

Machine Learning for Signal Processing

Independent Component Analysis

Instructor: Bhiksha Raj

A note on bits..

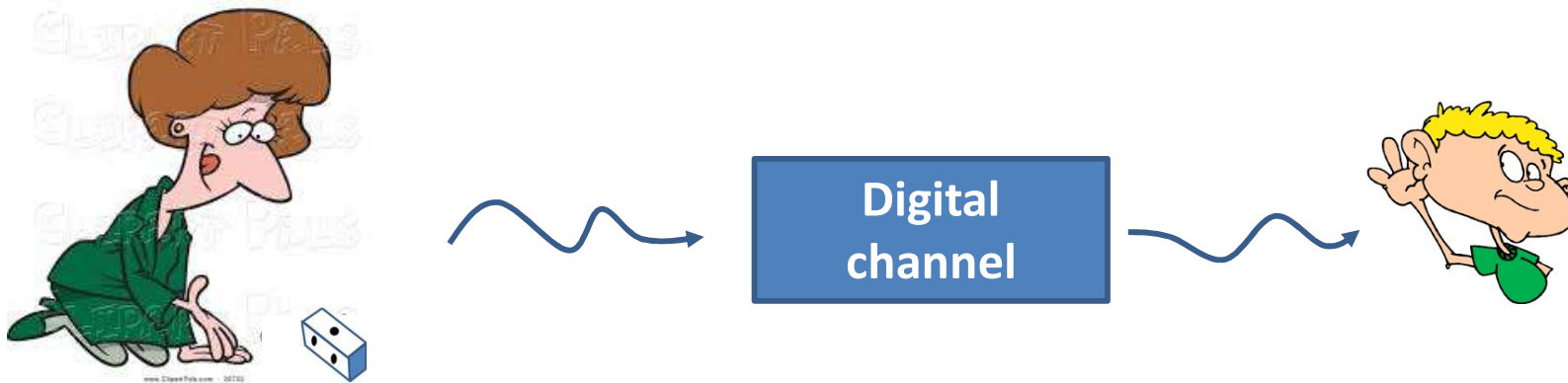
- You flip a coin. You must inform your friend in the next room about whether the outcome was heads or tails



- How many bits will you have to send?

A note on bits..

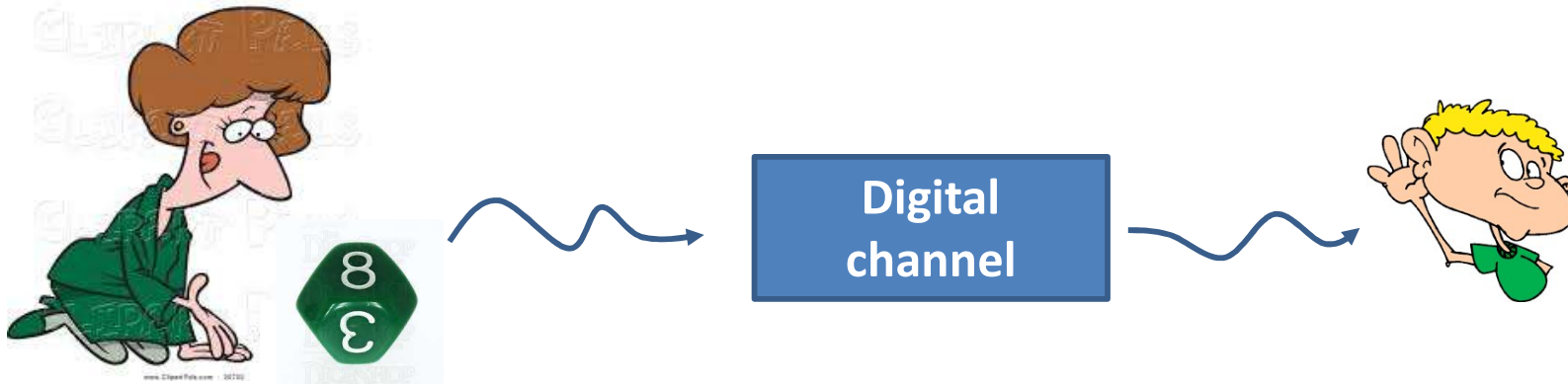
- You roll a four-side dice. You must inform your friend in the next room about the outcome



- How many bits will you have to send?

A note on bits..

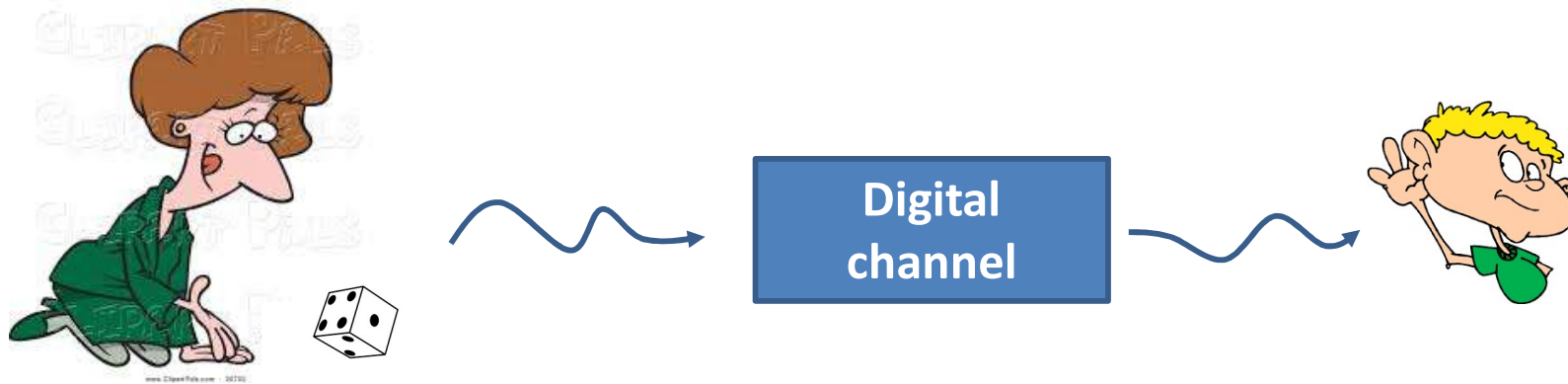
- You roll an *eight-sided octahedral* dice. You must inform your friend in the next room about the outcome



- How many bits will you have to send?

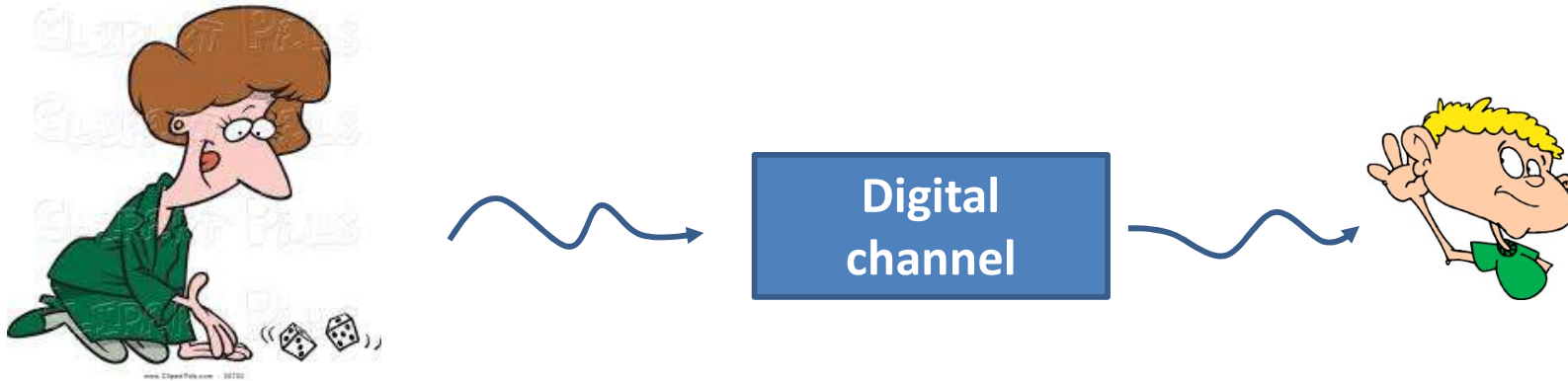
A note on bits..

- You roll a *six-sided* dice. You must inform your friend in the next room about the outcome



- How many bits will you have to send?

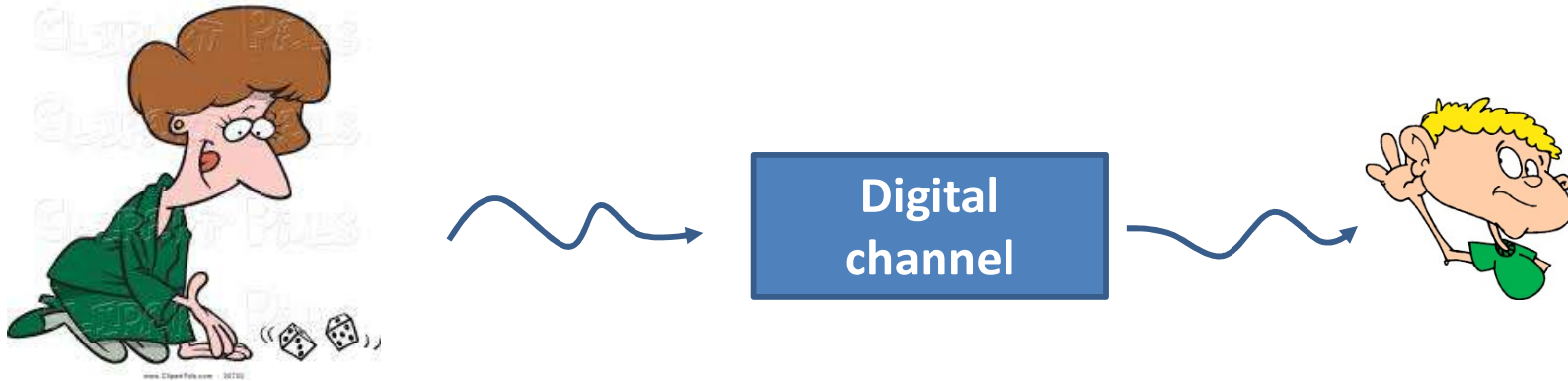
Batching up 6-sided dice rolls



- Instead of sending individual rolls, you roll the dice *twice*
 - And send the *pair* to your friend
- How many bits do you send *per roll*?

Roll 1	Roll 2
1	1
1	2
1	3
..	..
2	1
2	2
..	..
6	6

Batching up 6-sided dice rolls

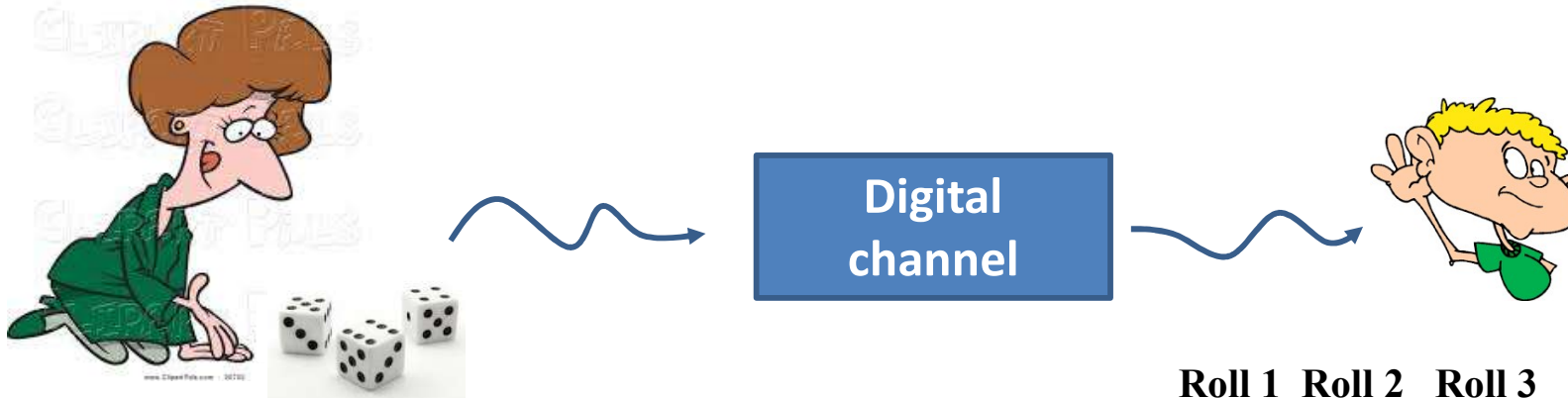


- Instead of sending individual rolls, you roll the dice *twice*
 - And send the *pair* to your friend
- How many bits do you send *per roll*?
- 36 combinations: 6 bits per pair of numbers
 - Still 3 bits per roll

Roll 1 **Roll 2**

1	1
1	2
1	3
..	..
2	1
2	2
..	..
6	6

Batching up 6-sided dice rolls



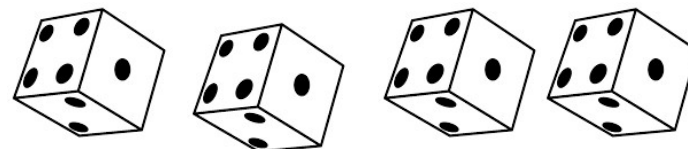
- Instead of sending individual rolls, you roll the dice ***three times***
 - And send the *triple* to your friend
- How many bits do you send *per roll*?
- 216 combinations: 8 bits per triple
 - Still 2.666 bits per roll
 - *Now we're talking!*

Roll 1	Roll 2	Roll 3
1	1	1
1	1	2
..
1	6	3
..		..
2	1	1
2	1	2
..		..
6	6	6

Batching up 6-sided dice rolls

- Batching *four rolls*

- 1296 combinations
- 11 bits per outcome (4 rolls)
- 2.75 bit per roll

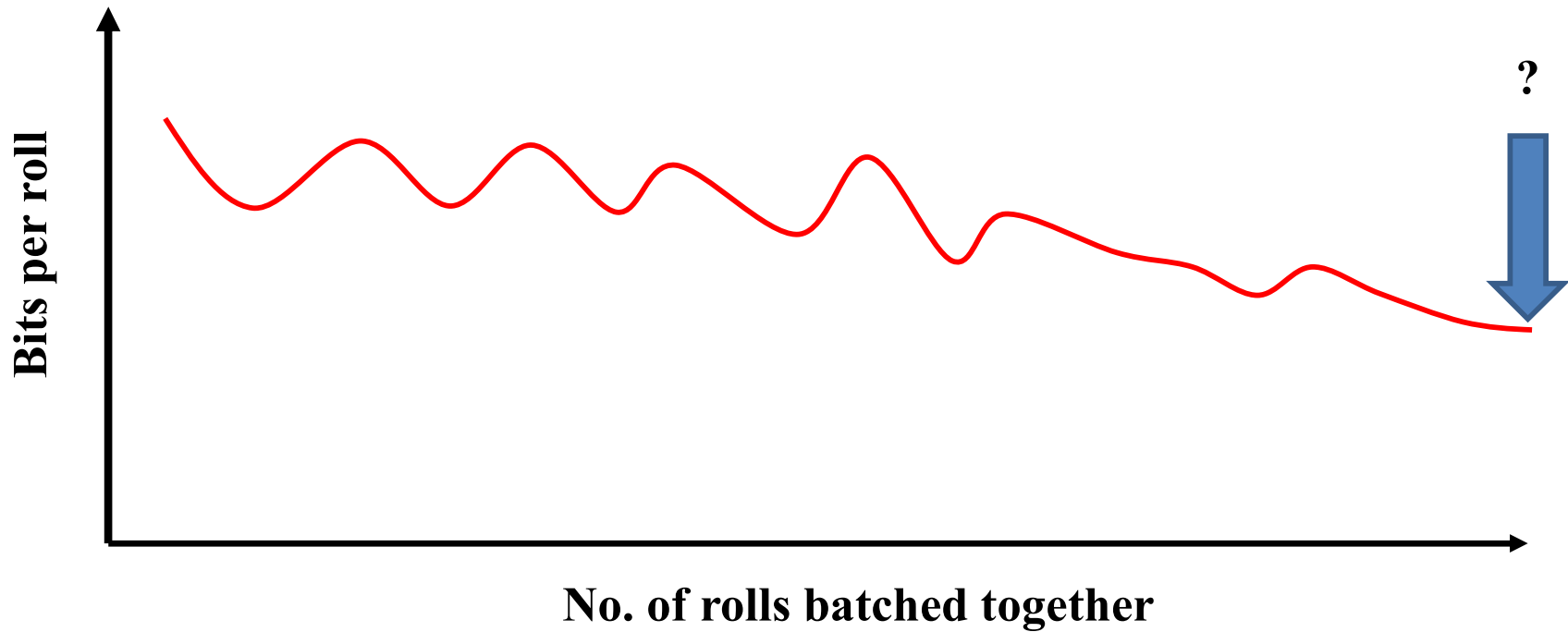


- Batching *five rolls*

- 7776 combinations
- 13 bits per outcome (5 rolls)
- 2.6 bits per roll

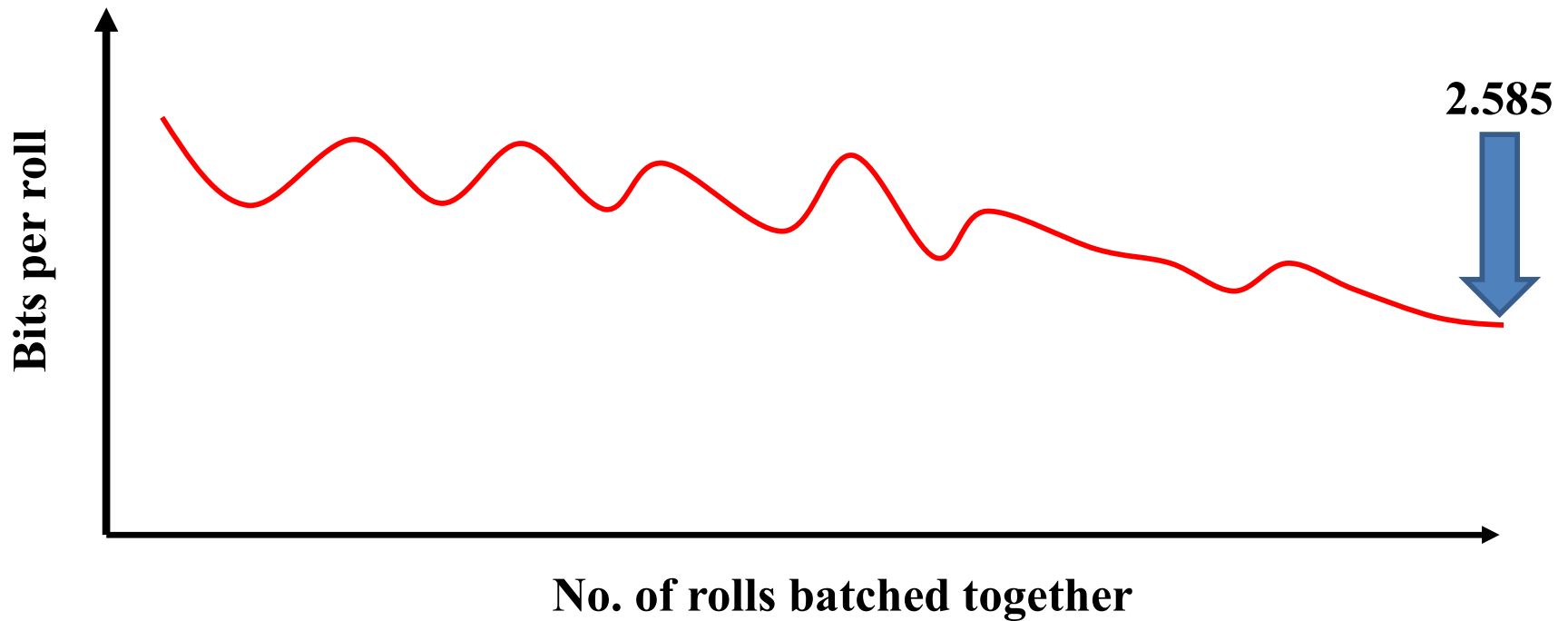


Batching up 6-sided dice rolls



- Where will it end?

Batching up 6-sided dice rolls



- Where will it end?
- $\lim_{k \rightarrow \infty} \frac{\lceil k \log_2(6) \rceil}{k} = \log_2(6)$ bits per roll in the limit
 - This is the absolute minimum – no simple batching will give you less than these many bits per outcome with this scheme

Poll 1

- The number of bits needed to send an individual outcome of the roll of an N-sided dice is $\log_2(N)$
 - True
 - False
- If we batch many outcomes (of the roll of an N-sided dice) together and transmit them, the average number of bits needed to per outcome tends to $\log_2(N)$ as the size of the batch increases to infinity
 - True
 - False

Poll 1

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 - True
 - **False**
- If we batch many outcomes (of the roll of an N-sided dice) together and transmit them, the average number of bits needed to per outcome tends to $\log_2(N)$ as the size of the batch increases to infinity
 - **True**
 - False

Can we do better?

- A four-sided die needs 2 bits per roll
- But then you find not all sides are equally likely



- $P(1) = 0.5, P(2) = 0.25, P(3) = 0.125, P(4) = 0.125$
- *Can you do better than 2 bits per outcome*

Can we do better?

- You have

$$P(1) = 0.5, P(2) = 0.25, P(3) = 0.125, P(4) = 0.125$$

1	0
2	1 0
3	1 1 0
4	1 1 1

- You use:

– Note receiver is *never in any doubt as to what they received*

- What is the average number of bits per outcome

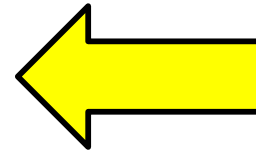
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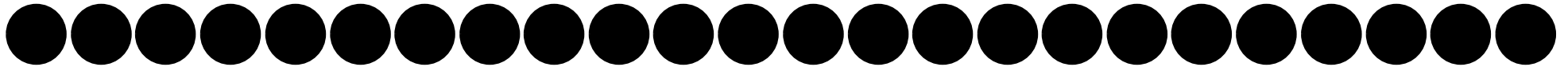
- You use:

1	0
2	1 0
3	1 1 0
4	1 1 1



- Note receiver is *never in any doubt as to what they received*
- How did we know to use three bits here for rows 3 and 4, 2 for row 2 and 1 for row 1?

In a loooong sequence of trials...



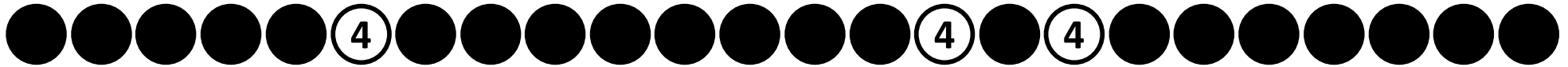
- What fraction of these trials will be “4”?
 - $P(4) = 0.125$

In a loooong sequence of trials...



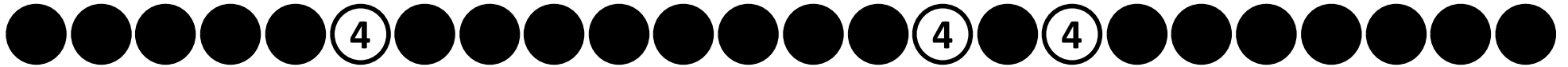
- What fraction of these trials will be “4”?
 - $P(4) = 0.125$
- From how many alternatives (on average) do we choose the 4
 - From the local perspective of 4

In a loooong sequence of trials...



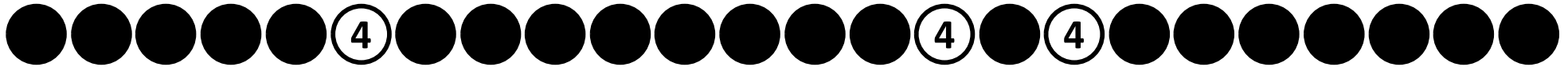
- What fraction of these trials will be “4”?
 - $P(4) = 0.125$
- From how many alternatives (on average) do we choose the 4
 - From the perspective of 4, you might as well have been rolling an *eight-sided dice*

In a loooong sequence of trials...



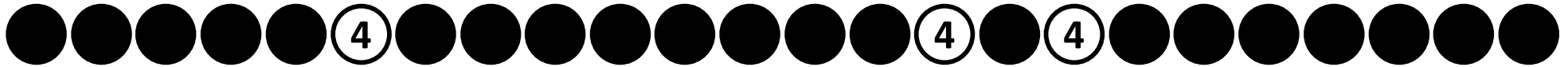
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 - From the perspective of 4 you might as well have been rolling an eight-sided dice
- How many bits to code each instance of 4?
 - When 4 is the outcome of rolls of an 8-sided dice

In a loooong sequence of trials...



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 - When 4 is the outcome of rolls of an 8-sided dice
- What is the average (expected) number of bits to transmit all instances of 4 in N rolls of the dice?

In a loooong sequence of trials...



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- How many bits to code each instance of 4?
 - When 4 is the outcome of rolls of an 8-sided dice
- What is the average (expected) number of bits to transmit all instances of 4 in N rolls of the dice?
 - Average per roll?

In a loooong sequence of trials...



- What fraction of these trials will be “1”?
 - $P(1) = 0.5$

In a loooong sequence of trials...



- What fraction of these trials will be “1”?
 - $P(1) = 0.5$
- From how many alternatives (on average) do we choose the 1
 - From the local perspective of 1

In a loooong sequence of trials...



- What fraction of these trials will be “1”?
 - $P(1) = 0.5$
- From how many alternatives (on average) do we choose the 1
 - From the perspective of 1, you might as well have been flipping a coin

In a loooong sequence of trials...



- What fraction of these trials will be “1”?
 - $P(1) = 0.5$
- From how many alternatives (on average) do we choose the 1
 - From the perspective of 1 you might as well have been flipping a coin
- How many bits to code each instance of 1?
 - When 1 is the outcome of a coin toss

In a loooong sequence of trials...



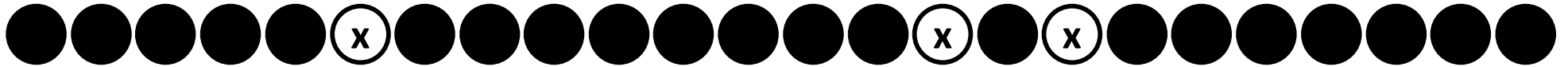
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In a loooong sequence of trials...



- What fraction of these trials will be “1”?
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 - When 1 is the outcome of rolls of a coin toss
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 - Average per roll?

In a loooong sequence of trials...

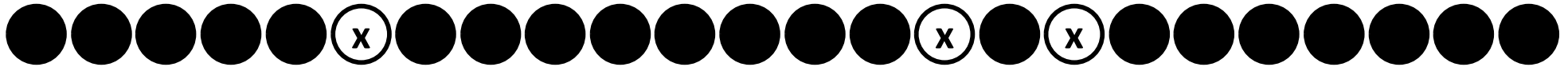


- An outcome x has probability $P(x)$
- From the perspective of x , how many-sided dice is it an outcome of?

In a loooong sequence of trials...

$$\frac{1}{P(x)}$$

$$\frac{1}{P(x)} \quad \frac{1}{P(x)}$$



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- From the perspective of x , how many-sided dice is it an outcome of?
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In a loooong sequence of trials...

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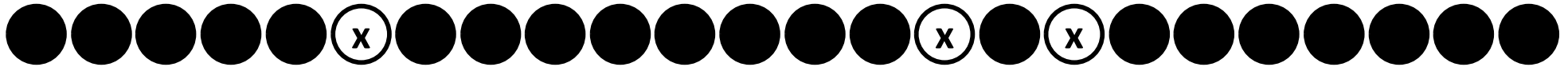


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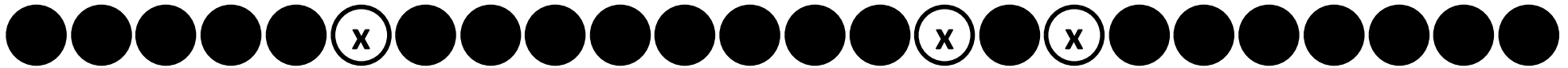


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- What is the average (expected) number of bits to transmit instances of x in N rolls of the dice?
- Expected number of bits per outcome for *any* outcome?

In a loooong sequence of trials...

$$\frac{1}{P(x)}$$

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- An outcome x has probability $P(x)$
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- How many bits to code an instance of x ?
- What is the average (expected) number of bits to transmit instances of x in N rolls of the dice?
- Expected number of bits per outcome for *any* outcome?
- Average per trial?

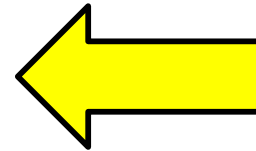
How we do better...

- You have

$$P(1) = 0.5, P(2) = 0.25, P(3) = 0.125, P(4) = 0.125$$

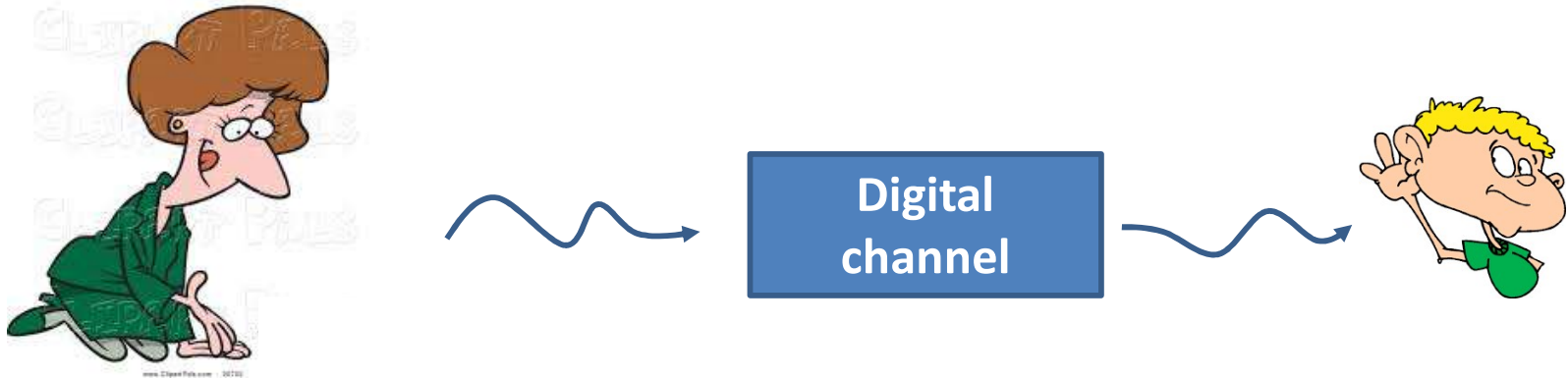
- You use:

1	0
2	1 0
3	1 1 0
4	1 1 1



- Note receiver is *never in any doubt as to what they received*
- An outcome with probability p is equivalent to obtaining one of $1/p$ equally likely choices
 - Requires $\log_2 \left(\frac{1}{p} \right)$ bits on average

Entropy



- The average number of bits per symbol required to communicate a random variable over a digital channel *using an optimal code* is

$$H(p) = \sum_i p_i \log \frac{1}{p_i} = - \sum_i p_i \log p_i$$

- You can't do better
 - Any other code will require more bits
- This is the *entropy of the random variable*

Poll 2

- Mark the true statements about transmitting the outcomes of draws from a distribution
 - The entropy of the distribution is the number of bits needed to transmit the outcome of a single draw
 - The entropy of the distribution is the average number of bits needed to transmit an outcome, when we batch infinite outcomes together

Poll 2

- Mark the true statements about transmitting the outcomes of draws from a distribution
 - The entropy of the distribution is the number of bits needed to transmit the outcome of a single draw
 - **The entropy of the distribution is the average number of bits needed to transmit an outcome, when we batch infinite outcomes together**

A brief review of basic info. theory



T(all), M(ed), S(hort)...

$$H(X) = \sum_X P(X)[- \log P(X)]$$

- Entropy: The *minimum average* number of bits to transmit to convey a symbol



T, M, S...

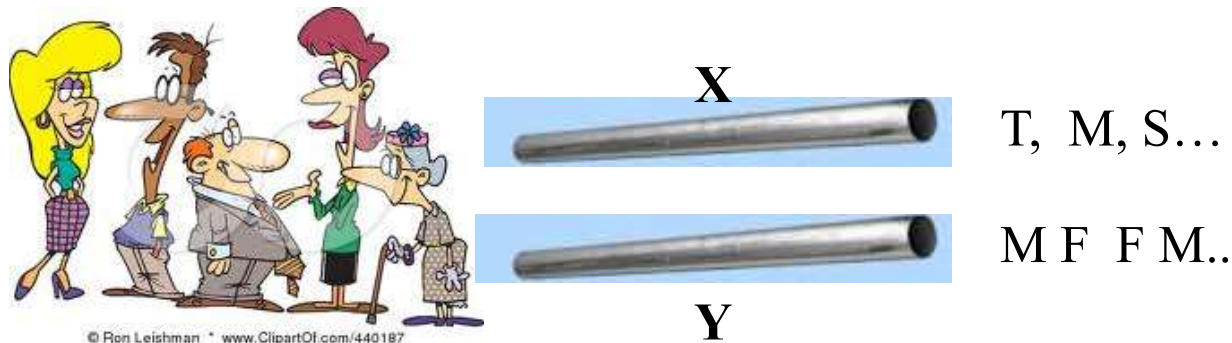


M F F M..

$$H(X, Y) = \sum_{X, Y} P(X, Y)[- \log P(X, Y)]$$

- Joint entropy: The *minimum average* number of bits to convey sets (pairs here) of symbols

A brief review of basic info. theory



$$H(X | Y) = \sum_Y P(Y) \sum_X P(X | Y) [-\log P(X | Y)] = \sum_{X,Y} P(X, Y) [-\log P(X | Y)]$$

- Conditional Entropy: The *minimum average* number of bits to transmit to convey a symbol X , after symbol Y has already been conveyed
 - Averaged over all values of X and Y

And now
for something
completely different...



The statistical concept of correlatedness

- Two variables X and Y are correlated if knowing X gives you an *expected* value of Y
- X and Y are uncorrelated if knowing X tells you nothing about the *expected* value of Y
 - Although it could give you other information
 - How?

Correlation vs. Causation

- The consumption of burgers has gone up steadily in the past decade



- In the same period, the penguin population of Antarctica has gone down

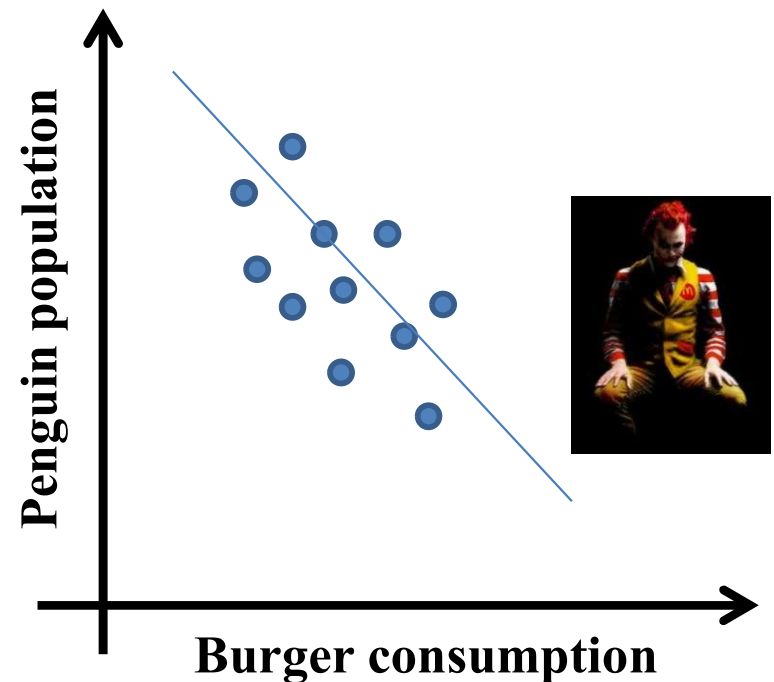
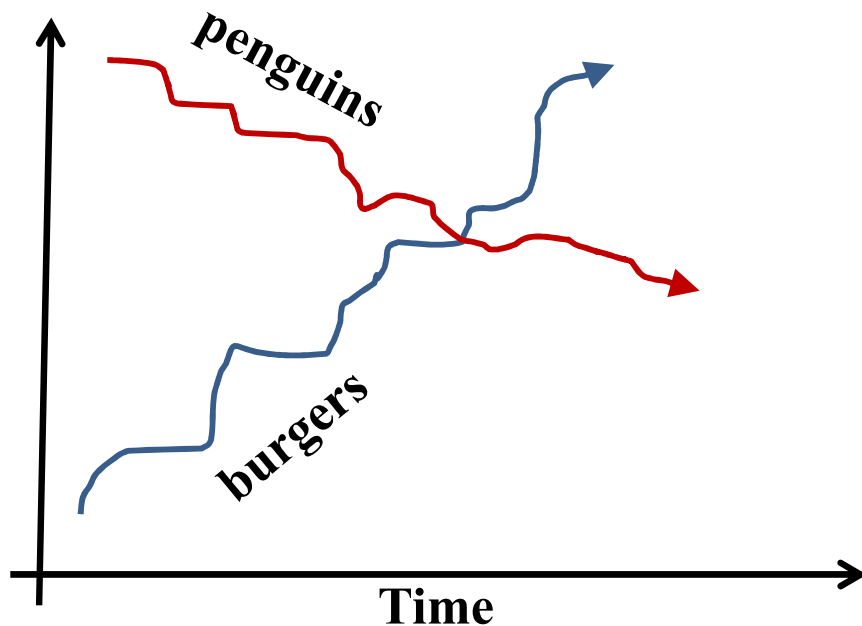


Correlation, not Causation
(unless McDonalds has a
top-secret Antarctica division)



The concept of *correlation*

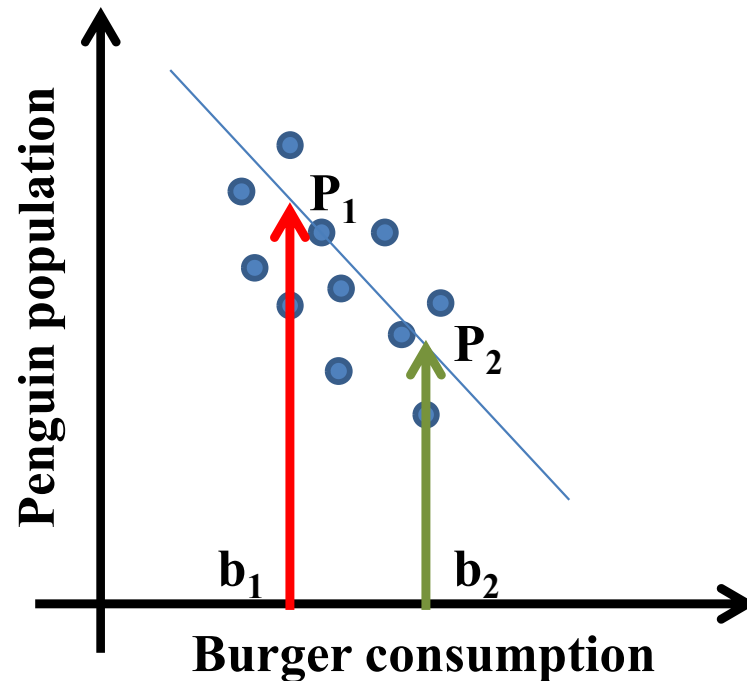
- Two variables are correlated if knowing the value of one gives you information about the ***expected value*** of the other



A brief review of basic probability

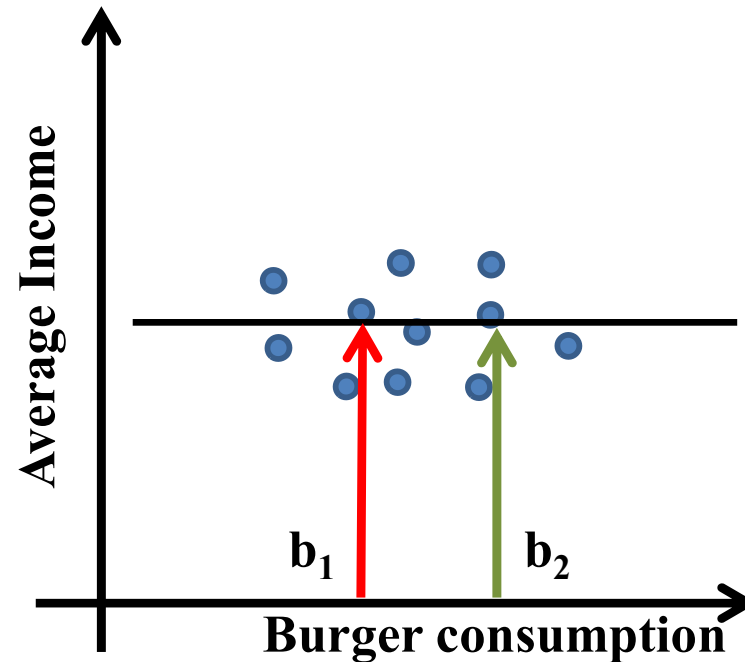
- *Uncorrelated*: Two random variables X and Y are uncorrelated iff:
 - The *average* value of the product of the variables equals the product of their individual averages
- Setup: Each draw produces one instance of X and one instance of Y
 - I.e one instance of (X, Y)
- $E[XY] = E[X]E[Y]$
- The average value of Y is the same regardless of the value of X

Correlated Variables



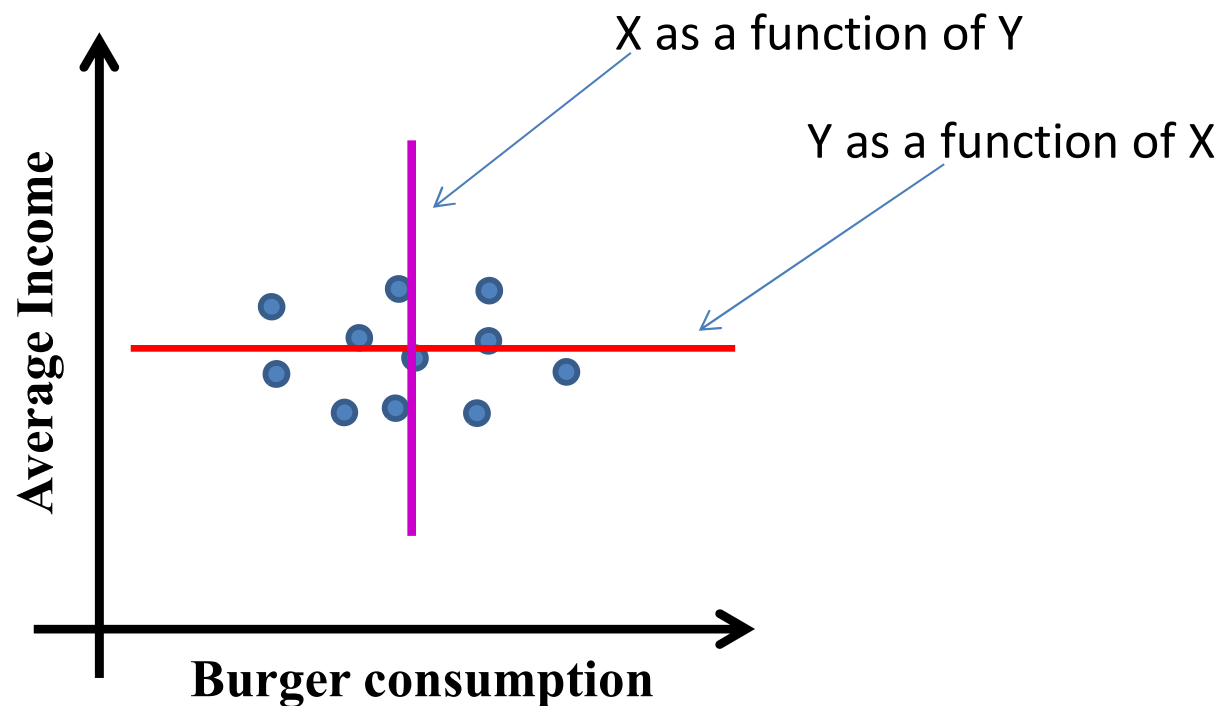
- Expected value of Y given X :
 - Find average of Y values of all samples at (or close) to the given X
 - If this is a function of X , X and Y are correlated

Uncorrelatedness



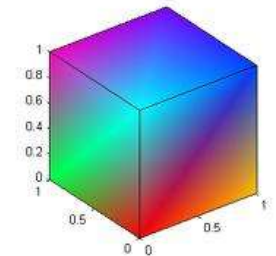
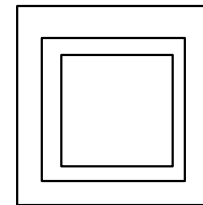
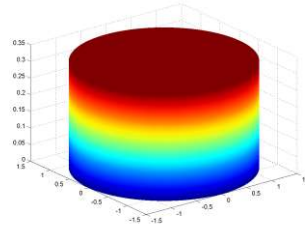
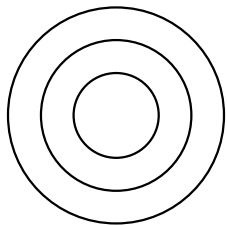
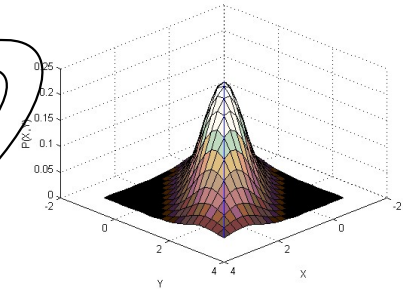
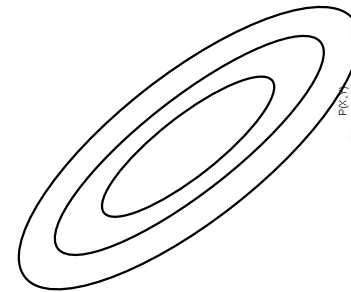
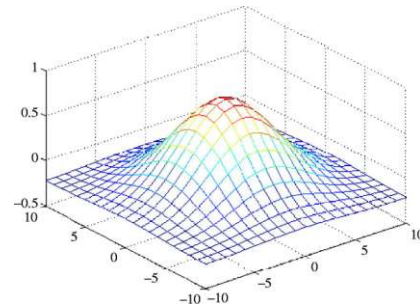
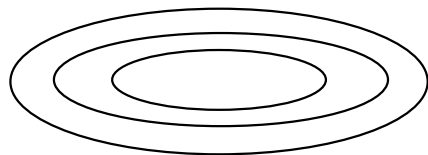
- Knowing X does not tell you what the *average* value of Y is
 - And vice versa

Uncorrelated Variables



- The average value of Y is the same regardless of the value of X and vice versa

Uncorrelatedness in Random Variables

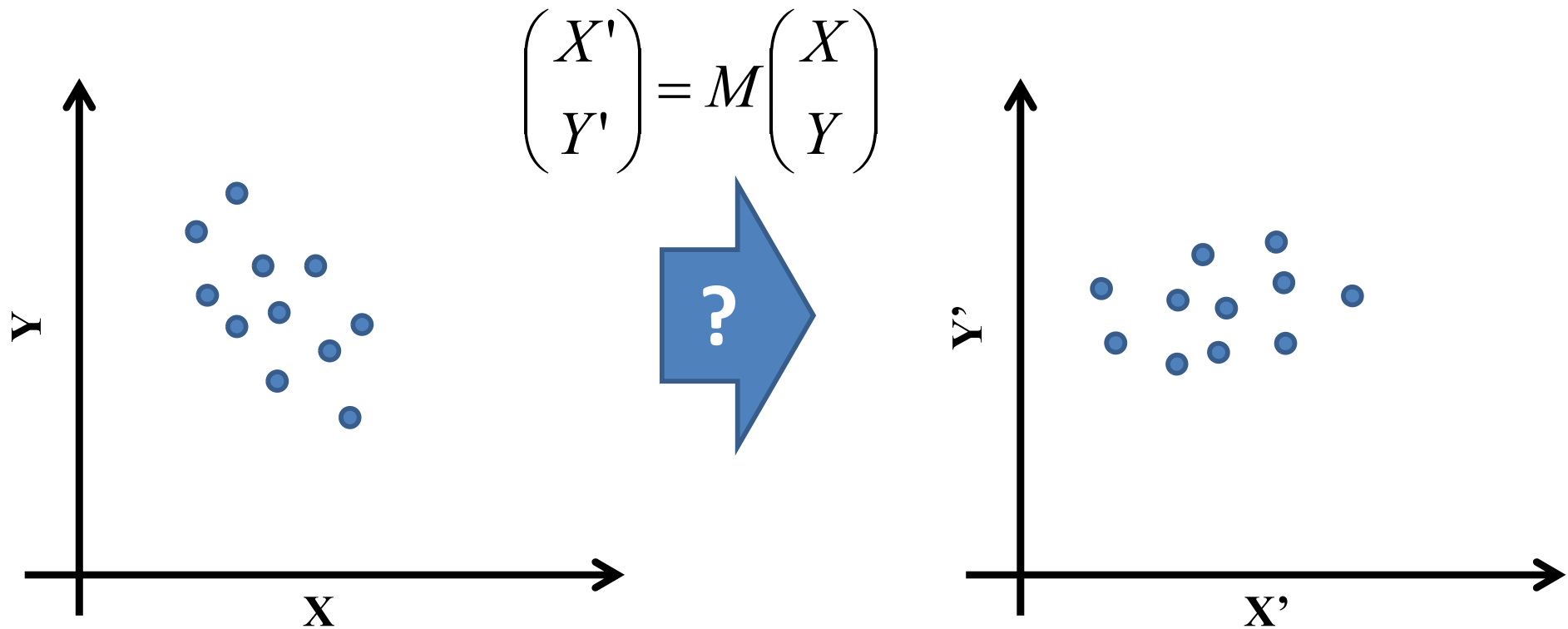


- Which of the above represent uncorrelated RVs?

Benefits of uncorrelatedness..

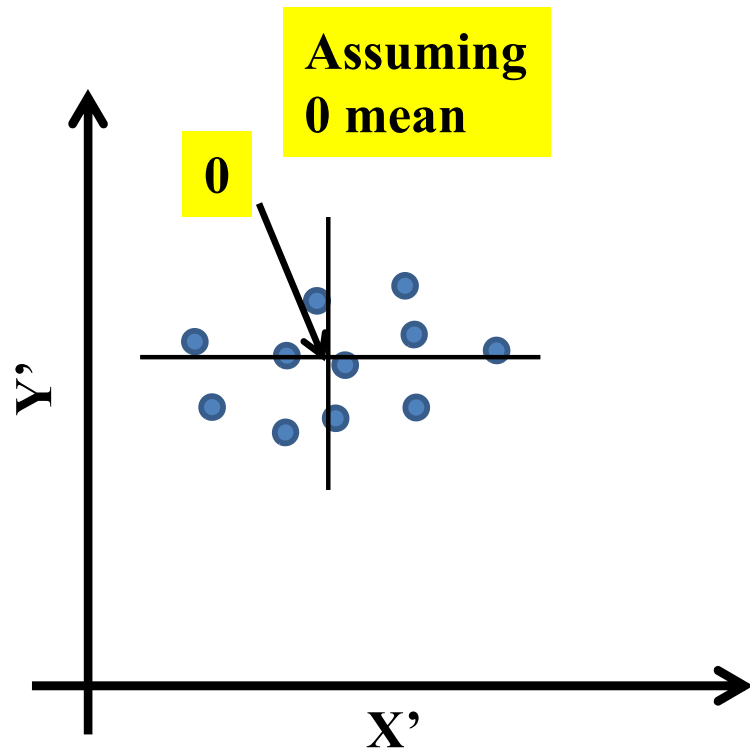
- Uncorrelatedness of variables is generally considered desirable for modelling and analyses
 - For Euclidean error based regression models and probabilistic models, uncorrelated variables can be separately handled
 - Since the value of one doesn't affect the average value of others
 - Greatly reduces the number of model parameters
 - Otherwise their interactions must be considered
- We will frequently transform correlated variables to make them uncorrelated
 - “Decorrelating” variables

The notion of *decorrelation*



- So how does one transform the correlated variables (X, Y) to the uncorrelated (X', Y')

What does “uncorrelated” mean



- $E[X'] = \text{constant}$
- $E[Y'] = \text{constant}$
- $E[Y'|X'] = \text{constant}$
- $E[X'Y'] = E[X']E[Y']$
- All will be 0 for centered data

$$E \left[\begin{pmatrix} X' \\ Y' \end{pmatrix} (X' \ Y') \right] = E \begin{pmatrix} X'^2 & X'Y' \\ X'Y' & Y'^2 \end{pmatrix} = \begin{pmatrix} E[X'^2] & 0 \\ 0 & E[Y'^2] \end{pmatrix} = \text{diagonal matrix}$$

- If \mathbf{Y} is a matrix of vectors, $\mathbf{Y}\mathbf{Y}^T = \text{diagonal}$

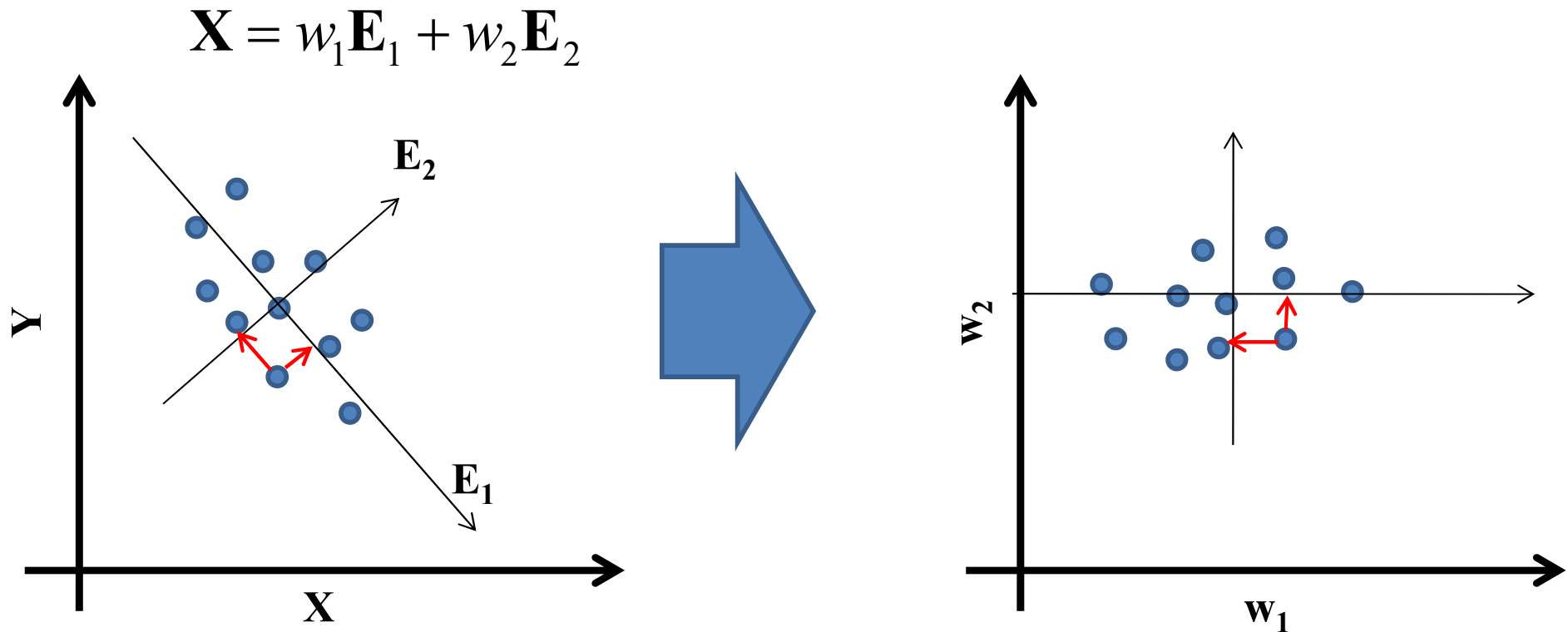
Decorrelation

- Let \mathbf{X} be the matrix of correlated data vectors
 - Each component of \mathbf{X} informs us of the mean trend of other components
- Need a transform \mathbf{M} such that if $\mathbf{Y} = \mathbf{MX}$ such that the covariance of \mathbf{Y} is diagonal
 - \mathbf{YY}^T is the covariance if \mathbf{Y} is zero mean
 - For uncorrelated components, $\mathbf{YY}^T = \mathbf{Diagonal}$
 - $\Rightarrow \mathbf{MXX}^T\mathbf{M}^T = \mathbf{Diagonal}$
 - $\Rightarrow \mathbf{M.Cov}(\mathbf{X}).\mathbf{M}^T = \mathbf{Diagonal}$

Decorrelation

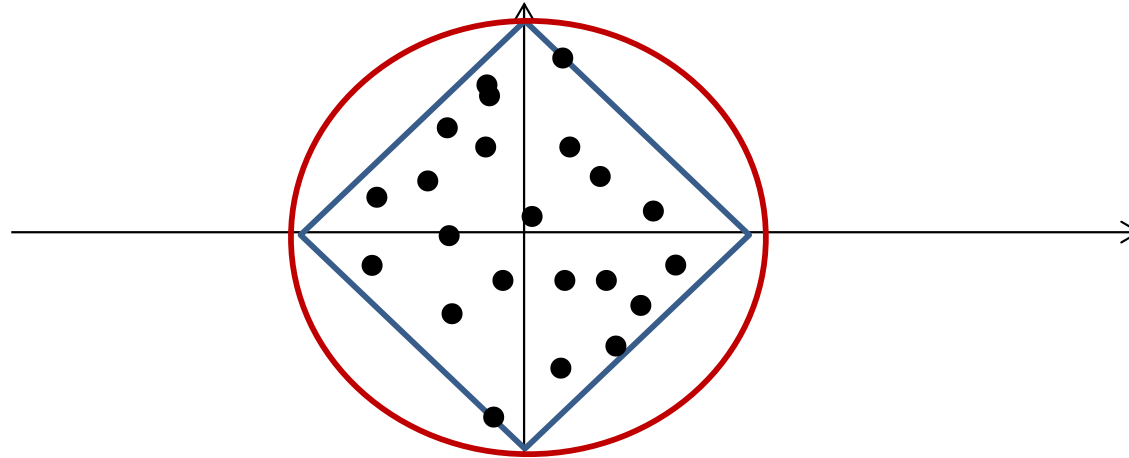
- Easy solution:
 - Eigen decomposition of $\text{Cov}(\mathbf{X})$:
$$\text{Cov}(\mathbf{X}) = \mathbf{E}\mathbf{\Lambda}\mathbf{E}^T$$
 - $\mathbf{E}\mathbf{E}^T = \mathbf{I}$
- Let $\mathbf{M} = \mathbf{E}^T$
- $\mathbf{M}\text{Cov}(\mathbf{X})\mathbf{M}^T = \mathbf{E}^T\mathbf{E}\mathbf{\Lambda}\mathbf{E}^T\mathbf{E} = \mathbf{\Lambda} = \text{diagonal}$
- **PCA: $\mathbf{Y} = \mathbf{E}^T\mathbf{X}$**
 - Projects the data onto the Eigen vectors of the covariance matrix
 - *Diagonalizes* the covariance matrix
 - “Decorrelates” the data

PCA



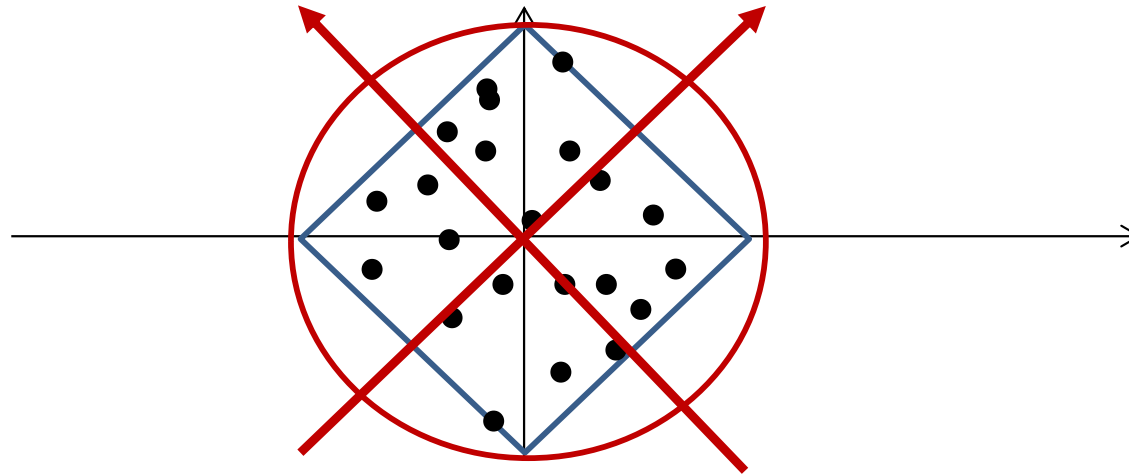
- PCA: $\mathbf{Y} = \mathbf{E}^T \mathbf{X}$
 - Projects the data onto the Eigen vectors of the covariance matrix
 - Changes the coordinate system to the Eigen vectors of the covariance matrix
 - *Diagonalizes* the covariance matrix
 - “Decorrelates” the data

Decorrelating the data



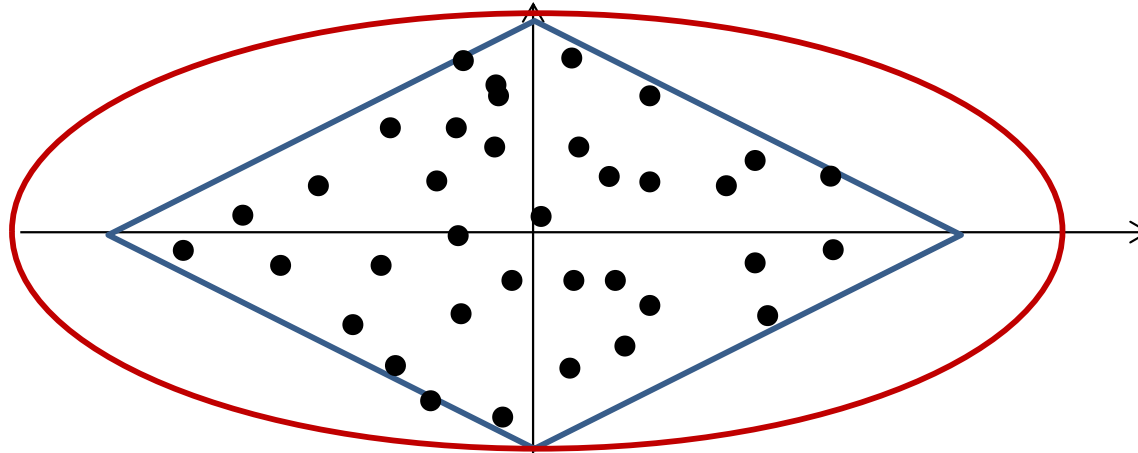
- Are there other decorrelating axes?

Decorrelating the data



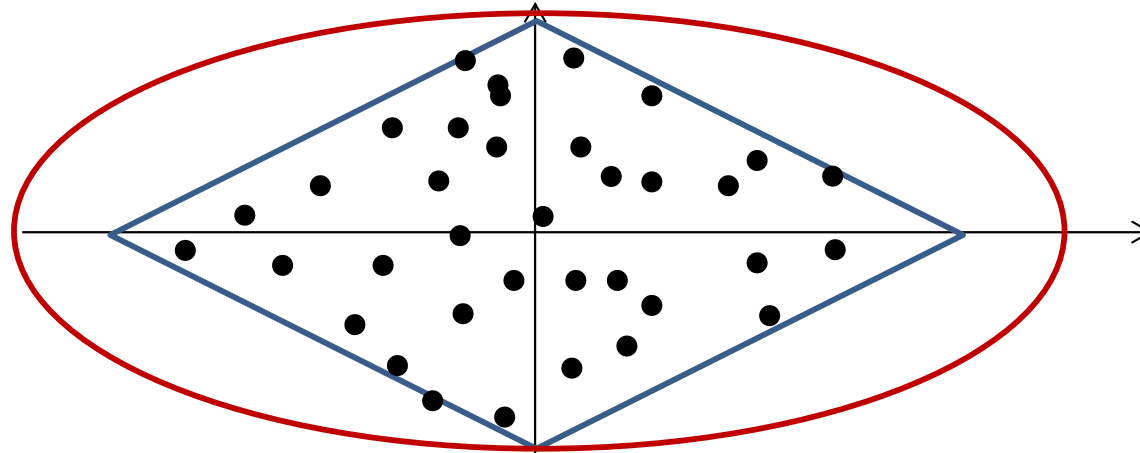
- Are there other decorrelating axes?

Decorrelating the data



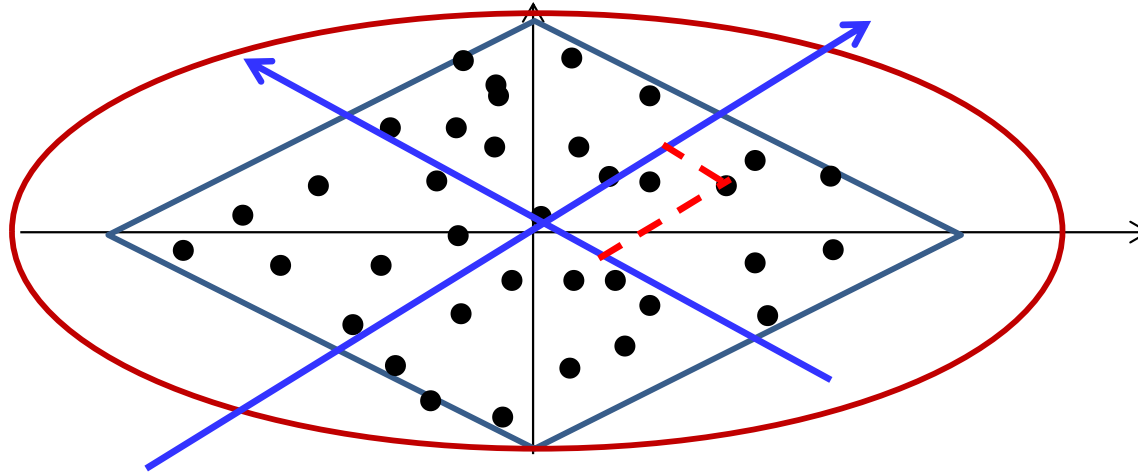
- Are there other decorrelating axes?

Decorrelating the data



- Are there other decorrelating axes?
- What about if we don't require them to be orthogonal?

Decorrelating the data



- Are there other decorrelating axes?
- What about if we don't require them to be orthogonal?
- What is special about these axes?

Poll 3

- Mark all true statements about decorrelation of the components of a vector
 - There is a unique set of orthogonal bases along which the components of a data are decorrelated
 - There may be many different sets of orthogonal bases along which the components of the data are decorrelated
 - The bases along which the components of the data are decorrelated are always orthogonal

Poll 3

- Mark all true statements about decorrelation of the components of a vector
 - There is a unique set of orthogonal bases along which the components of a data are decorrelated
 - **There may be many different sets of orthogonal bases along which the components of the data are decorrelated**
 - The bases along which the components of the data are decorrelated are always orthogonal

The statistical concept of *Independence*

- Two variables X and Y are *dependent* if knowing X gives you *any information about* Y
- X and Y are *independent* if knowing X tells you nothing at all of Y

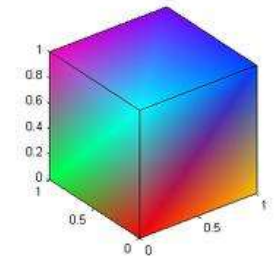
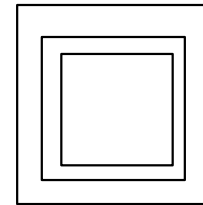
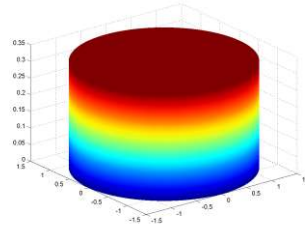
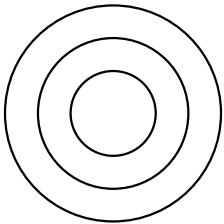
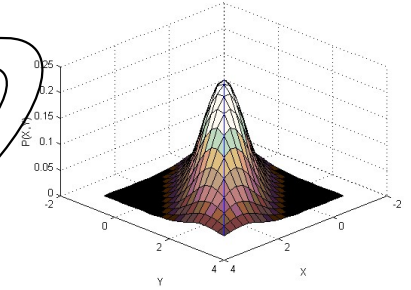
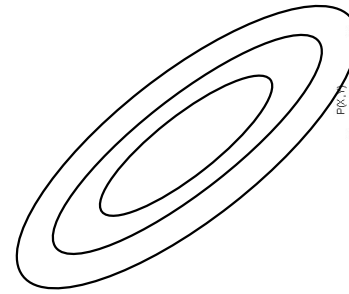
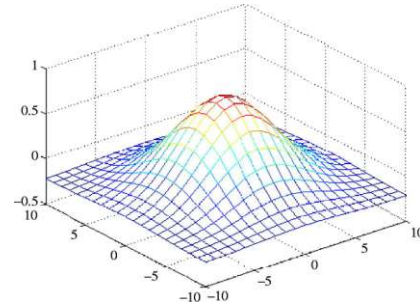
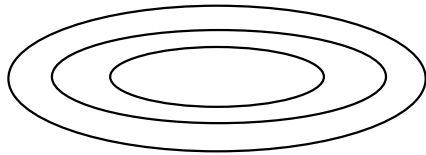
A brief review of basic probability

- ***Independence***: Two random variables X and Y are independent iff:
 - Their joint probability equals the product of their individual probabilities
- $P(X, Y) = P(X)P(Y)$
- Independence implies uncorrelatedness
 - The average value of X is the same regardless of the value of Y
 - $E[X|Y] = E[X]$
 - But uncorrelatedness does not imply independence

A brief review of basic probability

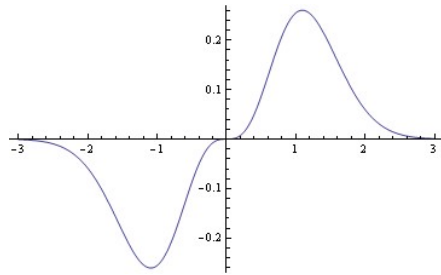
- *Independence*: Two random variables X and Y are independent iff:
- The average value of **any function** of X is the same regardless of the value of Y
 - Or any function of Y
- **$E[f(X)g(Y)] = E[f(X)] E[g(Y)]$ for all $f()$, $g()$**

Independence

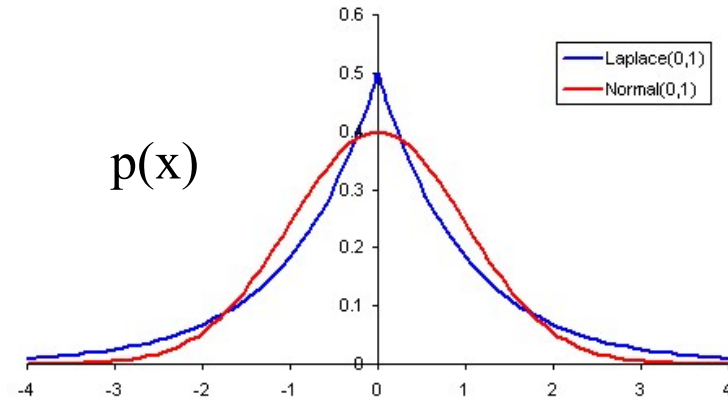
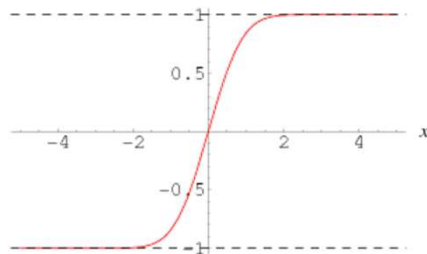


- Which of the above represent independent RVs?
- Which represent uncorrelated RVs?

A brief review of basic probability



$$y = f(x)$$



- The expected value of an odd function of an RV is 0 if
 - The RV is 0 mean
 - The PDF of the RV is symmetric around 0
- **$E[f(X)] = 0$ if $f(X)$ is odd symmetric**

A brief review of basic info. theory

- Conditional entropy of $X|Y = H(X)$ if X is independent of Y

$$H(X|Y) = \sum_Y P(Y) \sum_X P(X|Y) [-\log P(X|Y)] = \sum_Y P(Y) \sum_X P(X) [-\log P(X)] = H(X)$$

- Joint entropy of X and Y is the sum of the entropies of X and Y if they are independent

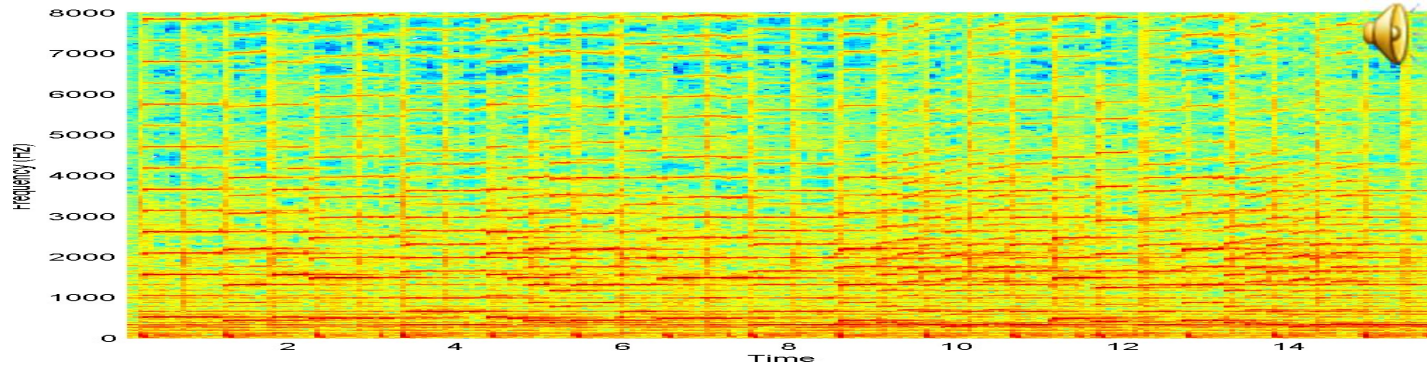
$$H(X, Y) = \sum_{X, Y} P(X, Y) [-\log P(X, Y)] = \sum_{X, Y} P(X, Y) [-\log P(X)P(Y)]$$

$$= -\sum_{X, Y} P(X, Y) \log P(X) - \sum_{X, Y} P(X, Y) \log P(Y) = H(X) + H(Y)$$

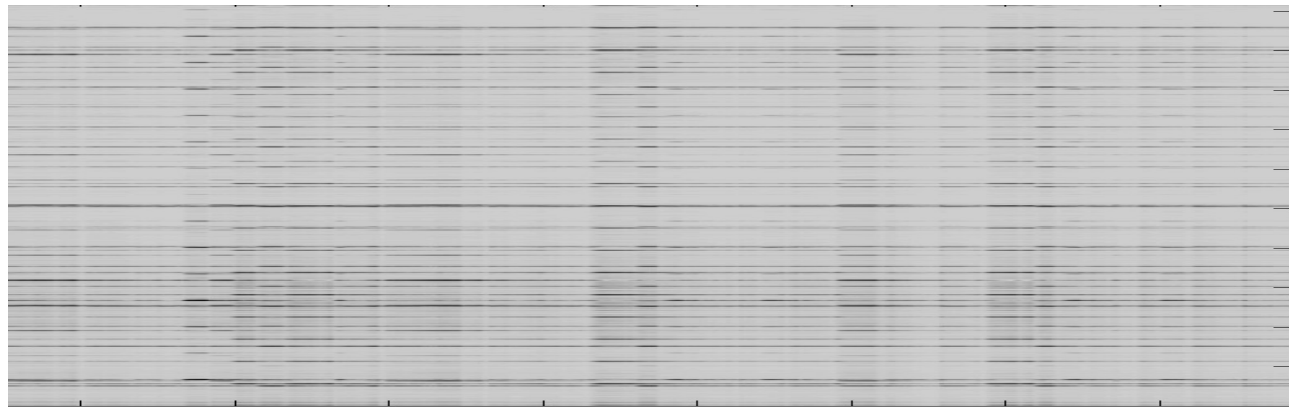
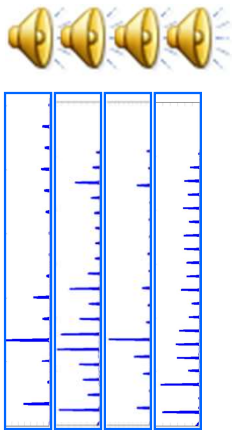
Onward..

Projection: multiple notes

M =



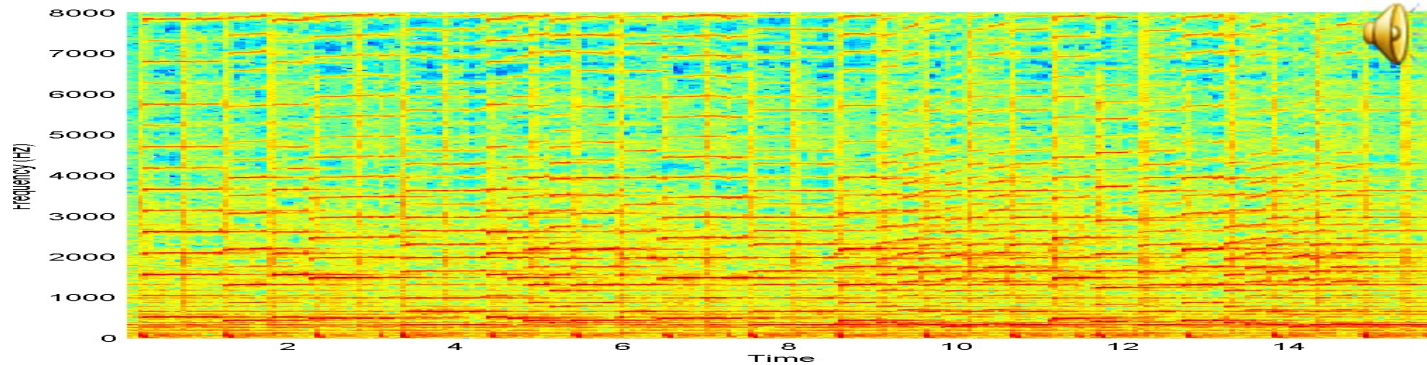
W =



- $\mathbf{P} = \mathbf{W} (\mathbf{W}^T \mathbf{W})^{-1} \mathbf{W}^T$
- Projected Spectrogram = $\mathbf{P}\mathbf{M}$

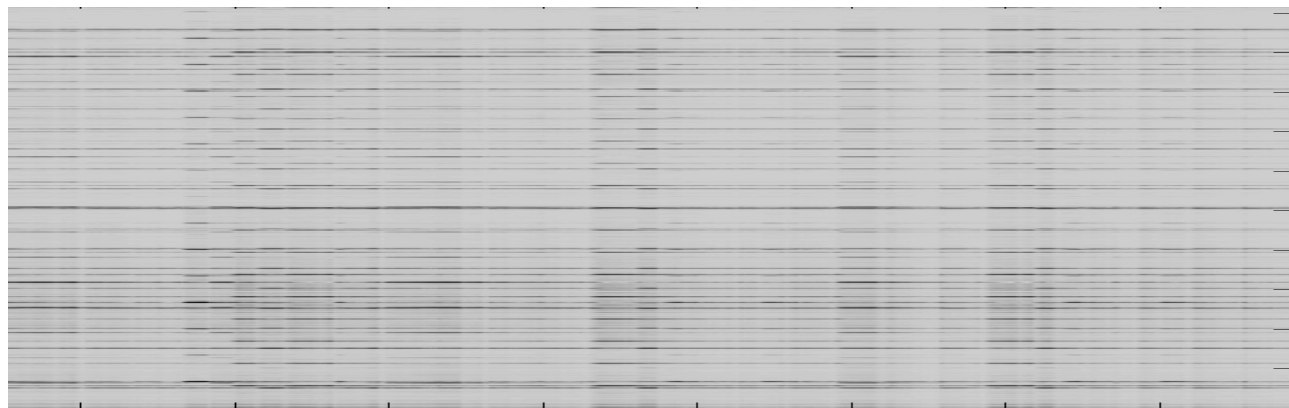
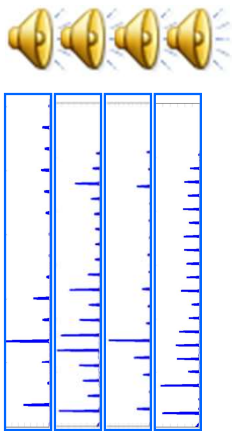
We're actually computing a score

$M =$



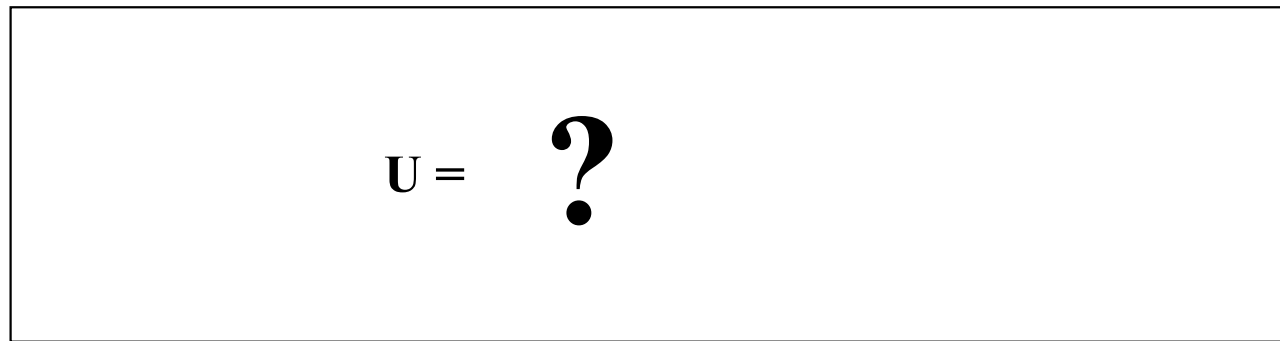
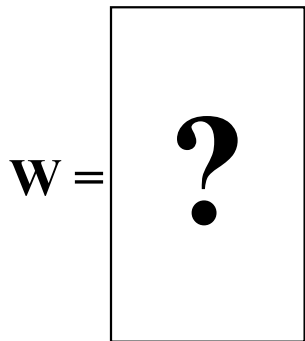
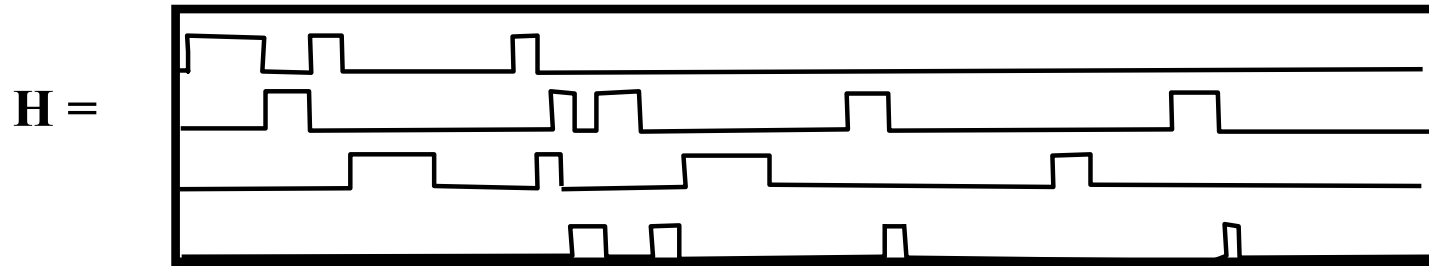
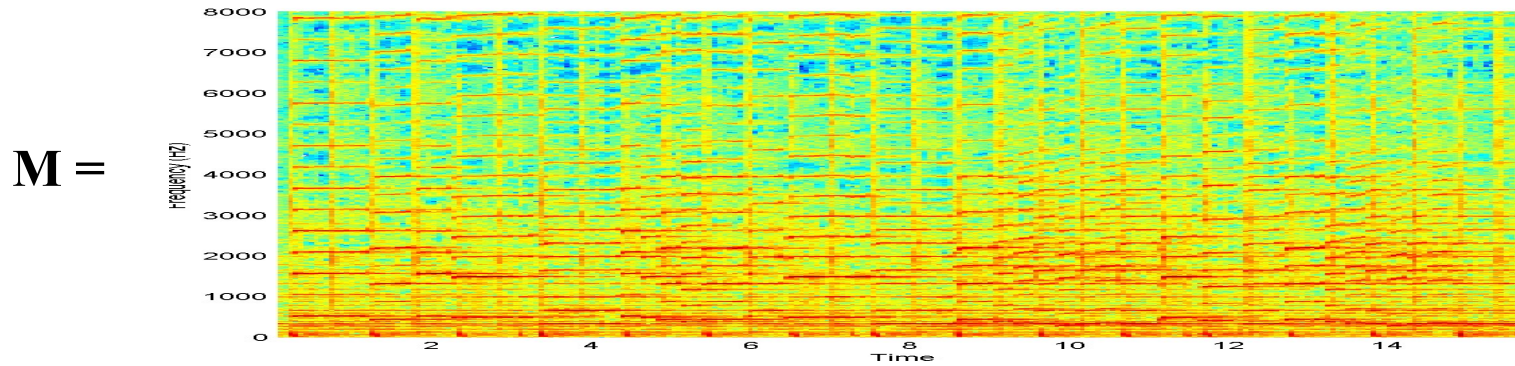
$H = ?$

$W =$



- $M \sim WH$
- $H = \text{pinv}(W)M$

How about the other way?



■ $M \sim WH$

$W = M \text{pinv}(H)$

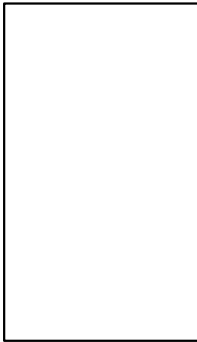
$U = WH$

When both parameters are unknown

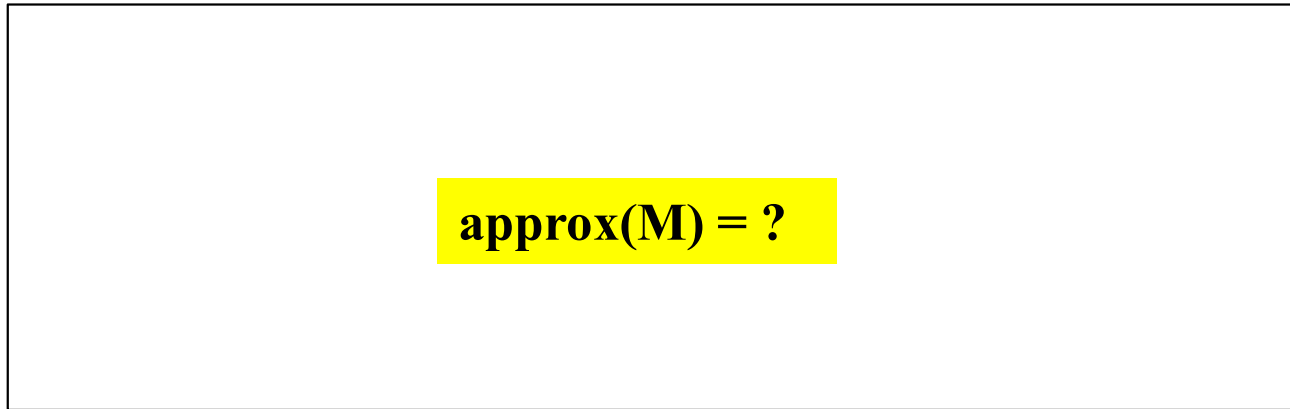
H = ?



W = ?



approx(M) = ?



- Must estimate both **H** and **W** to best approximate **M**
- Ideally, must learn *both* the *notes* and *their* transcription!

A least squares solution

$$\mathbf{W}, \mathbf{H} = \arg \min_{\overline{\mathbf{W}}, \overline{\mathbf{H}}} \|\mathbf{M} - \overline{\mathbf{W}}\overline{\mathbf{H}}\|_F^2 + \Lambda(\overline{\mathbf{W}}^T \overline{\mathbf{W}} - \mathbf{I})$$

- Constraint: \mathbf{W} is orthogonal
 - $\mathbf{W}^T \mathbf{W} = \mathbf{I}$
- The solution: \mathbf{W} are the Eigen vectors of $\mathbf{M}\mathbf{M}^T$
 - PCA!!
- $\mathbf{M} \sim \mathbf{W}\mathbf{H}$ is an approximation
- Also, the rows of \mathbf{H} are *decorrelated*
 - Trivial to prove that $\mathbf{H}\mathbf{H}^T$ is diagonal

PCA

$$\mathbf{W}, \mathbf{H} = \arg \min_{\overline{\mathbf{W}}, \overline{\mathbf{H}}} \|\mathbf{M} - \overline{\mathbf{W}\mathbf{H}}\|_F^2$$

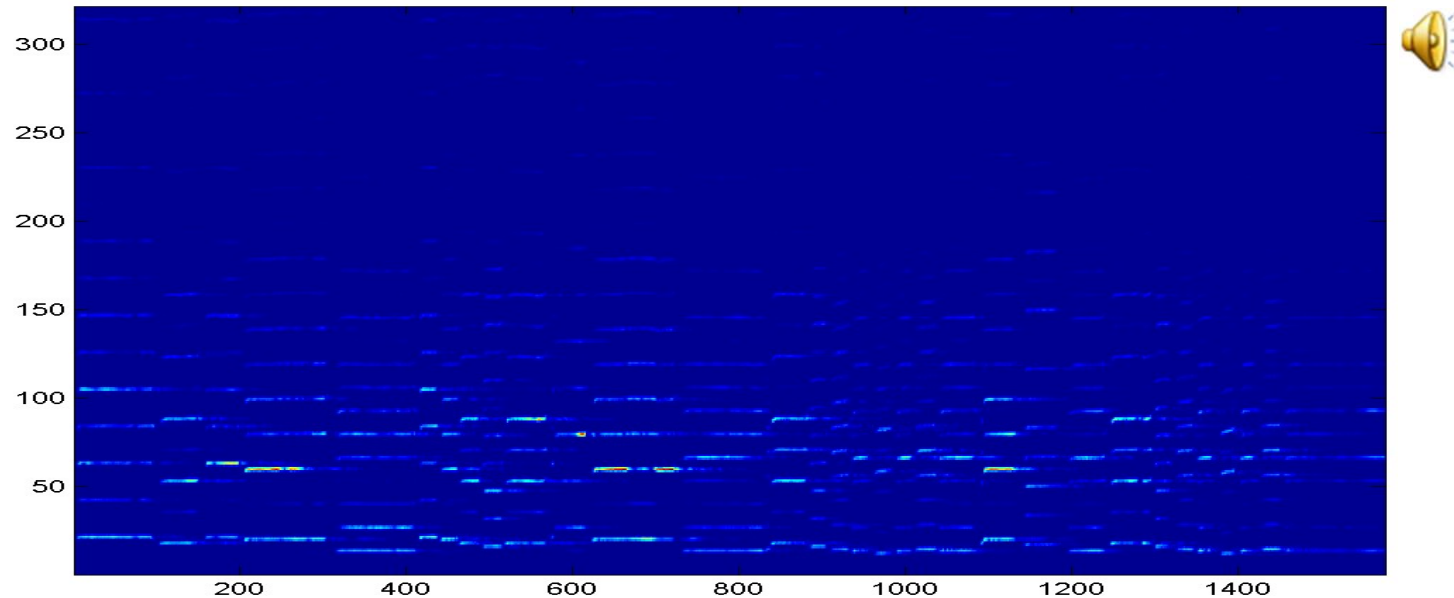
$$\mathbf{M} \approx \mathbf{W}\mathbf{H}$$

$$\mathbf{W}\mathbf{W}^T = \text{Diagonal} \quad \text{OR} \quad \mathbf{H}\mathbf{H}^T = \text{Diagonal}$$

The conditions are equivalent

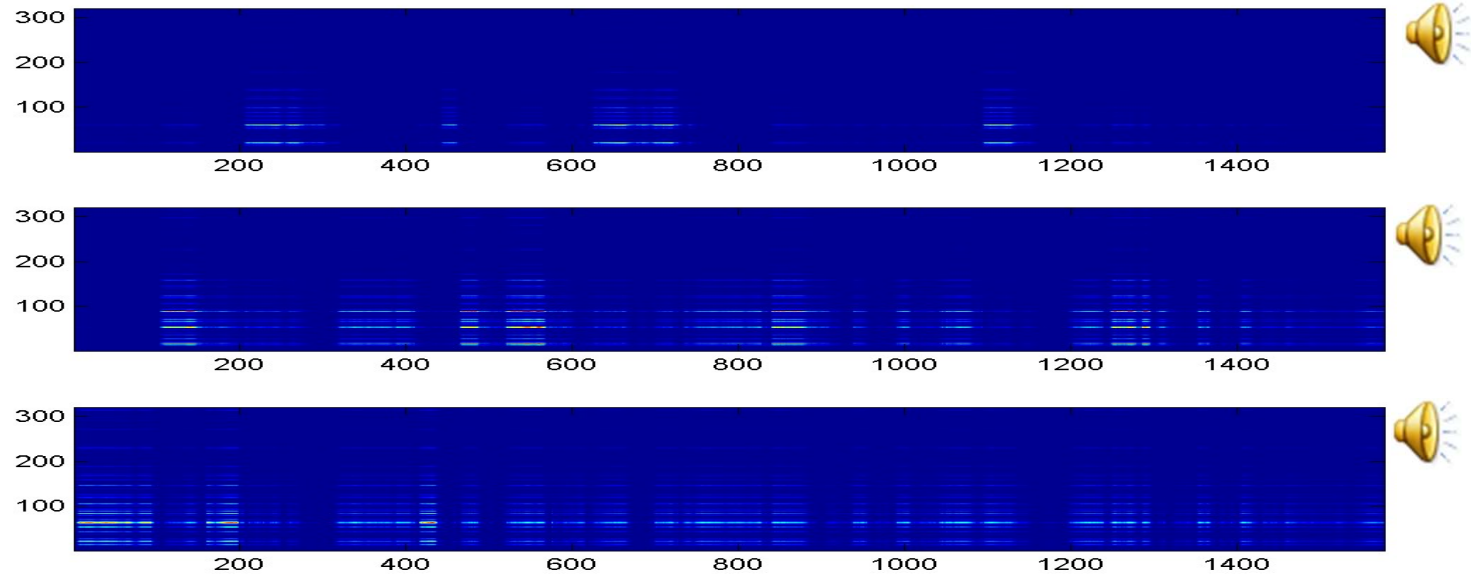
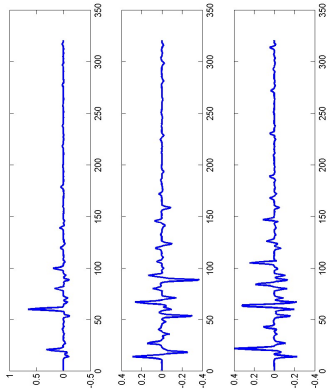
- The columns of \mathbf{W} are the bases we have learned
 - The linear “building blocks” that compose the music
- They represent “learned” notes
 - $\mathbf{w}_i \mathbf{h}_i$ is the contribution of the i th note to the music
 - \mathbf{w}_i is the i th column of \mathbf{W}
 - \mathbf{h}_i is the i th row of \mathbf{H}

So how does that work?



- There are 12 notes in the segment, hence we try to estimate 12 notes..

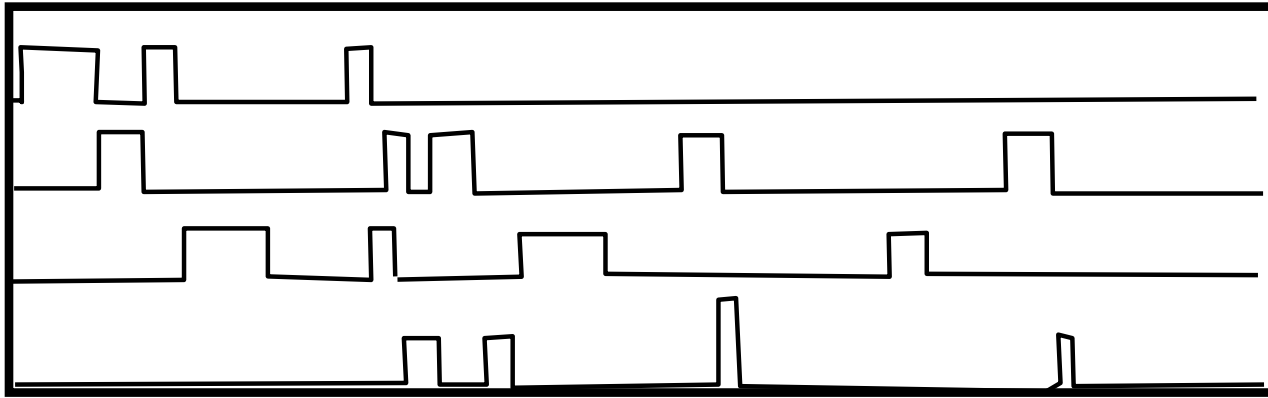
So how does that work?



- There are 12 notes in the segment, hence we try to estimate 12 notes..
- Results are not good

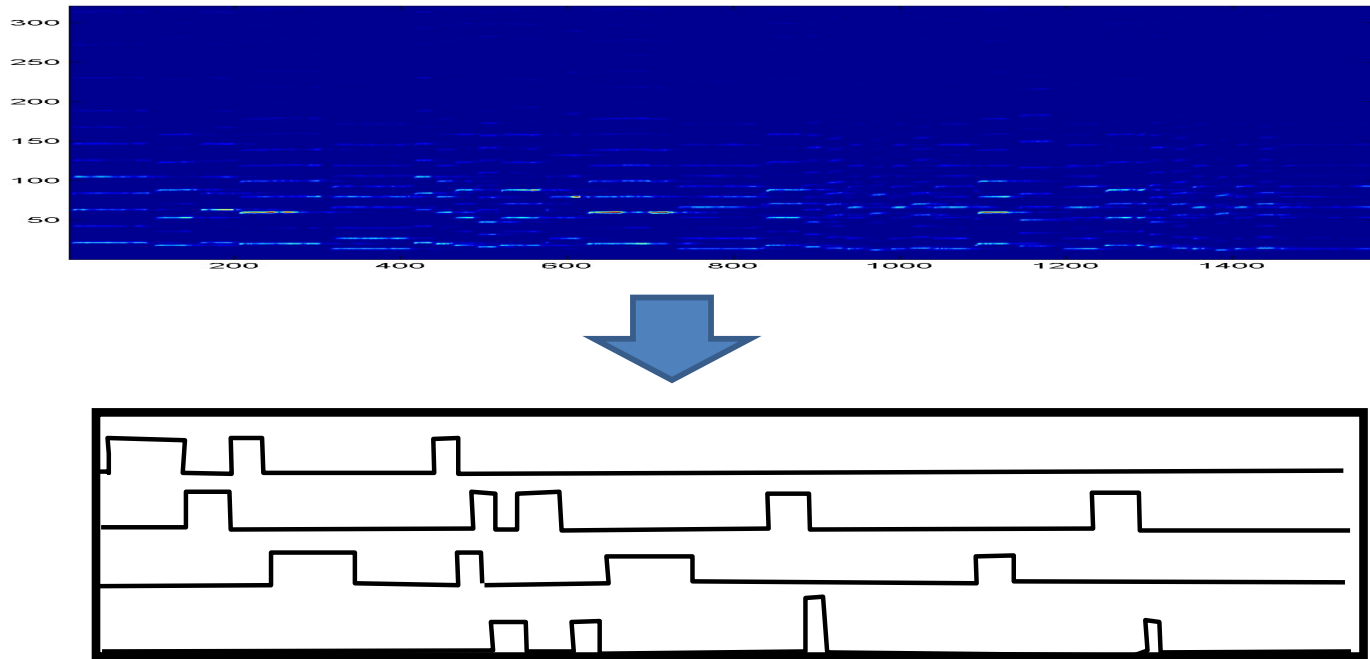
PCA through decorrelation of notes

$$\mathbf{W}, \mathbf{H} = \arg \min_{\overline{\mathbf{W}}, \overline{\mathbf{H}}} \|\mathbf{M} - \overline{\mathbf{H}}\|_F^2 + \Lambda(\overline{\mathbf{H}}\overline{\mathbf{H}}^T - \mathbf{D})$$



- Different constraint: Constraint \mathbf{H} to be decorrelated
 - $\mathbf{H}\mathbf{H}^T = \mathbf{D}$
- This will result exactly in PCA too
- Decorrelation of \mathbf{H} Interpretation: What does this mean?

Decorrelation



- Alternate view: Find a matrix \mathbf{B} such that the rows of $\mathbf{H}=\mathbf{B}\mathbf{M}$ are uncorrelated
- Will find $\mathbf{B} = \mathbf{W}^T$
- \mathbf{B} is the *decorrelating matrix* of \mathbf{M}

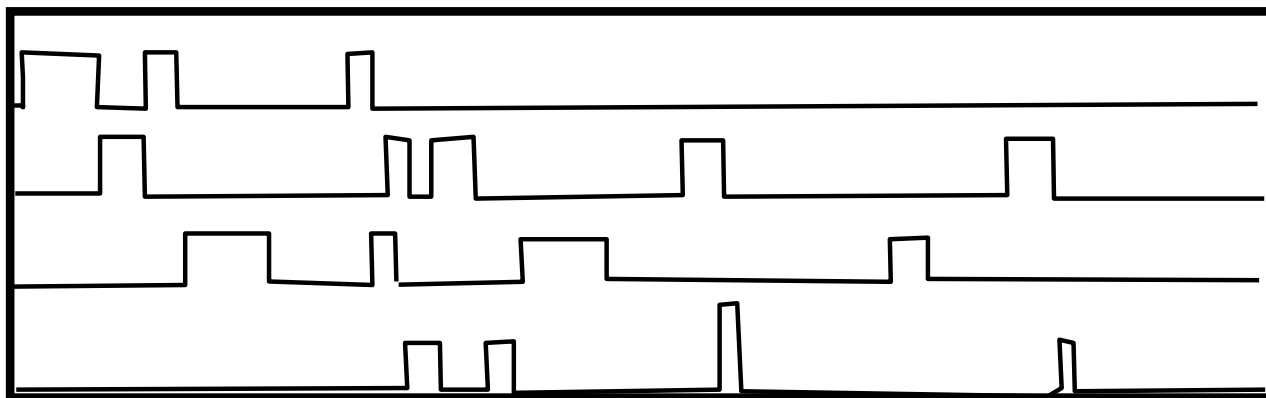
Poll 4

- Decorrelation is a sufficient criterion for determining semantically meaningful bases
 - T
 - F
- Non-negativity is a sufficient criterion for determining semantically meaningful bases
 - T
 - F

Poll 4

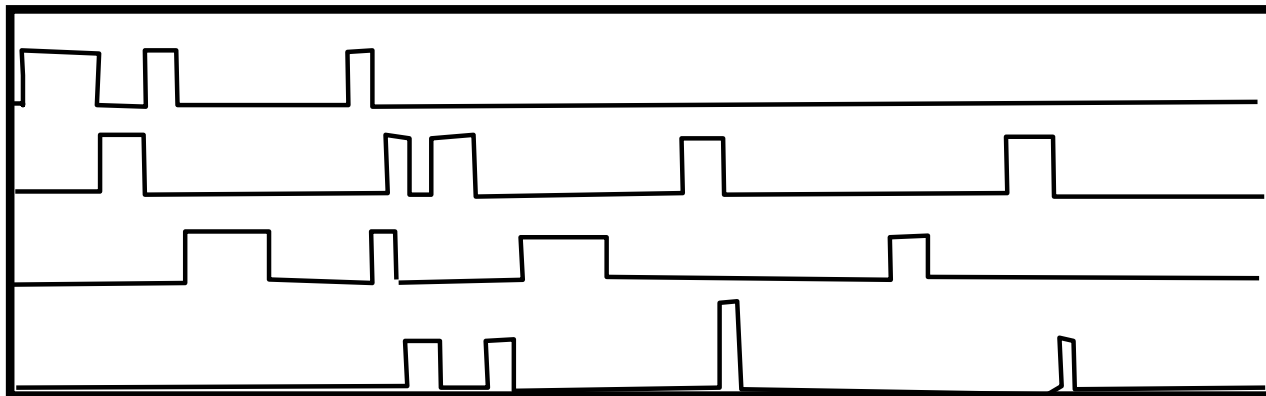
- Decorrelation is a sufficient criterion for determining semantically meaningful bases
 - T
 - **F**
- Non-negativity is a sufficient criterion for determining semantically meaningful bases
 - T
 - **F**

What *else* can we look for?



- Assume: The “transcription” of one note does not depend on what else is playing
 - Or, in a multi-instrument piece, instruments are playing independently of one another
- Not strictly true, but still..

What *else* can we look for?



- Assume: The “transcription” of one note does not depend on what else is playing
 - Or, in a multi-instrument piece, instruments are playing independently of one another
- **Attempting to find statistically independent components of the mixed signal**
 - ***Independent Component Analysis***

Formulating it with Independence

$$\mathbf{W}, \mathbf{H} = \arg \min_{\overline{\mathbf{W}}, \overline{\mathbf{H}}} \|\mathbf{M} - \overline{\mathbf{W}\mathbf{H}}\|_F^2 + \Lambda(\text{rows of } \mathbf{H} \text{ are independent})$$

- Impose statistical independence constraints on decomposition

Next Class

- Independent Component Analysis
- By Adnan Yunus