### The splitting necklace problem

Frédéric Meunier

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CERMICS, Optimisation et Systèmes

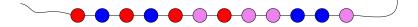
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#### Two thieves and a necklace

*n* beads, *t* types of beads,  $a_i$  (even) beads of each type.

Two thieves: Alice and Bob.

Beads fixed on the string.



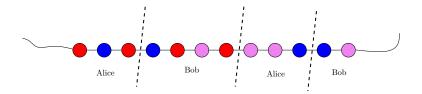
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Fair splitting = each thief gets  $a_i/2$  beads of type *i* 

#### The splitting necklace theorem

#### Theorem (Alon, Goldberg, West, 1985-1986)

There is a fair splitting of the necklace with at most t cuts.



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## t is tight

#### t cuts are sometimes necessary:



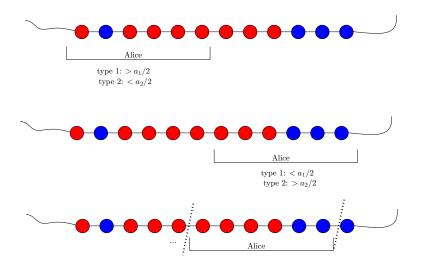


### Plan

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- 1. Proofs and algorithms
- 2. A special case: the binary necklace problem
- 3. Generalizations
- 4. Open questions

#### Easy proof when there are two types of beads



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## Definition

Binary Necklace Problem [Epping, Hochstättler, Oertel 2001]

**Input.** Necklace with *t* types of beads, 2 beads per type. i.e. n = 2t and  $a_i = 2$  for all *i*.

**Output.** Fair splitting minimizing the number of cuts.



Defined in an operations research context as the paintshop problem (automotive industry)

## Minimizing the number of cuts

Splitting necklace theorem:  $OPT \le t$ , but obvious: greedy algorithm.

Challenge here: optimization.

Proposition (Epping, Hochstättler 2006) The binary necklace problem is NP-hard.

Proof by MAX-CUT in 4-regular graphs.

M., Sebö 2009

APX-hard

Gupta et al. 2013

No polytime fixed-ratio approximation (assuming Unique Games Conjecture)

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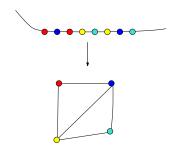
#### **Positive results**

M., Sebö 2009

OPT  $\leq \frac{3}{4}t + \frac{1}{4}\beta$ , where  $\beta =$  "forced cuts" ( $\bullet and \bullet \bullet \bullet$ ). Can be found in polytime.

Let G = ([t], E), with  $ij \in E$  if types *i* and *j* adjacent on necklace.

*G* planar  $\Rightarrow$  polytime.



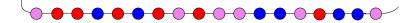
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## Generalizations

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#### q thieves and a necklace

*n* beads, *t* types of beads,  $a_i$  (multiple of *q*) beads of each type. *q* thieves: Alice, Bob, Charlie,...



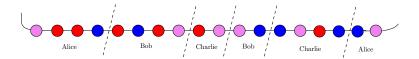
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### A generalization

Fair splitting = each thief gets  $a_i/q$  beads of type *i* 

#### Theorem (Alon 1987)

There is a fair splitting of the necklace with at most (q - 1)t cuts.



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# (q-1)t is tight

#### (q-1)t cuts are sometimes necessary:



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## **Open questions**

- Complexity of computing a fair splitting with at most *t* cuts when there are two thieves.
- Complexity of computing a fair splitting with at most (q-1)t cuts when there are q thieves.
- Existence of a fair splitting with choice of the advantaged thieves.
- Elementary proof of the splitting necklace theorem (any version).
- Simonyi's generalization for *q* thieves?
- Expected value of the optimum for the binary necklace problem.
- Extend results of the binary necklace problem to general necklaces (at least for two thieves).

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