

Lecture-12 & 13: Finite State Machine

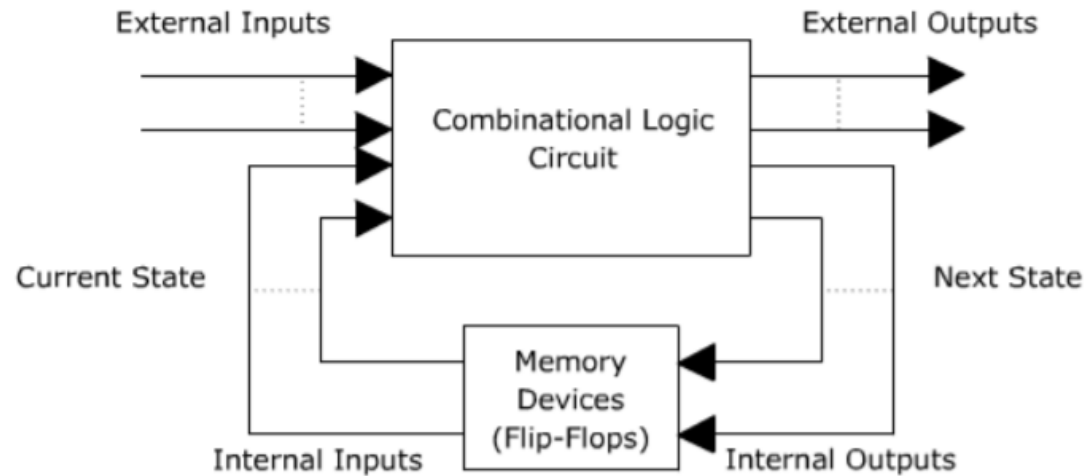
ECE-111 Advanced Digital Design Project

Vishal Karna

Winter 2022

Finite State Machine

- ❑ **Sequential circuits works on a clock cycle which may be synchronous or asynchronous.**
 - Sequential circuits use current inputs and previous inputs by storing the information and putting back into the circuit on the next clock cycle

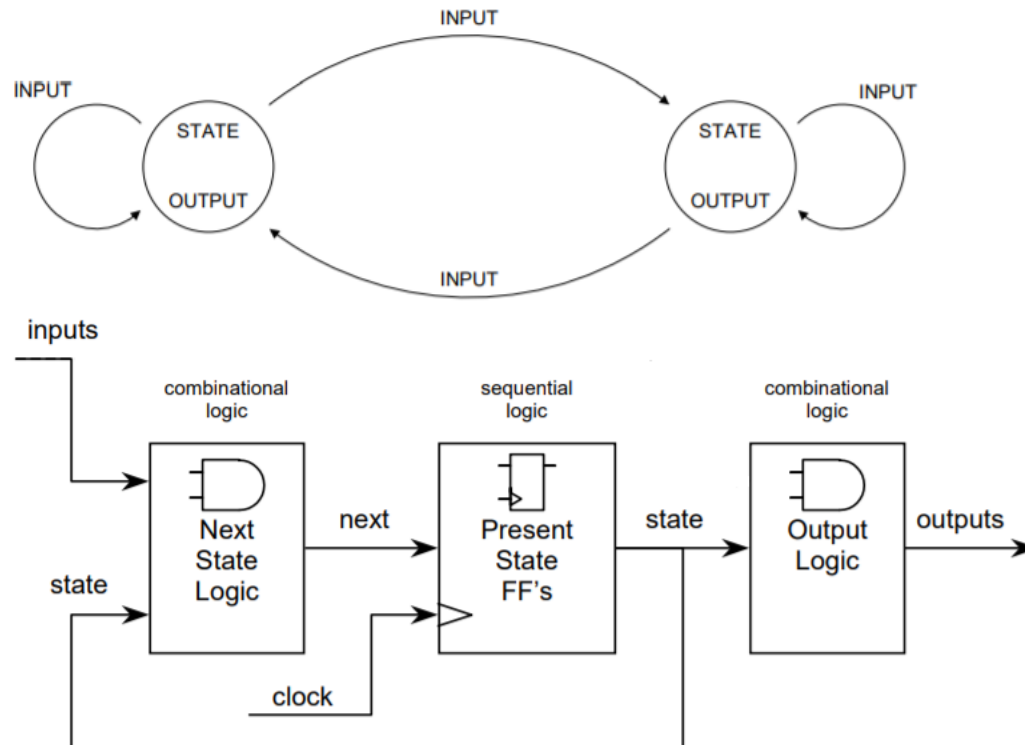


- ❑ **Finite State Machine (FSM) is a model used to design sequential logic circuits.**
 - It is conceived as an abstract machine that can be in one of a finite number of states.
 - The machine is in only one state at a time; the state it is in at any given time is called the current state.
 - It can change from one state to another when initiated by a triggering event or condition, this is called a transition.
 - A particular FSM is defined by a list of its states, and the triggering condition for each transition.
 - It can be implemented using models like Mealy and Moore machine.

Moore and Mealy Finite State Machine (FSM)

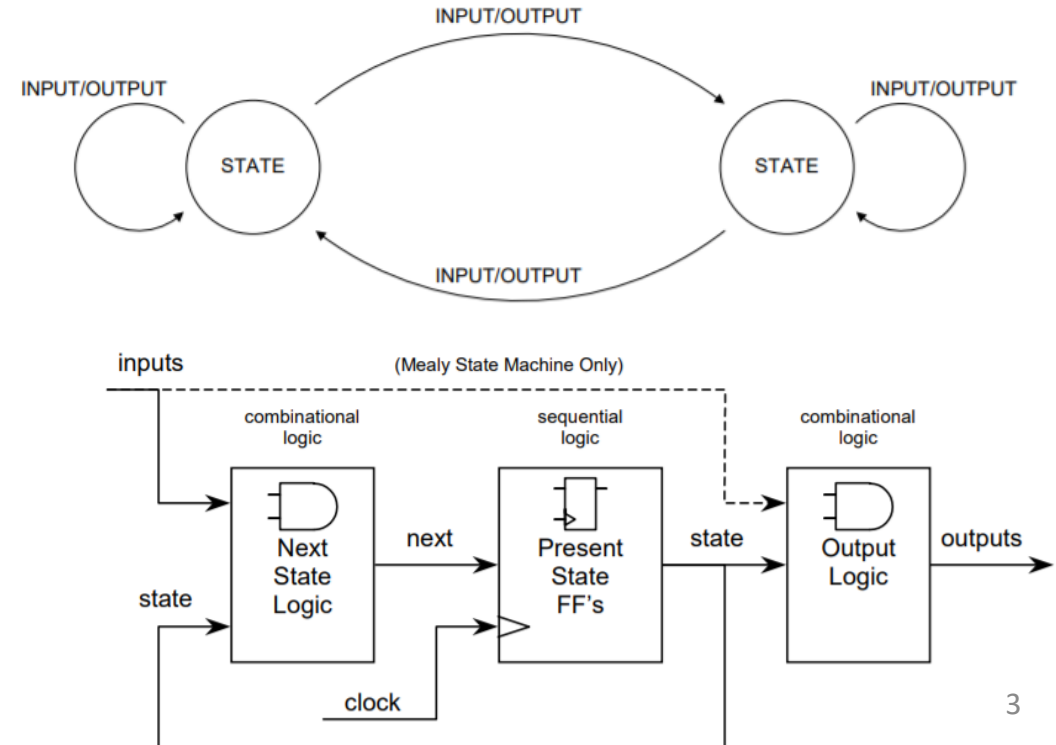
Moore FSM

- Output is solely based on present state of FSM
- Output is associated with a state
- Generally, more states than mealy (more hardware)
- More logic required to decode the outputs resulting in more circuit delays.
- Output react slower to input (one clock cycle later)
- Synchronous output and state generation



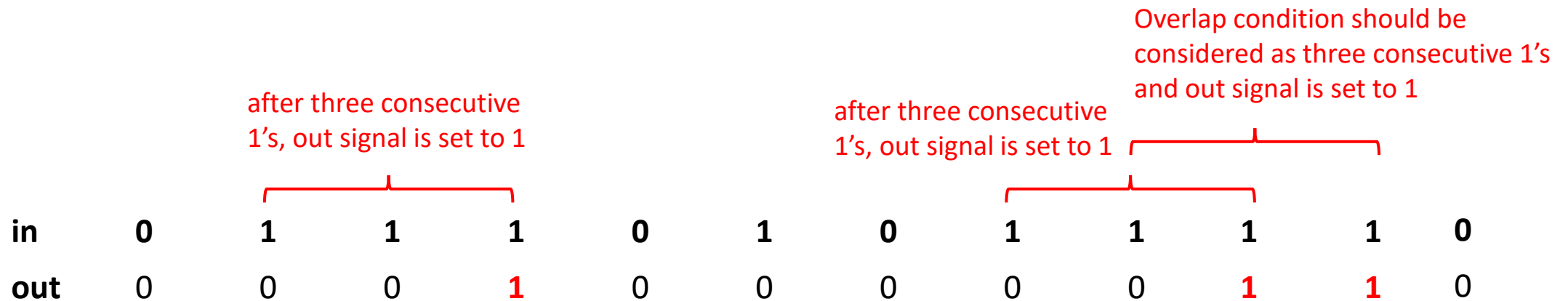
Mealy FSM

- Output based on present state and input(s)
- Output changes during transition of states
- Generally less states than moore
- Reacts faster to inputs and generally reacts in same cycle
- Asynchronous output generation
- Typically more complex to design than Moore



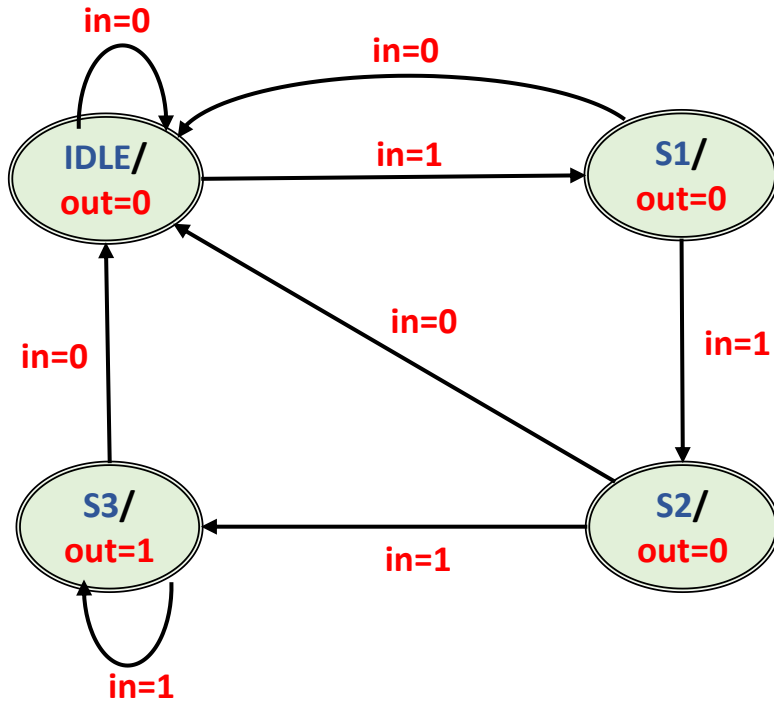
Moore and Mealy Finite State Machine

- ❑ FSM can be represented in form of state table or state transition diagram
- ❑ Example : Sequence Detector can be developed using Finite State Machine
 - Design a circuit to detect consecutive series of three or more '1's in serial input bit stream
 - Output will become '1' when three or more consecutive ones are detected
 - 4 states required to such sequence detector state machine :
 - State IDLE: reset state (zero 1s detected)
 - State S1: one 1 detected
 - State S2: two 1s detected
 - State S3: three 1s detected
 - Let's consider below mentioned input bit stream and observe output (out) behavior



Moore and Mealy Finite State Machine (FSM)

Moore FSM Diagram for Sequence Detector

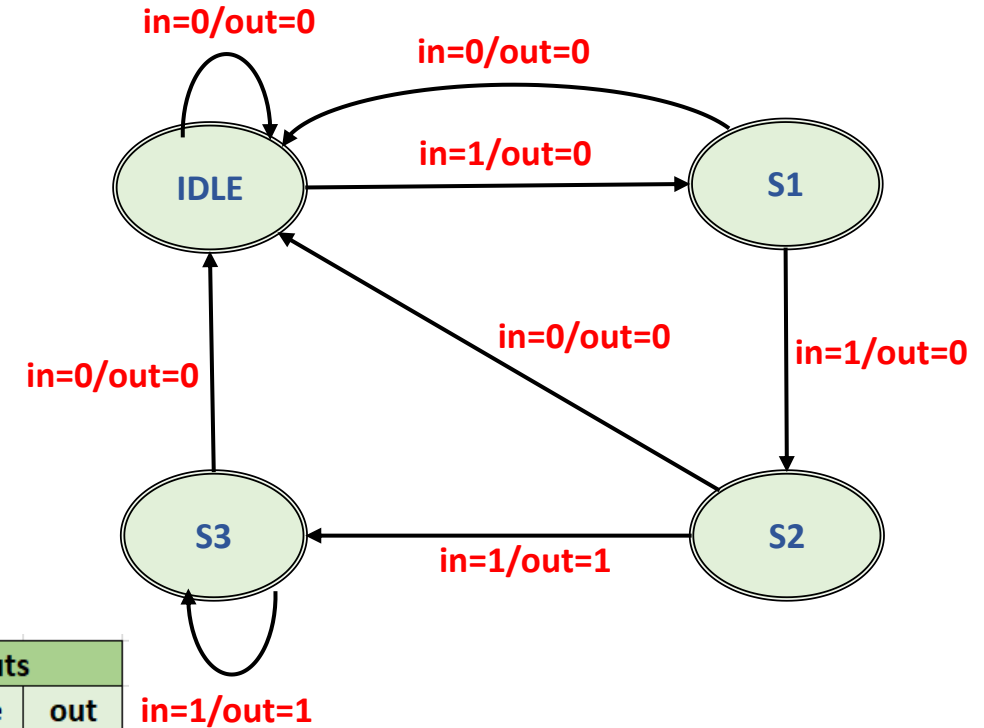


Note : In moore FSM, output is specified as part of present state

State Transition Table for Sequence Detector

inputs			outputs	
rstn	present_state	in	next_state	out
0	-	-	IDLE	0
1	IDLE	0	IDLE	0
1	IDLE	1	S1	0
1	S1	0	IDLE	0
1	S1	1	S2	0
1	S2	0	IDLE	0
1	S2	1	S3	1
1	S3	0	IDLE	0
1	S3	1	S3	1

Mealy FSM Diagram For Sequence Detector



Sequence Detector (Moore FSM – 2 always block approach)

```
module sequence_detector_moore(  
  input logic clk, rstn,  
  input logic in,  
  output logic out);
```

// Parameters to define FSM state encodings

```
localparam [1:0] IDLE=2'b00,  
                S1=2'b01,  
                S2=2'b10,  
                S3=2'b11;
```

State encodings are
declared as localparam
so that these cannot be
modified from outside

// Current state and next state variables

```
logic[1:0] present_state, next_state;
```

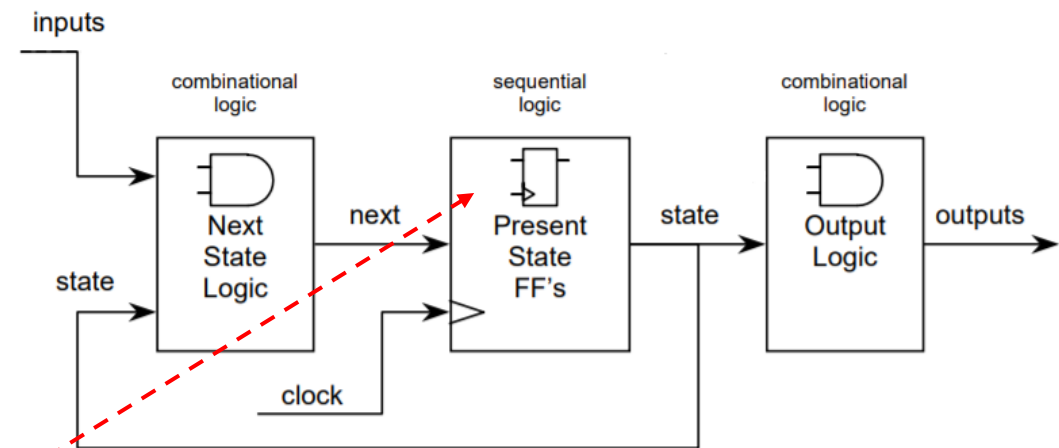
// Sequential Logic for present state

```
always_ff@(posedge clk) begin  
  if(!rstn)  
    present_state <= IDLE;  
  else  
    present_state <= next_state;  
end
```

Synthesis will
generate d-flipflop
for present_state

For sequential logic
Non-blocking
assignments used

(continued on next page....)



Sequence Detector (Moore FSM – 2 always block approach)

// Combination Logic for Next State and Output

```
always@(present_state,in) begin
```

```
case(present_state)
```

```
  IDLE: begin
```

```
    out = 0;
```

```
    if(in==1) next_state = S1;
```

```
    else next_state = IDLE;
```

```
  end
```

```
  S1: begin
```

```
    out = 0;
```

```
    if(in==1) next_state = S2;
```

```
    else next_state = IDLE;
```

```
  end
```

```
  S2: begin
```

```
    out = 0;
```

```
    if(in==1) next_state = S3;
```

```
    else next_state = IDLE;
```

```
  end
```

```
  S3: begin
```

```
    out = 1;
```

```
    if(in==1) next_state = S3;
```

```
    else next_state = IDLE;
```

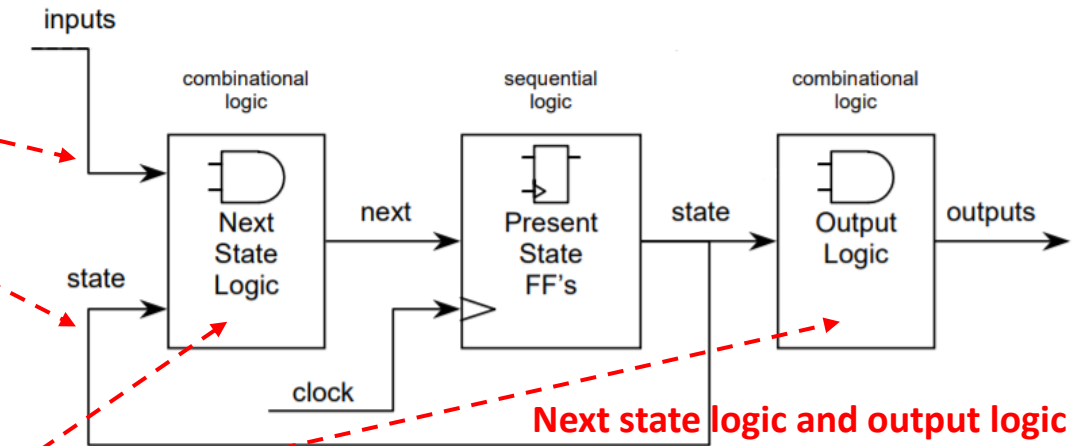
```
  end
```

```
(continued on next page....)
```

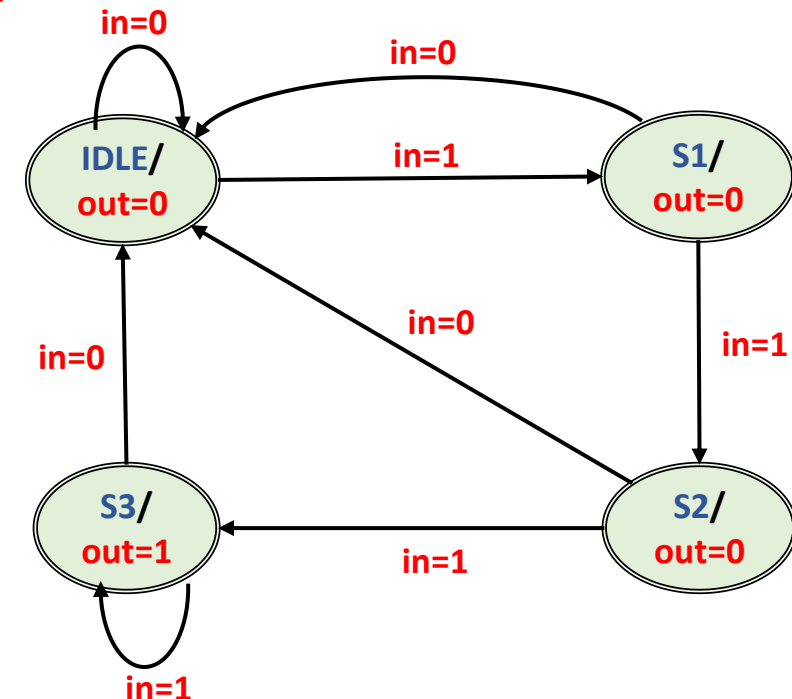
both present_state and
input "in" variable should
be listed in sensitivity list

For next state and
output combination
logic, blocking
assignments should
be used

For Moore FSM, output is set without
influence of input "in" signal and
purely based on present_state



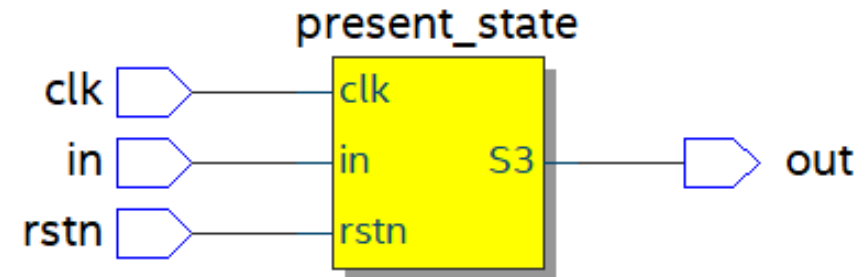
Next state logic and output logic
code in same always block !!!



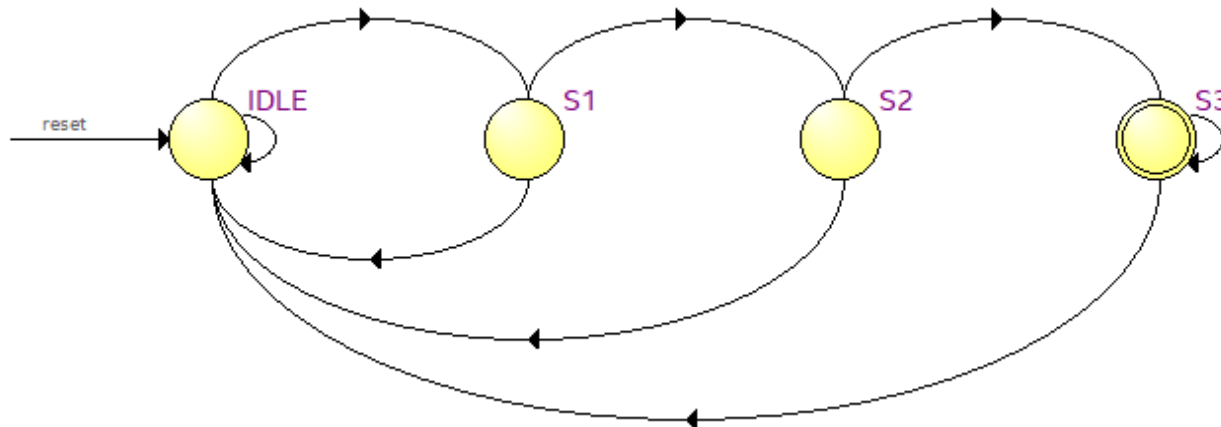
Sequence Detector (Moore FSM – 2 always block approach)

```
default: begin
  out = 0;
  next_state = IDLE;
end
endcase
end
endmodule: sequence_detector_moore
```

default is specified
in case a bad state
is reached



Sequence Detector RTL Netlist View
generated from Synthesizer



Sequence Detector State machine diagram generated by Synthesizer

	Source State	Destination State	Condition
1	IDLE	IDLE	(!in) + (in).(!rstn)
2	IDLE	S1	(in).(rstn)
3	S1	IDLE	(!in) + (in).(!rstn)
4	S1	S2	(in).(rstn)
5	S2	IDLE	(!in) + (in).(!rstn)
6	S2	S3	(in).(rstn)
7	S3	IDLE	(!in) + (in).(!rstn)
8	S3	S3	(in).(rstn)

State machine Transition Table

Sequence Detector (Moore FSM – 3 always block approach)

```
module sequence_detector_moore(  
  input logic clk, rstn,  
  input logic in,  
  output logic out);
```

// Parameters to define FSM state encodings

```
localparam [1:0] IDLE=2'b00,  
                S1=2'b01,  
                S2=2'b10,  
                S3=2'b11;
```

State encodings are
declared as localparam
so that these cannot be
modified from outside

// Current state and next state variables

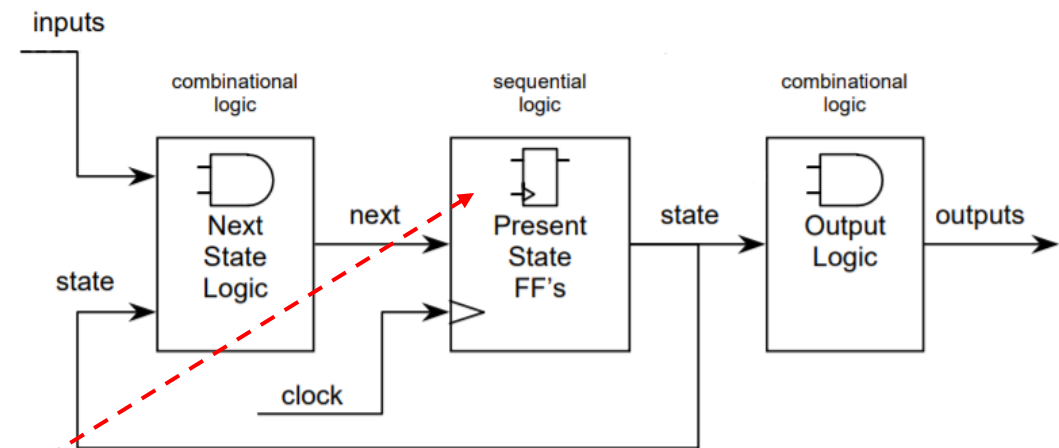
```
logic[1:0] present_state, next_state;
```

// Sequential Logic for present state

```
always_ff@(posedge clk) begin  
  if(!rstn)  
    present_state <= IDLE;  
  else  
    present_state <= next_state;  
end
```

1st always block for
present state
sequential logic

(continued on next page....)



Sequence Detector (Moore FSM – 3 always block approach)

// Combination Logic for Next State and Output

```
always@(present_state,in) begin
```

```
case(present_state)
```

```
  IDLE: begin
```

```
    if(in==1) next_state = S1;
```

```
    else next_state = IDLE;
```

```
  end
```

```
  S1: begin
```

```
    if(in==1) next_state = S2;
```

```
    else next_state = IDLE;
```

```
  end
```

```
  S2: begin
```

```
    if(in==1) next_state = S3;
```

```
    else next_state = IDLE;
```

```
  end
```

```
  S3: begin
```

```
    if(in==1) next_state = S3;
```

```
    else next_state = IDLE;
```

```
  end
```

```
  default: next_state = IDLE;
```

```
endcase
```

```
end
```

(continued....)

both present_state and
input "in" variable should
be listed in sensitivity list

```
always@(present_state) begin
```

```
case(present_state)
```

```
  S3: out = 1;
```

```
  default: out = 0;
```

```
endcase
```

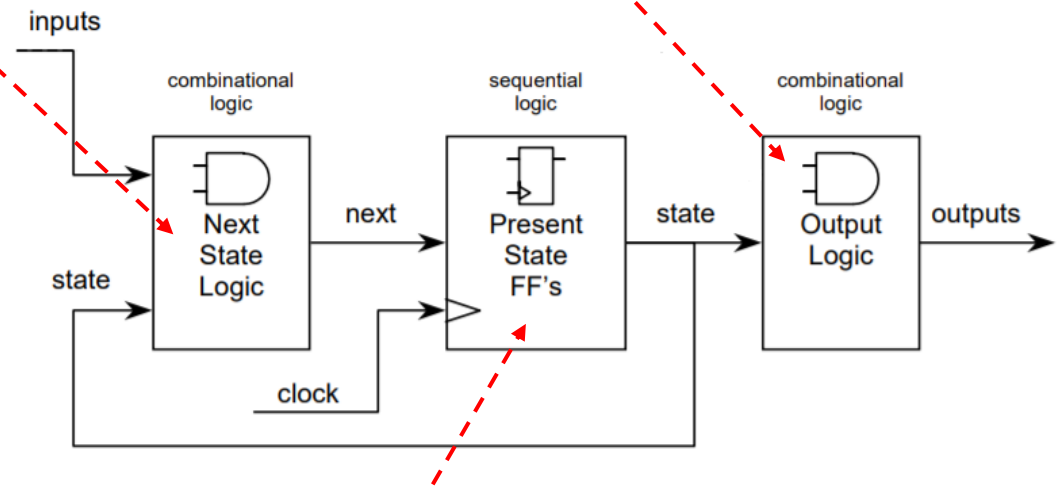
```
end
```

```
endmodule: sequence_detector_moore
```

Output is dependent
Only on present state
in moore

2nd always block for
separate combinational
block for next state logic

3rd always block
separate combinational
block for output
logic



1st always block for present
state FF (see previous slide)

Sequence Detector (Mealy FSM)

```
module sequence_detector_mealy(  
  input logic clk, rstn,  
  input logic in,  
  output logic out);
```

// Parameters to define FSM state encodings

```
localparam [1:0] IDLE=2'b00,  
                S1=2'b01,  
                S2=2'b10,  
                S3=2'b11;
```

State encodings are
declared as localparam
so that these cannot be
modified from outside

// Current state and next state variables

```
logic[1:0] present_state, next_state;
```

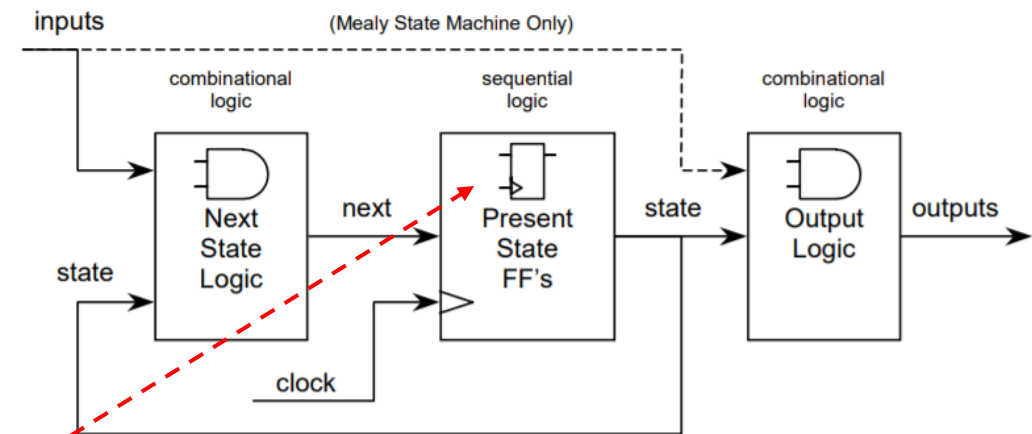
// Sequential Logic for present state

```
always_ff@(posedge clk) begin  
  if(!rstn)  
    present_state <= IDLE;  
  else  
    present_state <= next_state;  
end
```

Synthesis will
generate d-flipflop
for present_state

For sequential logic
Non-blocking
assignments used

(continued on next page....)



Sequence Detector (Mealy FSM)

// Combination Logic for Next State and Output

```
always@(present_state, in) begin
```

```
case(present_state)
```

```
  IDLE: begin
```

```
    if(in==1) begin
```

```
      next_state = S1;
```

```
      out = 0;
```

```
    end
```

```
  else begin
```

```
    next_state = IDLE;
```

```
    out = 0;
```

```
  end
```

```
end
```

```
S1: begin
```

```
  if(in==1) begin
```

```
    next_state = S2;
```

```
    out = 0;
```

```
  end
```

```
  else begin
```

```
    next_state = IDLE;
```

```
    out = 0;
```

```
  end
```

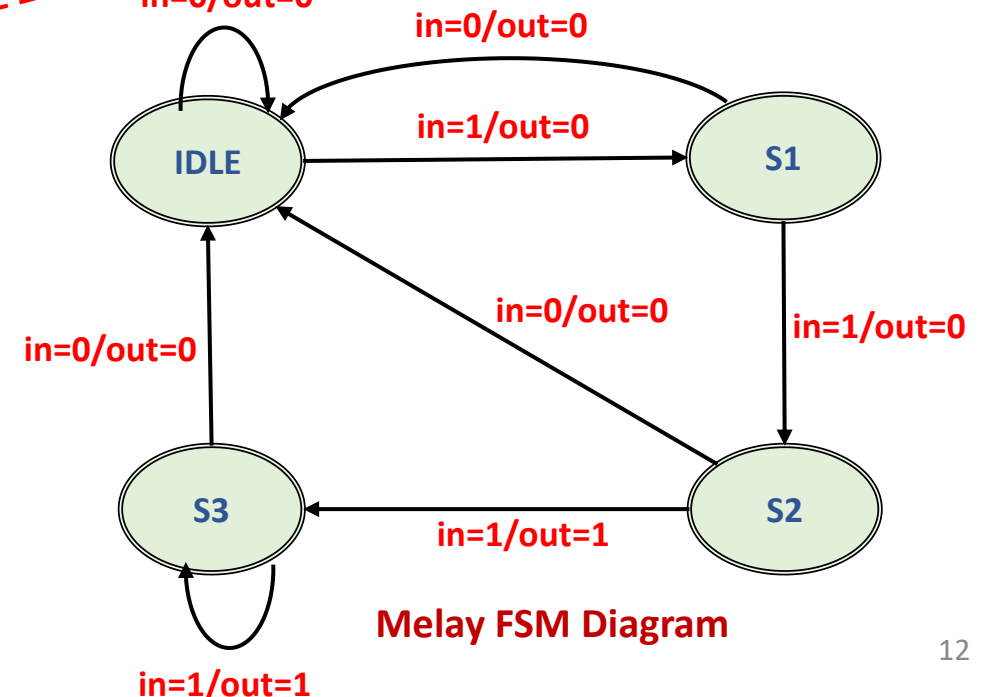
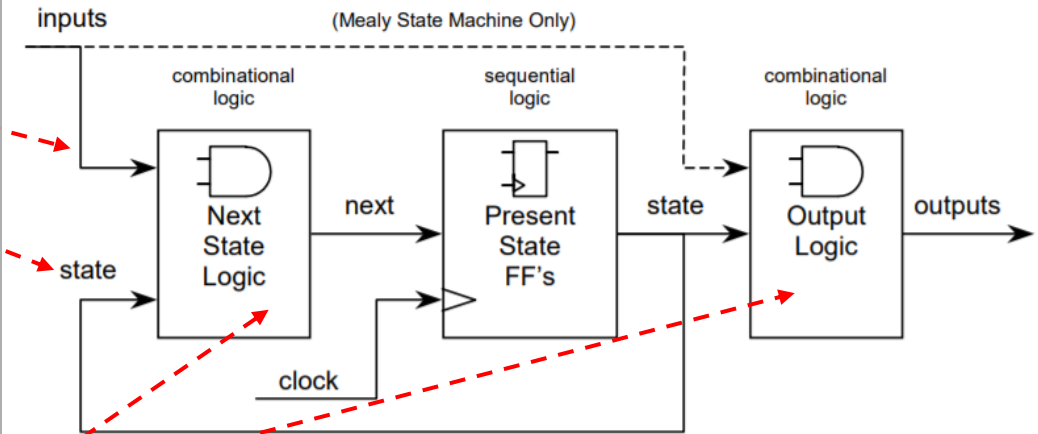
```
end
```

(continued on next page....)

both present_state and input "in" variable should be listed in sensitivity list

For next state and output combination logic, blocking assignments should be used

For Mealy FSM, output is set based on influence of both input "in" signal and present_state
Note : out = 0; statement is specified within if(in == 1) condition in mealy. And in moore, out = 0; is specified outside if(in == 1) condition



Mealy FSM Diagram

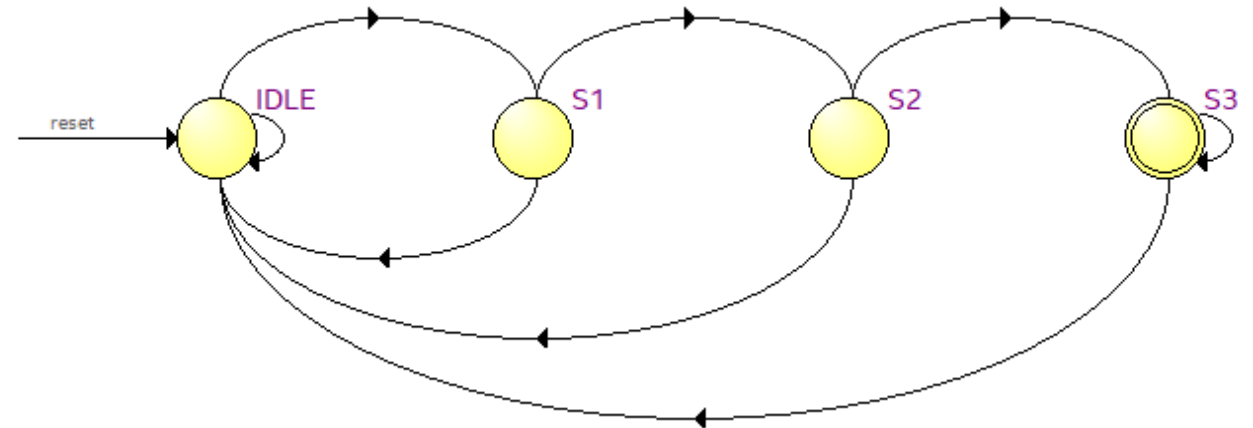
Sequence Detector (Mealy FSM)

```

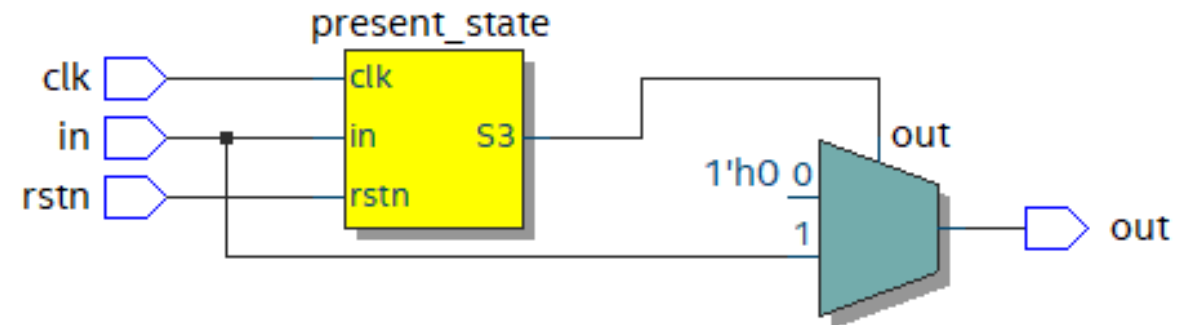
S2: begin
  if(in==1) begin
    next_state = S3;
    out = 1;
  end
  else begin
    next_state = IDLE;
    out = 0;
  end
end
S3: begin
  if(in==1) begin
    next_state = S3;
    out = 1;
  end
  else begin
    next_state = IDLE;
    out = 0;
  end
  default: begin out = 0; next_state = IDLE; end
endcase
end
endmodule: sequence_detector_mealy
    
```

For each state, under else condition if "out=0" is not present then Synthesis compiler will create latch for the output "out" signal

default is specified in case a bad state is reached



Sequence Detector State machine diagram generated by Synthesizer

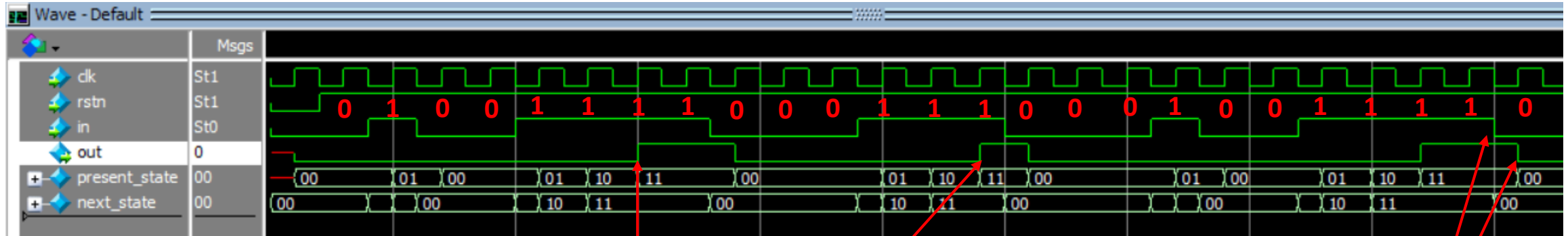


Sequence Detector RTL Netlist view generated by Synthesizer

Sequence Detector Simulation Snapshot(Moore FSM vs Mealy FSM)

Sequence Detector Moore FSM Simulation Snapshot

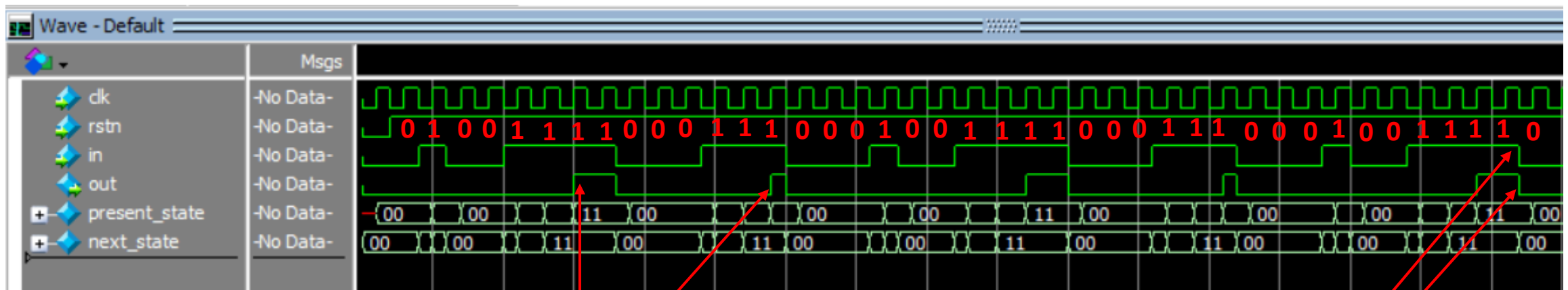
(Note : simulation result for 2 always vs 3 always block approach is same)



After three 1' detected out goes to '1'

Output did not react to change in input immediately in moore

Sequence Detector Mealy FSM Simulation Snapshot



After three 1' detected out goes to '1'

Output reacted to change in input immediately in mealy

Sequence Detector (One always block approach)

```
module sequence_detector_one_always_block(
    input logic clk, rstn,
    input logic in,
    output logic out);

// Parameters to define FSM state encodings
localparam [1:0] IDLE=2'b00,
               S1=2'b01, S2=2'b10, S3=2'b11;

// Current state and next state variables
logic[1:0] state; // Does not need two separate state variable

// Use of same clocked always block
always_ff@(posedge clk) begin
    if(!rstn)
        state <= IDLE;
        out <= 0;
    else
        case(state)
            IDLE: begin
                out <= 0;
                if(in==1) state <= S1;
                else state <= IDLE;
            end
```

Inputs are no longer asynchronously sampled.
Change in input is captured with respect to clock edge event

Use of non-blocking assignment statement in case encoding branches

// Combination Logic for Next State and Output

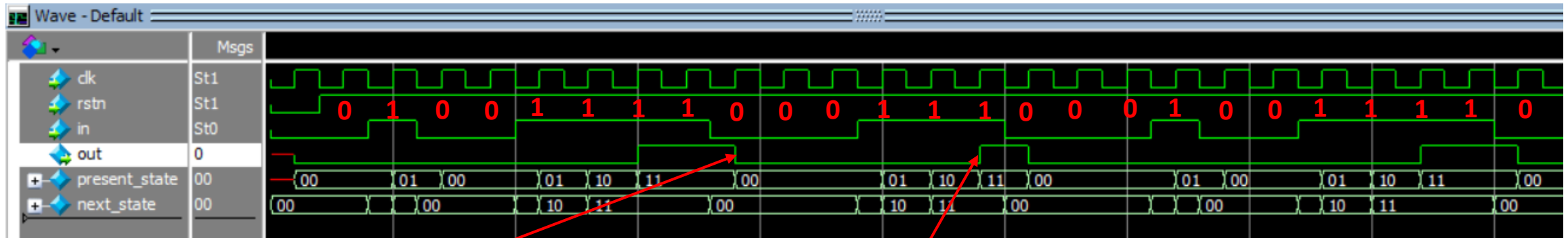
```
S1: begin
    out <= 0;
    if(in==1) state <= S2;
    else state <= IDLE;
end
S2: begin
    out <= 0;
    if(in==1) state <= S3;
    else state <= IDLE;
end
S3: begin
    out <= 1;
    if(in==1) state <= S3;
    else state <= IDLE;
end
default: begin
    out <= 0;
    state <= IDLE;
end
endcase
end
endmodule: sequence_detector_one_always_block
```

Output will stay asserted longer even when state has transitioned to another state where output should have gone back to reset or some other value

Sequence Detector Simulation Snapshot (One block vs Two block approach)

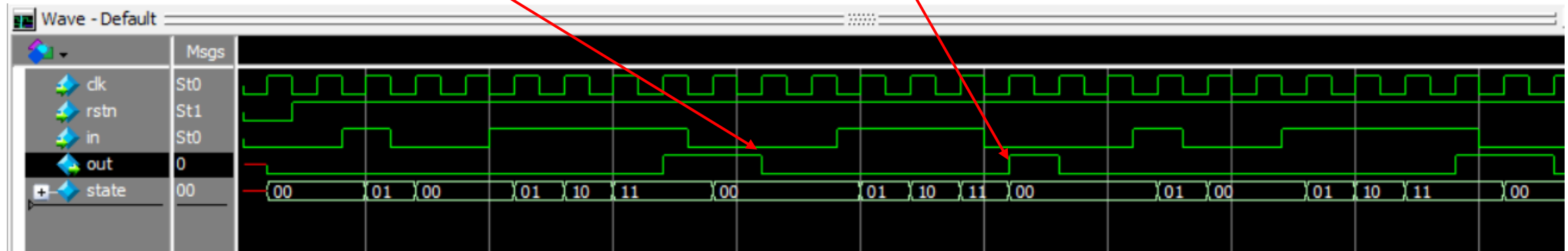
Sequence Detector Moore FSM Simulation Snapshot

(Note : simulation result for 2 always vs 3 always block approach is same)



In case of 1 always block implementation out = 1 stays for additional one additional clock cycle then 2 or 3 always block approach

out = 1 took additional one cycle in case of 1 always block approach compared to 2 or 3 always block implementation



Sequence Detector FSM Simulation Snapshot

(Note : simulation result for 1 always block

One Always Block Approach For FSM Modeling

- ❑ **One always block state machine is slightly more simulation-efficient than the two always block state machine :**
 - Since the inputs are only examined on clock changes (less work for simulator)

- ❑ **There some dis-advantages of one always block approach :**
 - RTL simulation does not model accurate gate level implementation
 - At the gate level combinational logic output updates when any input changes
 - In one process RTL simulation, combinational logic mixed with sequential logic, it is only evaluated at the clock edge
 - State machine can be more difficult to modify and debug since all sequential and combinational logic is mixed in one always block
 - Placing output assignments inside of the always block will infer output flip-flops.
 - This might lead to late availability of output. Sometimes not desirable in some applications.

- ❑ **Output assignments inside of a sequential always block cannot be Mealy outputs**
 - Hence cannot model Mealy FSM using one always block approach !

State Encoding

❑ States in FSM are represented by encoded value. There are multiple choices:

- Binary encoding, One-hot encoding, Gray encoding, Johnson encoding and more.
- Binary encoding and one-hot encoding are two commonly choices.

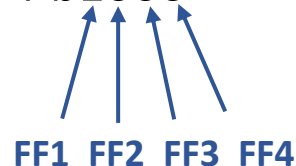
❑ Binary encoding of states

- binary encoding For N states, use $\text{ceil}(\log_2 N)$ bits to encode the state with each state represented by a unique combination of the bits.
- Tradeoffs:
 - Most efficient use of state registers (less number of Flipflops are required)
 - Slower and more complicated combinational logic to detect when in a particular state.
- Example :
 - In case of 4 states, 2 Flipflops are required to represent each state
 - **enum** logic [1:0] {RESET = 2'b00, WAIT = 2'b01, LOAD = 2'b10, DONE = 2'b11} next_state

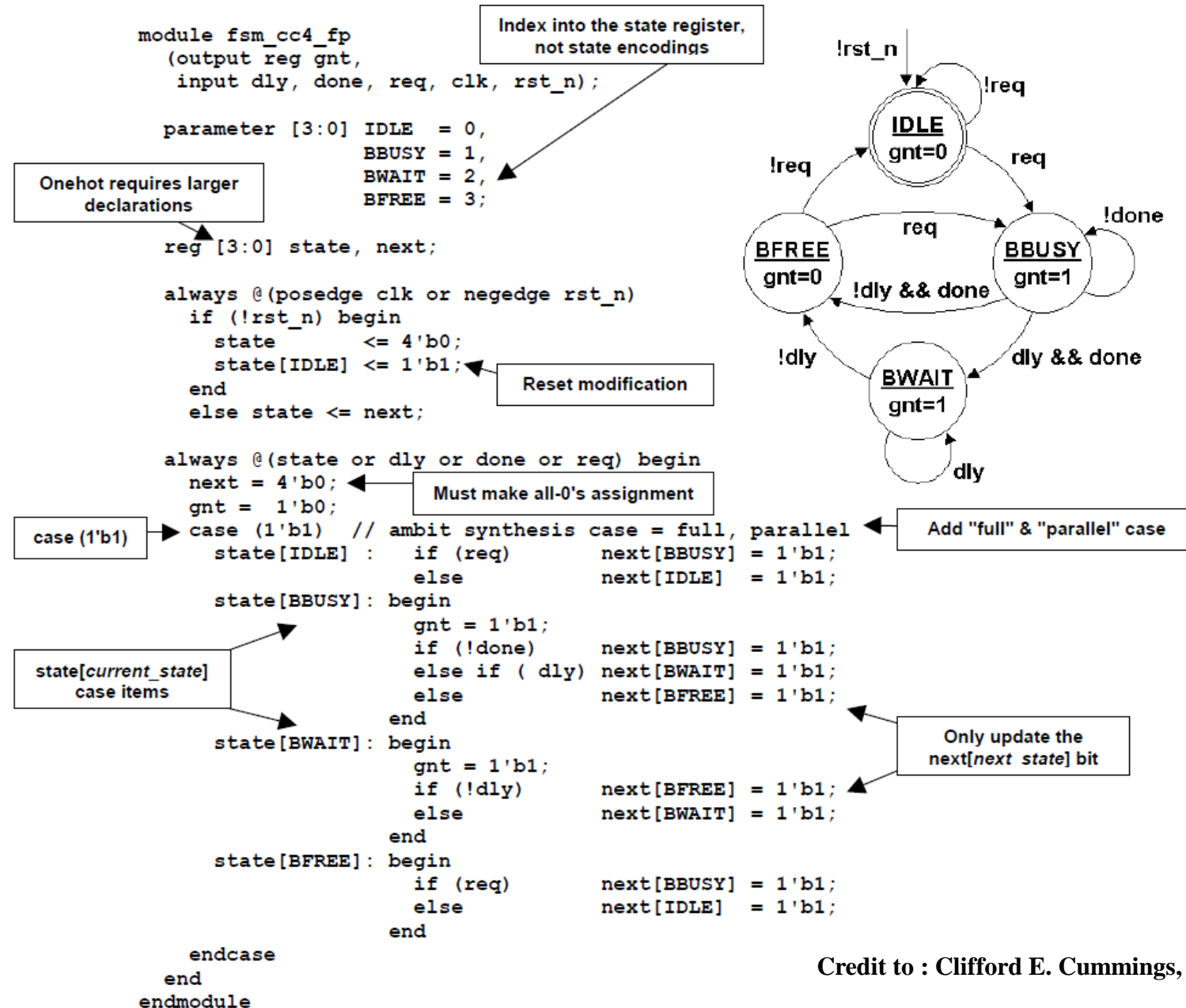
State Encoding

❑ One-hot encoding of states :

- For **N** states, use **N** bits to encode the state
 - bit corresponding to the current state is 1, all the others 0.
- **Tradeoffs:**
 - Leads to simpler and Faster design. Much less combinational logic for state decoding
 - Good alternative when trying to optimize speed or to reduce power dissipation.
 - Can be costly in terms of FFs for FSMs with large number of states
 - FPGAs have larger number for Flipflops. Therefore one-hot state machine encoding is often a very appropriate with most FPGA targets
- **Example :** In case of 4 states, 4 Flipflops are required for state encoding (one FF per state)
 - **STATE S1** : 4'b0001
 - **STATE S2** : 4'b0010
 - **STATE S3** : 4'b0100
 - **STATE S4** : 4'b1000



One-Hot FSM Implementation Example

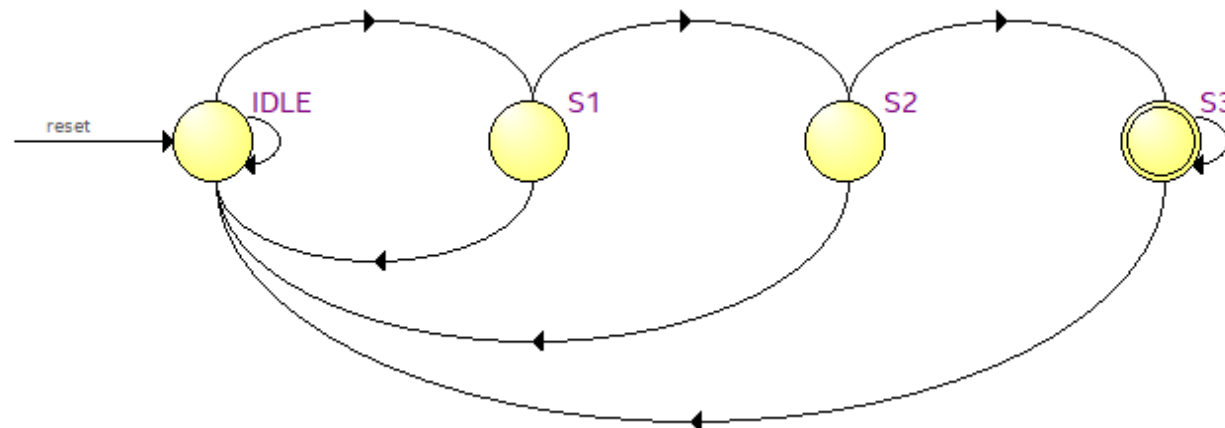


One Hot State Encoding

- ❑ In case of one-hot encoding style FSM, Quartus Prime will not auto-generate FSM state diagrams !
 - **Example:** if sequence detector state machine state encoding is changed from binary counting to one hot encoding style as described below, no state machine diagram generated however it will still work as an FSM

```
parameter[3:0] IDLE=4'b0001, S1=4'b0010, S2=4'b0100, S3=1000;  
logic[3:0] present_state, next_state;
```

Sequence Detector State machine diagram will not be generated by Quartus Synthesizer
Functionally logic will still behave as a correct FSM !



State Encoding Constants and Variables Declaration

❑ There are multiple ways to declare FSM state encoding constants and variables :

- **Using enumeration :**

- **enum** logic[1:0] {**WAIT**=2'b00, **EDGE**=2'b01} present_state, next_state;

- **Using Parameter :**

- **parameter** logic [1:0] **RESET** = 2'b00, **WAIT** = 2'b01, **LOAD** = 2'b10, **DONE** = 2'b11;
logic [1:0] present_state, next_state;

- **Using Local Parameter :**

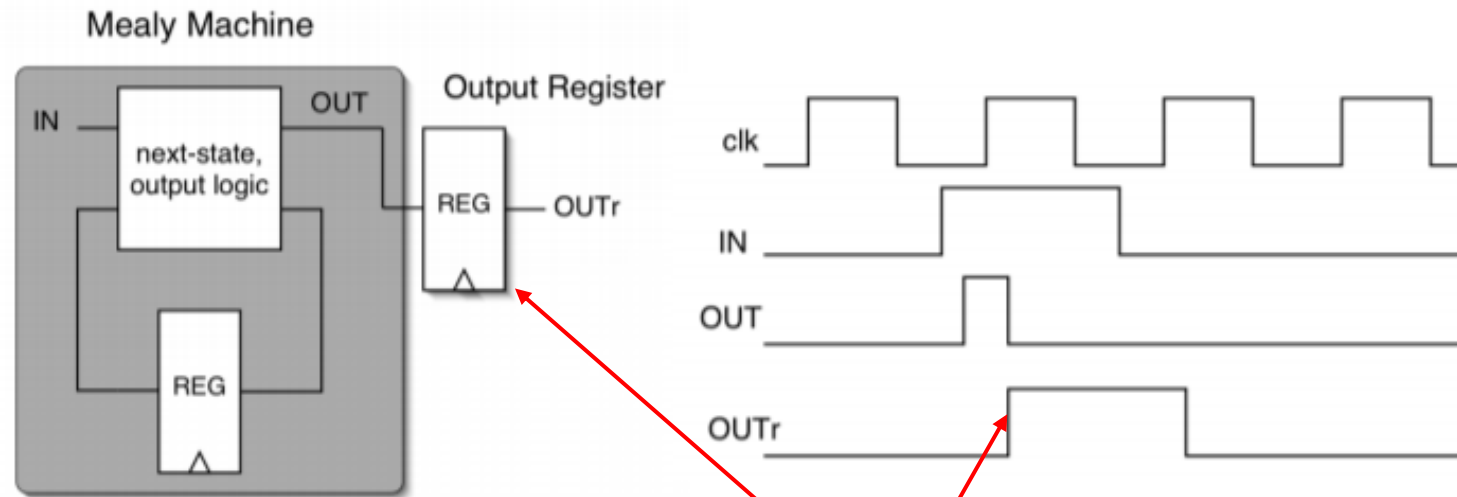
- **localparam** logic [1:0] **RESET** = 2'b00, **WAIT** = 2'b01, **LOAD** = 2'b10, **DONE** = 2'b11;
logic [1:0] present_state, next_state;

- **Using Using typedef : (Introduced in SystemVerilog. Not supported in Verilog)**

- **typedef enum** logic [1:0] {**WAIT**=2'b00, **EDGE**=2'b01} **e_states**;
e_states present_state, next_state;

Final Note on Moore vs Mealy FSM Modeling

- ❑ A given state machine can have both Moore and Mealy style outputs.
 - Nothing wrong with either implementation, but one to be aware of the timing differences between the two types.
 - The output timing behavior of the Moore machine can be achieved in a Mealy machine by “**registering**” the Mealy output values
 - **Note : registering output or registered output means having flipflop at the output signal**



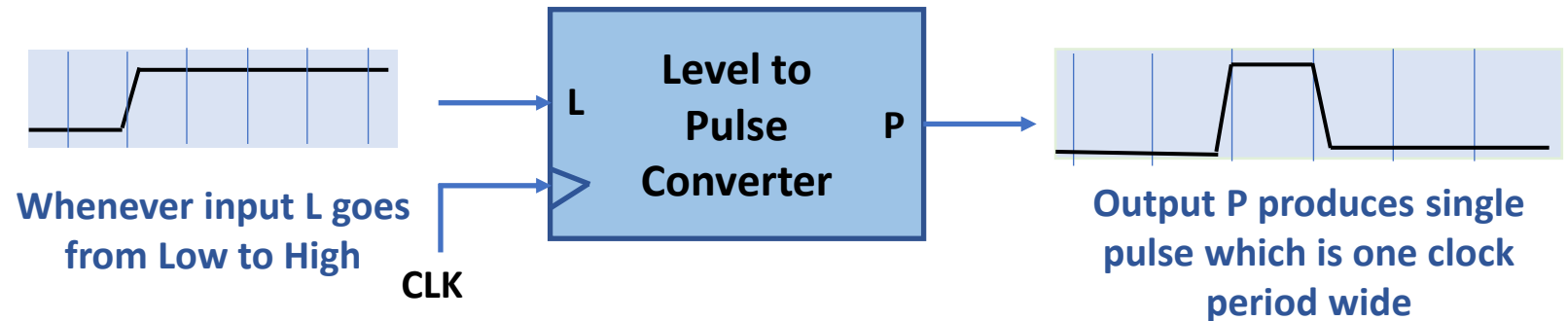
**Registered Output of Mealy matches
the output timing behavior of Moore
machine**

FSM Design Steps Using SystemVerilog

- ☐ Specify circuit function
- ☐ Draw state transition diagram
- ☐ Minimize number of States
- ☐ Derive state transition table
- ☐ Determine next state and output function
- ☐ Assign encodings (bit patterns) to symbolic states
- ☐ Implement State Machine using SystemVerilog :
 - Use either of the parameters/enum/localparam/typedef to represent encoded states.
 - Use separate always blocks for next state register assignment and combinational logic blocks
 - Use **always_comb** for all combinational block specification
 - Use **case** within combinational block (including default case item expression).
 - Within each **case** section assign all outputs and next state value based on inputs.
 - Ensure **default** case item is specified to avoid latches
- ☐ Note:
 - For Moore style machine make outputs dependent only on state not dependent on inputs.
 - Try not to mix up combinational logic and sequential logic inside same always block

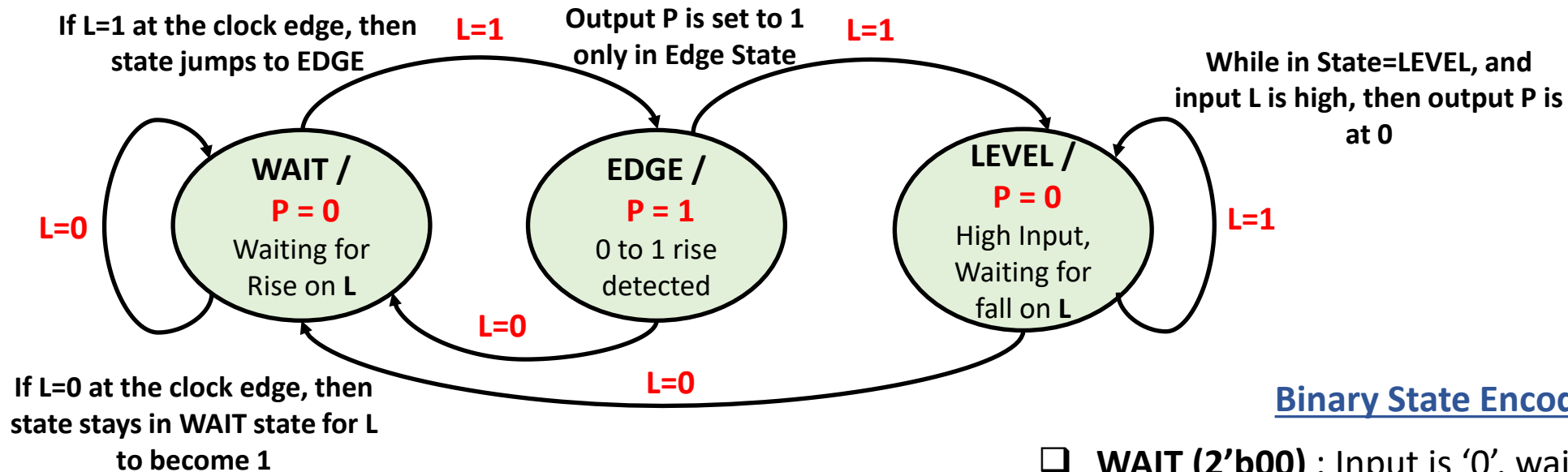
Level to Pulse Converter

- ❑ A level-to-pulse converter produces a single cycle pulse each time its input goes high
 - It is also known as synchronous rising edge detector
 - **Usage:** Pushing a button of a signal traffic light controller at pedestrian crossing
 - Pressing button for arbitrary period of time by pedestrian should generate single-cycle enable signals for counters in traffic light controller system



Level to Pulse Converter : Moore and Mealy State Diagrams

Moore State Transition Diagram

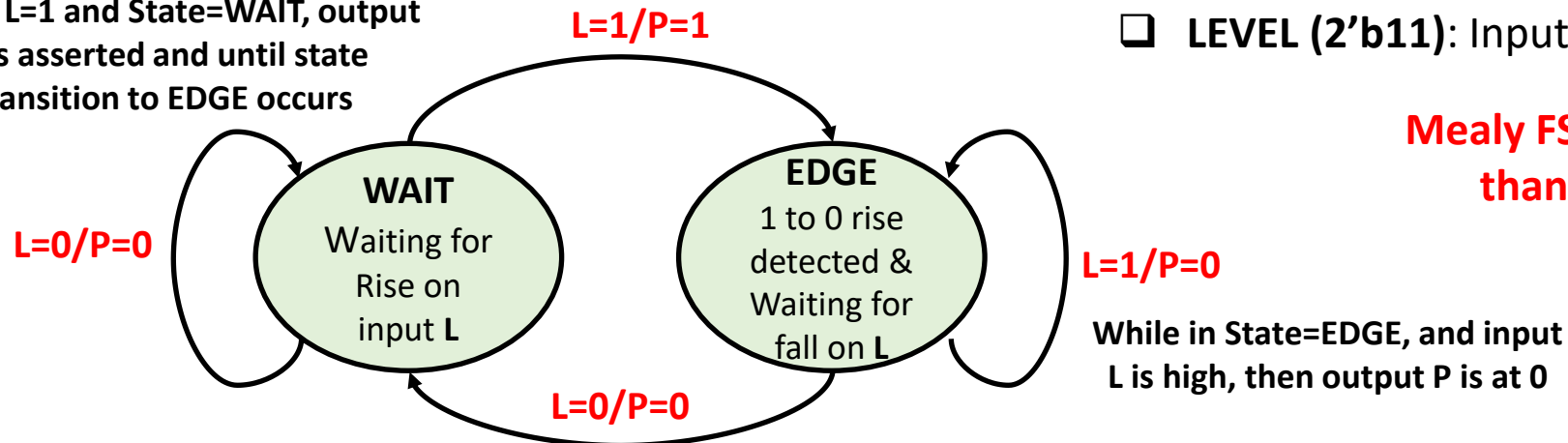


Binary State Encoding

- ❑ WAIT (2'b00) : Input is '0', wait for '1'
- ❑ EDGE (2'b01) : '0' to '1' edge detected on Input
- ❑ LEVEL (2'b11) : Input is stable at '1'

Mealy State Transition Diagram

When L=1 and State=WAIT, output P is asserted and until state transition to EDGE occurs



Mealy FSM has 1 less state than Moore FSM !!

Level to Pulse Converter (Moore FSM)

```
module level_to_pulse_converter_moore(  
    input logic clk, rstn,  
    input logic L,  
    output logic P);  
  
// FSM state encodings and state registers declaration  
enum logic[1:0] {WAIT=2'b00,  
                 EDGE=2'b01,  
                 LEVEL=2'b11} present_state, next_state;
```

State variables are declared as enumerated encoded logic

// Sequential Logic for present state

```
always_ff@(posedge clk) begin  
    if(!rstn)  
        present_state <= WAIT;  
    else  
        present_state <= next_state;  
end
```

Synthesis will
generate d-flipflop
for present_state

For sequential logic
Non-blocking
assignments used

(continued....)

// Combination Logic for Next State and Output

```
always_comb begin  
    case(present_state)  
        WAIT: begin  
            P = 0;  
            if(L==1) next_state = EDGE;  
            else next_state = WAIT;  
        end  
        EDGE: begin  
            P = 1;  
            if(L==1) next_state = LEVEL;  
            else next_state = WAIT;  
        end  
        LEVEL: begin  
            P = 0;  
            if(L==1) next_state = LEVEL;  
            else next_state = WAIT;  
        end  
        default: begin  
            P = 0; next_state = WAIT;  
        end  
    endcase  
end  
endmodule: level_to_pulse_converter_moore
```

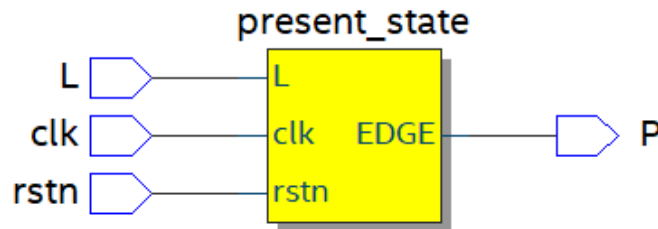
for next state combinational logic
always_comb specified to
automatically infer sensitivity list

Output P is set to '1' as soon as rising edge
on P is detected

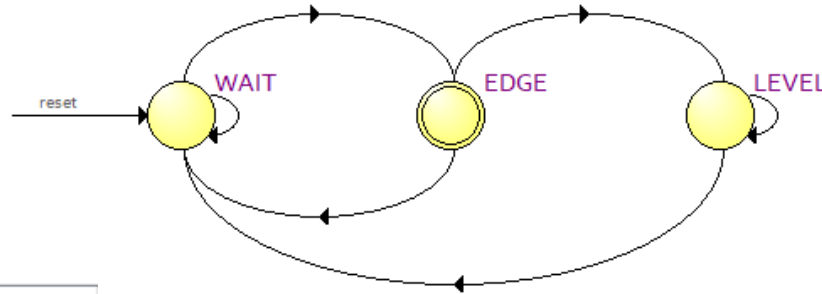
Since output P requirement is single cycle
pulse, P is set to '0' if input stays at level '1'
after rising edge detection

Level to Pulse Converter Simulation and Synthesis Results (Moore FSM)

Post Synthesis RTL Netlist Schematic



Moore State Diagram

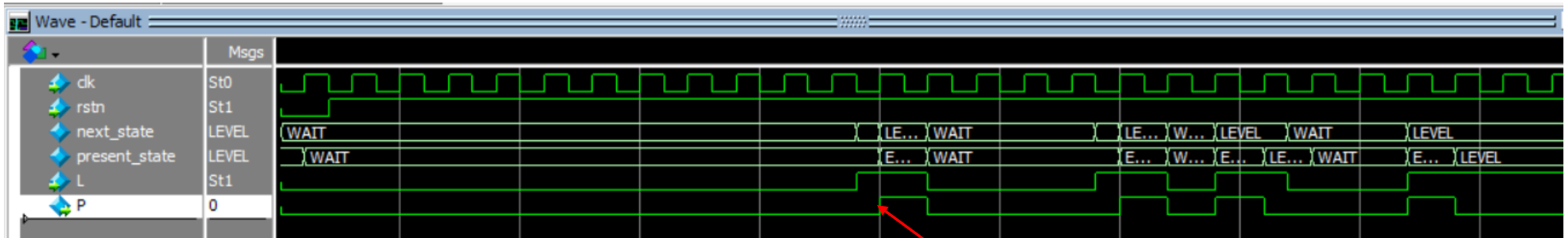


Mealy State Transition Table

	Source State	Destination State	Condition
1	EDGE	LEVEL	(L).(rstn)
2	EDGE	WAIT	(!L) + (L).(!rstn)
3	LEVEL	LEVEL	(L).(rstn)
4	LEVEL	WAIT	(!L) + (L).(!rstn)
5	WAIT	WAIT	(!L) + (L).(!rstn)
6	WAIT	EDGE	(L).(rstn)

2 Flipflops to represent three states WAIT EDGE, LEVEL

Simulation Waveform



Output **P** is single cycle pulse, available after sometime there is rising edge on input **L**

Level to Pulse Converter (Mealy FSM with Reduced States)

```
module level_to_pulse_converter_mealy(  
    input logic clk, rstn,  
    input logic L,  
    output logic P);
```

// FSM state encodings and state registers declaration

```
enum logic[1:0] WAIT=2'b00,  
                EDGE=2'b01} present_state, next_state;  
  
Only two states required for Mealy compared to Moore FSM
```

// Sequential Logic for present state

```
always_ff@(posedge clk) begin  
    if(!rstn)  
        present_state <= WAIT;  
    else  
        present_state <= next_state;  
end
```

(continued....)

// Combination Logic for Next State and Output

```
always_comb begin  
    case(present_state)  
        WAIT: begin  
            if(L==1) begin  
                next_state = EDGE; P = 1;  
            end  
            else begin  
                next_state = WAIT; P = 0;  
            end  
        end  
        EDGE: begin  
            if(L==1) begin  
                next_state = EDGE; P = 0;  
            end  
            else begin  
                next_state = WAIT; P = 0;  
            end  
        end  
        default: begin P = 0; next_state = WAIT; end  
    endcase  
end  
endmodule: level_to_pulse_converter_mealy
```

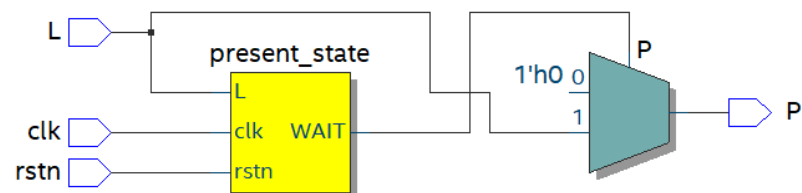
for next state combinational logic **always_comb** specified to automatically infer sensitivity list

Output **P** is set to '1' as soon as input L is detected as '1' in **WAIT** state

Since output **P** requirement is single cycle pulse, **P** is set to '0' if input stays at level '1' after rising edge detection

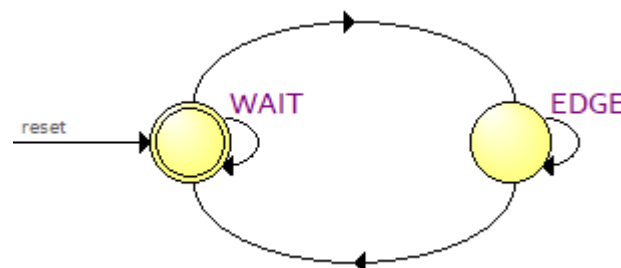
Level to Pulse Converter Simulation and Synthesis Results (Mealy FSM)

Post Synthesis RTL Netlist Schematic



Resource	Usage
Estimated ALUTs Used	2
-- Combinational ALUTs	2
-- Memory ALUTs	0
-- LUT_REGS	0
Dedicated logic registers	1

Moore State Diagram



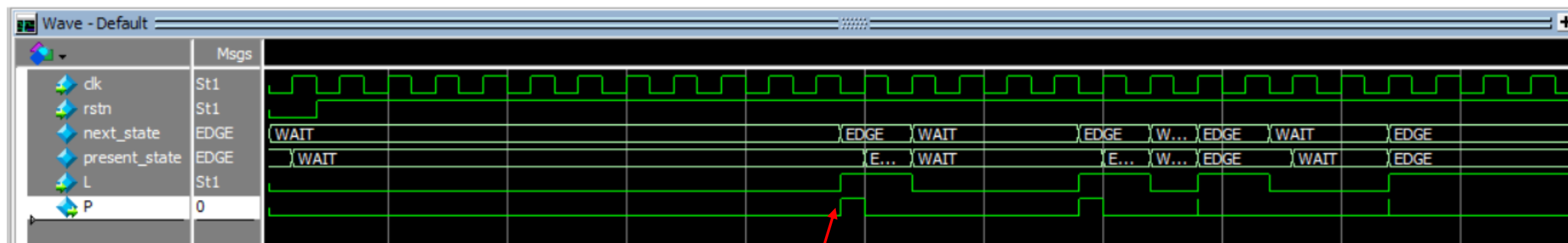
Mealy State Transition Table

	Source State	Destination State	Condition
1	EDGE	EDGE	(L).(rstn)
2	EDGE	WAIT	(!L) + (L).(!rstn)
3	WAIT	EDGE	(L).(rstn)
4	WAIT	WAIT	(!L) + (L).(!rstn)

1 Flipflop to represent two states WAIT and EDGE.

1 less D-flipflop required in Mealy compared to Moore FSM

Simulation Waveform




Output **P** is available as soon as input **L** changed from '0' to '1'. However output **P** pulse is not stable for 1 cycle. Which does not meet design requirement for **P** to be 1 cycle pulse !!

Level to Pulse Converter (Mealy FSM with Reduced States + Registered output)

```
module level_to_pulse_converter_mealy(  
    input logic clk, rstn,  
    input logic L,  
    output logic P);
```

// FSM state encodings and state registers declaration

```
enum logic[1:0] {WAIT=2'b00,  
                 EDGE=2'b01} present_state, next_state;
```

Only two states required for Mealy compared to Moore FSM

```
logic r_P; // local variable declaration
```

// Sequential Logic for present state

```
always_ff@(posedge clk) begin
```

```
    if(!rstn) begin
```


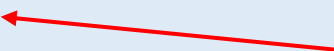
```
        present_state <= WAIT;
```

```
        P <= 0;
```

```
    end
```

```
    else begin
```

```
        present_state <= next_state;
```

```
        P <= r_P; 
```

```
    end
```

```
end
```

(continued....)

Synthesizer will create D-flipflop to provide registered output P.

Output P is registered to ensure P is at least 1 cycle pulse


// Combination Logic for Next State and Output

```
always_comb begin
```

```
    case(present_state)
```


```
        WAIT: begin
```

```
            if(L==1) begin
```

```
                next_state = EDGE; r_P = 1;
```

```
            end
```

```
        else begin
```

```
            next_state = WAIT; r_P = 0;
```

```
        end
```

```
    end
```


```
    EDGE: begin
```

```
        if(L==1) begin
```

```
            next_state = EDGE; r_P = 0;
```


```
        end
```

```
    else begin
```

```
        next_state = WAIT; r_P = 0;
```

```
    end
```

```
end
```

```
default: begin r_P = 0; next_state = WAIT; end
```

```
endcase
```

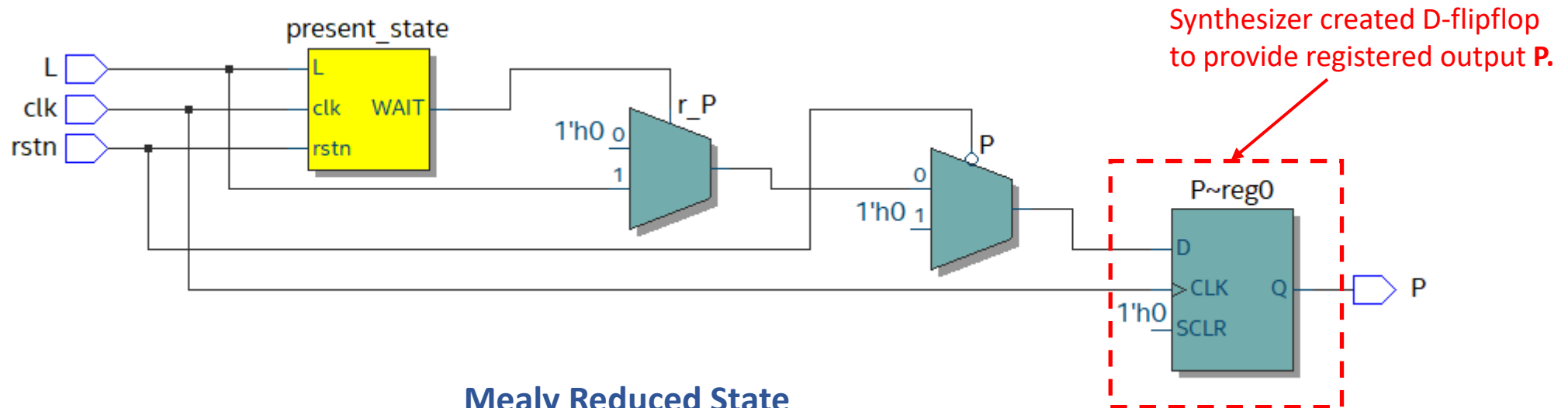
```
end
```

```
endmodule: level_to_pulse_converter_mealy
```

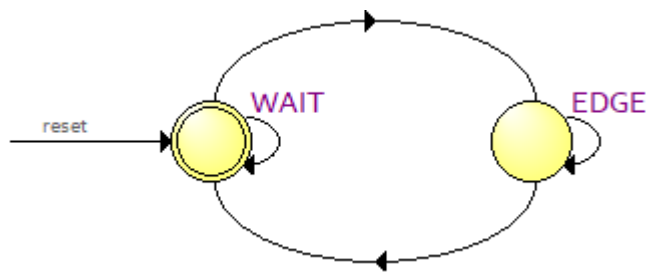
Output is assigned to local variable r_P

Level to Pulse Converter Synthesis Results (Mealy FSM) : With Registered Output

Post Synthesis RTL Netlist Schematic



Mealy Reduced State Diagram



Only 2 states in Mealy FSM compared to Moore which as 3 States

Mealy Reduced State Transition Table

	Source State	Destination State	Condition
1	EDGE	WAIT	(!L) + (L).(!rstn)
2	EDGE	EDGE	(L).(rstn)
3	WAIT	WAIT	(!L) + (L).(!rstn)
4	WAIT	EDGE	(L).(rstn)

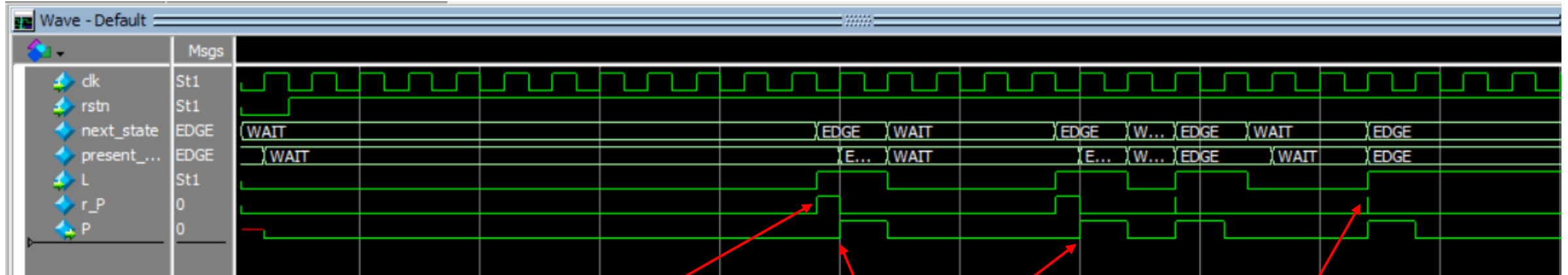
Post Synthesis Resource Utilization

Resource	Usage
Estimated ALUTs Used	2
-- Combinational ALUTs	2
-- Memory ALUTs	0
-- LUT_REGS	0
Dedicated logic registers	2

1 Flipflop for registering output P and 1 Flipflop for representing two states WAIT and EDGE

Level to Pulse Converter Simulation Result(Mealy FSM) : With Registered Output

Post Synthesis RTL Netlist Schematic



Unregistered output **r_P** is available as soon as input **L** changed from '0' to '1'

Registered Output **P** which is single cycle pulse when there is rising edge on **L** detected.

Glitches in Simulation on unregistered output `r_P`

References

- ❑ For FSM design and synthesis coding guidelines refer to below mentioned white paper :
 - http://www.sunburst-design.com/papers/CummingsSNUG1998SJ_FSM.pdf