditionalized marginalization}
 \end{aligned}$$

Denominator is just numerator summed over x .

Focus on numerator.

* Numerator:

$$\begin{aligned}
 P(X, E_x^- | E_x^+) &= P(X | E_x^+) P(E_x^- | X, E_x^+) \quad \text{conditionalized product rule} \\
 &= P(X | E_x^+) P(E_x^- | X) \quad \text{conditional (case I) independent}
 \end{aligned}$$

goal: "upstream" recursion ↙ "downstream" recursion ↘

* "Upstream" recursion

$$\begin{aligned}
 P(X | E_x^+) &= \sum_{\vec{u}} P(X, \vec{U}=\vec{u} | E_x^+) \quad \text{marginalization over parents} \\
 &= \sum_{\vec{u}} P(\vec{U}=\vec{u} | E_x^+) P(X | \vec{U}=\vec{u}, E_x^+) \quad \text{conditionalized product rule} \\
 &= \sum_{\vec{u}} P(\vec{U}=\vec{u} | E_x^+) P(X | \vec{U}=\vec{u}) \quad \text{d-separation case I or II} \\
 &= \sum_{\vec{u}} P(X | \vec{U}=\vec{u}) \prod_{i=1}^m P(U_i=u_i | E_x^+) \quad \text{d-separation case III}
 \end{aligned}$$

Let $E_{u_i|x}$ denote evidence inside i^{th} parent's box:

evidence connected to U_i via path not containing X .

$$\begin{aligned}
 P(X | E_x^+) &= \sum_{\vec{u}} P(X | \vec{U}=\vec{u}) \prod_{i=1}^m P(U_i=u_i | E_x^+) \\
 &= \sum_{\vec{u}} \underbrace{P(X | \vec{U}=\vec{u})}_{\text{CPT at node } X} \prod_{i=1}^m \underbrace{P(U_i=u_i | E_{u_i|x})}_{\text{recursive instance of original problem}} \quad \text{d-separation case III}
 \end{aligned}$$

* "Downstream" recursion

How to compute $P(E_x^- | X)$?

Possible but somewhat more complicated.

* Termination conditions:

- root node (no parents)
 - leaf node (no child)
 - evidence node (trivial)
- } algorithm terminates b/c polytree has no loops.

* Running Time

- linear in # nodes of network

- linear in size of CPTs (b/c we must sum over the parent values)

$$\sum_{\vec{u}} P(x | \vec{u} = \vec{u}) \dots$$

Loopy BNs - how to perform inference

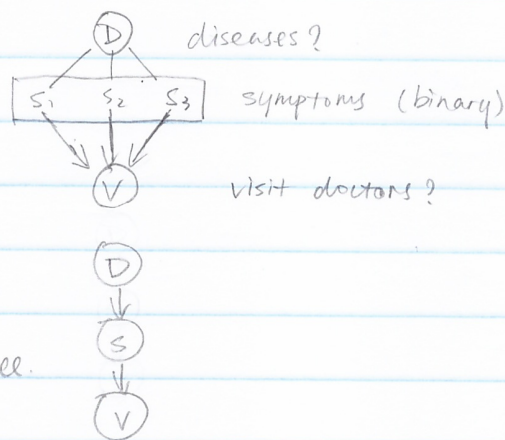
Ex: Medical Diagnosis

2-layer network



Ex: Simpler Example

How to do inference?



Turn Loopy BN into

polytree by clustering nodes.

Merge nodes to form polytree.

- Merge S_1, S_2, S_3 into some mega-node S

- S takes on 2^3 possible values

- Merge CPTs $P(S_1|D)$, $P(S_2|D)$, $P(S_3|D)$ into $P(S|D)$

- Apply polytree algorithm.

S_1	S_2	S_3	S
0	1	0	1
1	0	0	2
0	1	0	3
0	0	1	4
$\vdots$	$\vdots$	$\vdots$	$\vdots$

$P(S=2|D)$

$= P(S_1=1, S_2=0, S_3=0 | D)$

$P(V|S=2) = P(V|S_1=1, S_2=0, S_3=0)$

* Polytree algorithm

linear in size of CPTs

But CPT size grows exponentially with # nodes that are clustered.

How to choose optimal clustering?

Computationally hard results