Constraints and Triggers

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Reference: A First Course in Database Systems, 3rd edition, Chapter 7 (but not section 7.4)

Important Notices

• **Reminder:** Midterm is on **Monday, Nov 6** is usual classroom, at usual time.

- No make-ups, no early/late exams.
- You may bring a **single two-sided 8.5" x 11" sheet of paper** with as much info written (or printed) on it as you can fit and read unassisted.
 - No sharing of these sheets will be permitted.
 - Hand in these sheets at the end of the exam.
- Midterm will cover all material up to and including Lecture 7 (Views/Indexes).
- No devices are permitted during exam.
- You must show your UCSC ID at the end of the exam.
- Hope that everyone who needs DRC accommodation has already submitted form to me.
- Winter 2017 Midterm has been posted on Piazza; solution will be posted next week.
- Gradiance #3 is due on Friday, November 3; one problem uses Views!
- Sign-up for LSS tutoring with Alexander Ou ... if there's room.

Constraints and Triggers

- A *constraint* is a relationship among data elements that the DBMS is required to enforce.
 - Example: key constraints.
- *Triggers/Rules* are only executed when a specified condition occurs, e.g., insertion of a tuple.
 - Easier to implement than complex constraints.

Kinds of Constraints

- Keys/Unique constraints
- Foreign-key, or referential-integrity constraints
- Value-based constraints
 - Constrain values of a particular attribute
- Tuple-based constraints
 - Relationship among components of tuple
- Assertions
 - Any SQL boolean expression (<u>not implemented</u> in most relational DBMS, not discussed in this lecture)

Review: Single-Attribute Keys

- Place PRIMARY KEY or UNIQUE after the type in the declaration of the attribute.
- Example:

```
CREATE TABLE Beers (
    name CHAR(20) UNIQUE,
    manf CHAR(20)
);
```

Review: Multi-Attribute Key

• The bar and beer together are the key for Sells: CREATE TABLE Sells (

bar CHAR(20),

beer VARCHAR(20),

price REAL,

PRIMARY KEY (bar, beer)

);

Review: NULL

CREATE TABLE Sells (bar CHAR(20), beer VARCHAR(20), price REAL NOT NULL, PRIMARY KEY (bar, beer));

If the CREATE statement <u>didn't</u> include NOT NULL for price: ALTER TABLE Sells ALTER COLUMN price SET NOT NULL; ALTER TABLE Sells ALTER COLUMN price DROP NOT NULL;

Foreign Keys

- Values appearing in attributes of one relation must also appear together in specific attributes of another relation.
- Example:
 - In Sells(bar, beer, price), we might expect that a beer value also appears in Beers.name (the name column of the Beers table, the primary key for that table).
- Like a link/pointer, but based on value.

Expressing Foreign Keys

- Use keyword REFERENCES, either:
 - 1. After an attribute (for one-attribute keys)
 - As an element of the schema:
 FOREIGN KEY (<list of attributes>)
 REFERENCES <relation> (<attributes>)
- Referenced attributes must be declared as either PRIMARY KEY or UNIQUE.
 - (Why?)

Example: With Attribute

```
CREATE TABLE Beers (
 name CHAR(20) PRIMARY KEY,
 manf CHAR(20) );
CREATE TABLE Sells (
 bar CHAR(20),
 beer CHAR(20) REFERENCES Beers(name),
 price REAL );
```

Example: As Schema Element

CREATE TABLE Beers (name CHAR(20) PRIMARY KEY, manf CHAR(20)); CREATE TABLE Sells (bar CHAR(20), beer CHAR(20), price REAL, FOREIGN KEY (beer) REFERENCES Beers (name));

Example: Adding Foreign Key

CREATE TABLE Beers (

name CHAR(20) PRIMARY KEY,

manf CHAR(20));

CREATE TABLE Sells (

bar CHAR(20),

beer CHAR(20),

price REAL);

ALTER TABLE Sells ADD FOREIGN KEY (beer) REFERENCES Beers(name);

Enforcing Foreign-Key Constraints (Referential Integrity, RI)

- If there is a foreign-key constraint from *referring relation R* to *referenced relation S*, then violations may occur two ways:
 - 1. An insert or update to *R* introduces values that are not found in *S*, or
 - A deletion or update to S causes some tuples of *R* to "dangle", referencing a value that no longer exists

Actions Taken --- (1)

- Example: suppose *R* = Sells, *S* = Beers.
 - That is, Sells refers to Beers
- An insert or update to Sells that introduces a nonexistent beer must be rejected.
- A deletion or update to Beers that removes a beer value found in some tuples of Sells can be handled in one of three ways (next slide).

Actions Taken --- (2)

- 1. Default : Reject the modification.
- 2. Cascade : Make the same changes in Sells.
 - **Deleted beer**: delete Sells tuple.
 - Updated beer: also change value in Sells ...
 - ... so that Sells.beer has the same new value as Beers.name
- *3. Set NULL* : Change Sells.beer to NULL.

Example: Cascade

- Upon Delete of the Bud tuple from Beers:
 - Delete all tuples from Sells that have beer = 'Bud'
- Upon Update of the Bud tuple by changing 'Bud' to 'Budweiser':
 - Change all Sells tuples that have beer = 'Bud' to have beer = 'Budweiser'

Example: Set NULL

- Upon Delete of the Bud tuple from Beers:
 - Change all tuples of Sells that have beer = 'Bud' to have beer = NULL.
- Upon Update of the Bud tuple, changing 'Bud' to 'Budweiser':
 - Make the same change to tuples of Sells that have beer='Bud' as for deletion (making beer=NULL).

Choosing a Referential Integrity Policy

- When we declare a foreign key, we may choose policies SET NULL or CASCADE independently for deletions and updates.
- Follow the foreign-key declaration with: ON [UPDATE, DELETE][SET NULL, CASCADE]
- Two such clauses may be used, one for UPDATE and one for DELETE
- Otherwise, the DEFAULT (Reject) is used.

Example: Setting Policy

CREATE TABLE Sells (

bar CHAR(20),

beer CHAR(20),

price REAL,

FOREIGN KEY(beer)

REFERENCES Beers (name)

ON DELETE SET NULL

ON UPDATE CASCADE

);

<u>Attribute</u>-Based Check

- Constraint on the value of a particular attribute.
- CHECK(<condition>) may be added to the declaration for the attribute.
 - Condition must evaluate to TRUE or UNKNOWN; can't be FALSE.
- The condition may refer to the attribute of the relation that is being checked.
- But for the condition to reference <u>any other tuples or</u> <u>relations</u>, a subquery must be used.
 - Note: PostgreSQL <u>does not</u> support CHECK with subquery. Sigh.

Example: <u>Attribute</u>-Based Check

CREATE TABLE Sells (
 barCHAR(20),
 beer CHAR(20) CHECK (beer IN
 (SELECT name FROM Beers)),
 price REAL CHECK (price <= 5.00)
);</pre>

Example: Named Constraints

```
CREATE TABLE Sells (
   bar CHAR(20),
   beer CHAR(20) CHECK ( beer IN
      (SELECT name FROM Beers)),
   price REAL
   CONSTRAINT price_is_cheap
      CHECK ( price <= 5.00 )
);</pre>
```

```
ALTER TABLE Sells DROP CONSTRAINT price_is_cheap;
ALTER TABLE Sells ADD CONSTRAINT price_is_cheap
CHECK ( price <= 5.00 );
```

Timing of <u>Attribute</u>-Based Check

- Attribute-based checks are performed <u>only</u> when a value for <u>that attribute</u> is inserted or updated.
 - Example: CHECK (price <= 5.00) checks every new price and rejects the modification (for that tuple) if the price in Sells is more than \$5.
 - Example: CHECK (beer IN (SELECT name FROM Beers)) is <u>not</u> checked if a beer is deleted from Beers (unlike foreign-keys).

Tuple-Based Checks

- CHECK (<condition>) may be added as a relationschema element.
- The condition may refer to any attribute of the relation (same tuple).
- But for the condition to reference <u>any other</u> <u>tuples or relations</u>, a subquery must be used.
 - Condition is checked **only** on INSERT or UPDATE into relation that has the CHECK.

Example: <u>Tuple</u>-Based Check

• Only Joe's Bar can sell beer for more than \$5: CREATE TABLE Sells (

bar CHAR(20),

beer CHAR(20),

price REAL,

```
CHECK (bar = 'Joe''s Bar' OR
price <= 5.00)
```

);

When are Constraints Checked and Handled?

- In the CONSTRAINT clause, you may specify a constraint to be DEFERRABLE or <u>NOT DEFERRABLE</u>
 - DEFERRABLE INITIALLY DEFERRED or
 - DEFERRABLE <u>INITIALLY IMMEDIATE</u>
- Within a transaction, **SET CONSTRAINTS** determines whether:
 - SET CONSTRAINTS { ALL | name [, ...] } <u>IMMEDIATE</u>;
 - Deferrable constraints are checked immediately for each SQL statement, or
 - SET CONSTRAINTS { ALL | name [, ...] } DEFERRED;
 - Checking is deferred until the end of the transaction.

Example: Deferred Constraint

```
CREATE TABLE Sells (
  bar CHAR(20),
  beer CHAR(20) CHECK ( beer IN
    (SELECT name FROM Beers)),
  price REAL
  CONSTRAINT price_is_cheap
    CHECK ( price <= 5.00 )
    DEFERRABLE INITIALLY DEFERRED
);</pre>
```

SET CONSTRAINTS price_is_cheap IMMEDIATE;

Assertions

- These are database-schema elements, like relations or views.
- Defined by: CREATE ASSERTION <name> CHECK (<condition>);
- Condition may refer to any relation or attribute in the database schema.
- (Not implemented in most Relational DBMS because they're too complicated and expensive!)

Triggers: Motivation

- Assertions are powerful, but the DBMS often can't tell when they need to be checked ...
 - ... and they're probably not implemented by the DBMS.
- Attribute- and tuple-based checks are checked at known times, but they are not that powerful.
- Triggers let the user (often the DBA) decide when to check for any condition.

Event-Condition-Action Rules

- Another name for "trigger" is an ECA Rule, or Event-Condition-Action Rule
- Event: typically a type of database modification, e.g., "insert on Sells"
- *Condition* : Any SQL boolean-valued expression
- Action : Any SQL statements

Preliminary Example: A Trigger

 Instead of using a foreign-key constraint and rejecting insertions into Sells(bar, beer, price) with unknown beers, a trigger can add that beer to Beers, with a NULL manufacturer.

Example: Trigger Definition



CREATE TRIGGER

• Either:

CREATE TRIGGER <name>

• Or:

CREATE OR REPLACE TRIGGER <name>

• Useful if there is a trigger with that name and you want to modify the trigger.

Options: The Event

- AFTER INSERT can be BEFORE INSERT.
 - Also, can be INSTEAD OF, if the relation is a view.
 - A clever way to execute view modifications is to have triggers translate them to appropriate modifications on the base tables.
- INSERT can be DELETE or UPDATE.
 - And UPDATE can be UPDATE ON a particular attribute.

Options: FOR EACH ROW

- Triggers are either "row-level" or "statement-level."
- FOR EACH ROW indicates row-level; its absence indicates statement-level.
- *Row level triggers* : Execute once for each modified tuple.
- *Statement-level triggers* : Execute once for a SQL statement, regardless of how many tuples are modified.

Options: REFERENCING

- INSERT statements imply a new tuple (for row-level) or new table (for statement-level).
 - The "table" is the set of inserted tuples.
- DELETE implies an old tuple or table.
- UPDATE implies both.
- Refer to these by

[NEW OLD] [TUPLE TABLE] AS <name>

Options: The Condition

- Any boolean-valued condition.
- Evaluated on the database as it would exist before or after the triggering event, depending on whether BEFORE or AFTER is used.
 - But always before the changes take effect.
- Access the new/old tuple/table through the names in the REFERENCING clause.

Options: The Action

- There can be more than one SQL statement in the action.
 - Surround by BEGIN . . . END if there is more than one.
- But queries make no sense in an action, so we are really limited to modifications.

Another Example

• Using Sells(bar, beer, price) and a unary relation RipoffBars(bar), maintain a list of RipoffBars that raise the price of some beer by more than \$1.

